Pronouns, Second Edition (Python Version)

Kelly Roach kellybrianroach@outlook.com https://www.kellyroach.com

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Preface

Pronouns, Second Edition is a 2024 LaTeX-formatted version of the author's original 1980 Caltech M.S. thesis, *Pronouns* [27]. The content of *Pronouns, Second Edition* is substantially the same as the original *Pronouns* with the following principal differences:

- OCR'd content of *Pronouns* converted to modern LaTeX style.
- Misspellings, minor grammar points, numberings, and minor technical points are corrected.
- Capitalization changes.
- Figure captions shortened.
- PEP 8 Coding Style identifier spellings.
- Tikz figures replace partially hand-drawn TXT figures.
- LaTeX tabular tables replace TXT tables.
- LaTeX References replace TXT References.
- LaTeX Index added.
- Preface added.

Caltech's "Usage Policy", inherited from the original Pronouns, states:

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Postscript and PDF versions of the original *Pronouns* are also available on the author's PLANETQUANTUM.COM website [26].

The author's M.S. thesis and undergraduate adviser was Frederick B. Thompson [34].

Introduction

Certain substitutions and abbreviations occur in English which are not well understood yet that we would like to understand better so that we may implement them in computer natural language systems intended for man-machine communication. These include **pronouns** and other **function words** like those below in Figure 0.1 acting both in isolation and with each other.

| Ι | me | my | myself | mine | we |
|------------|----------|------------|-----------|---------|--------|
| us | our | ours | ourselves | you | your |
| yours | yourself | yourselves | he | him | his |
| himself | she | her | hers | herself | it |
| its | itself | they | them | their | theirs |
| themselves | this | that | these | those | one |
| ones | oneself | other | others | all | none |
| some | any | each | which | what | who |
| whom | whose | another | do | does | did |
| done | doing | SO | | | |

Figure 0.1. Pronouns and Other Function Words

As well, we have noun phrases modified by **demonstratives**, Head Deletion, and Equi-NP Deletion. Bloomfield [2] defined **substitution** as a replacement operation.

A substitute is a linguistic form or grammatical feature which, under certain conventional circumstances, replaces any one of a class of linguistic forms. Thus, in English, the substitute I replaces any singular-number substantive expression, provided that this substantive expression denotes the speaker of the utterance in which the substitute is used.

In this thesis we will be concerned with pronouns. Possibly because this will be the only chance we get, we should note the wide variety of substitution mechanisms in general. Examples (0.2)-(0.11) are from Sag [30].

(0.2) Do It Anaphor

Jerry won't prove that theorem; Alice will do it. [do it = prove that theorem]

(0.3) Sentential It Anaphor

I believe that she means business and you'd better believe it too. [it = that she means business]

(0.4) Null Complement

They asked me to leave but I refused ϕ . $\phi = \text{to leave}$

(0.5) **Ones Pronominalization**

Betsy has a blue car, and Randy has a red one. [one = car]

(0.6) Verb Phrase Deletion

Joan wouldn't eat a Quarter Pounder, but Annie would ϕ . $[\phi = \text{eat a Quarter Pounder}]$

(0.7) Sluicing

Someone has drunk my entire six-pack of Schlitz Light, but I don't know who ϕ .

 $\phi = has drunk my entire six-pack of Schlitz-Light$

(0.8) Stripping

Gwendolyn snorts cocaine, but ϕ_1 not ϕ_2 in her own apartment. $[\phi_1 = \text{Gwendolyn (does)}, \phi_2 = \text{snort cocaine}]$

(0.9) Gapping

Erichman duped Haldeman and Nixon ϕ Mitchell. $[\phi = duped]$

(0.10) Conjunction Reduction

Mitchell lied to the committee and ϕ was sentenced last year. $[\phi = \text{Mitchell}]$

(0.11) So Anaphor

Mitchell said he was innocent and Nixon said so too. [so = he was innocent]

To this list we can add pronominalizations. Examples (0.12)-(0.14) are from Lees and Klima [22].

(0.12) Reflexive Pronominalization Mary's father supported himself. [himself = Mary's father]

(0.13) **Pronominalization** Mary's father supported her.

[her = Mary]

(0.14) **Reciprocal Pronominalization**

John and Mary kissed each other. [each other = John and Mary]

And we might add (0.15) and (0.16) as well.

(0.15) Head Deletion Joan's cat pures but Mary's ϕ doesn't. $[\phi = \text{cat}]$

(0.16) Equi-NP Deletion

John is a fraid of ϕ cutting himself. $[\phi = \text{John's}]$

Clearly, this list starts to grow very large with addition or refinement and it is probably safe to say that many volumes could be written on substitution processes without putting it to bed. This thesis is about pronouns and chaining of pronouns, and so is much narrower in scope. But this is not much comfort if the goals are not clearly in sight. We are just as lost in the middle of Lake Michigan as we are in the middle of the Pacific Ocean if we don't have a horizon to steer us by.

Part of the problem with investigations of anaphora today is that there is no horizon to steer by. Even though work on anaphora continues in an intelligent way, little progress is being made towards a really comprehensive theory. Instead we have a lot of scattered and independent results.

One goal of this thesis, besides talking about pronouns, is to seek out an algorithmic framework on which to build theory. Accordingly, various data structures such as nodes, C-S-N trees, and chaining tables are created for this purpose. Hopefully, the reader will recognize these data structures as too simplistic and will be moved to improve upon them. This thesis is, by no means at all, a solution to pronouns. At best, it may be a small compass in the middle of Lake Michigan, but this is our approach.

1 Fundamentals

1.1 Introduction

This chapter describes notation and basic ideas that will be used throughout this thesis. Hopefully, most of the notation described in this chapter is already familiar to the reader, but if not, then this chapter should be self-contained enough to be understandable by a reader with less experience.

1.2 Sentences

Sentences are numbered and are kept separate from the text of discussion for ease of reference. For example, (1.1) is from Huddleston [13] and is an example of a **Bach Peters sentence**.

(1.1) The boy who was fooling her kissed the girl who loved him.

Ungrammatical sentences are prefixed with an asterisk (*) and **sentences of questionable grammaticality** are prefixed with a question mark (?). Here, (1.3) is from Chomsky [5].

- (1.2) *John killed herself.
- (1.3) ?Colorless green ideas sleep furiously.

Subscripts are used to indicate identity between constituents, meaning roughly that they mean the same thing or denote the same referent. More properly, we may think of constituents having the same subscript as being chained together. Below, (1.4) and (1.5) are from Bresnan [3].

- (1.4) Some students think they are smarter than they are.
- (1.5) *Some studentsI think some studentsI are smarter than some studentsI are.

Sometimes we enclose information in **brackets** at the beginning or end of a sentence. This same notation is also sometimes used as an alternative to subscripts in identifying constituents. Here, (1.6) is from Bresnan [3], (1.7) and (1.8) are from Roberts [28] and (1.9) is from Bloom and Hayes [1].

- (1.6) My uncle has never ridden a camel but his brother has, although it was lame. [it = camel]
- (1.7) Men are mortal. [All men are mortal]
- (1.8) Men are waiting. [Some men are waiting]
- (1.9) [Seeing a picture of John Smith] That's John Smith.

A deletion site is indicated by a ϕ . Example (1.10) is from Hockett [12].

(1.10) I like the fresh candy better than the stale ϕ . $[\phi = \text{candy}]$

Deletion sites arising from transformations like **Equi-NP Deletion** are treated similar to pronouns in this paper. Although there are many different kinds of deletion sites with distinct properties, we won't pay attention to this distinction in this thesis.

The symbol = is used between sentences to indicate that they are equivalent, while the symbol \neq is used between sentences to indicate that they are not equivalent. Below, (1.11)-(1.14) are from Ross [29].

- (1.11) If John can, he will do it. =
- (1.12) If he can, John will do it.
- (1.13) John will do it if he can. \neq
- (1.14) He will do it if John can.

1.3 Noun Phrases

Quantified noun phrases are **noun phrases** modified by **quantifiers**. Examples (1.15)-(1.18) are quantified noun phrases.

- (1.15) all female astronauts
- (1.16) at least 10 sexual perverts
- (1.17) many notorious criminals
- (1.18) nearly a dozen Unicorns

Genitives are possessive noun phrases. Examples (1.19)-(1.22) are genitives.

- (1.19) Uncle Iggy's
- (1.20) my cobra's
- (1.21) the Nazi war criminal's
- (1.22) the alien creatures'

A noun phrase can be **generic**, **specific**, or **nonspecific**, indicated respectively by (1.23)-(1.25) from Kuno [16].

- (1.23) A cat is a malicious animal. [generic]
- (1.24) I have a cat at home, but hate it. [specific]
- (1.25) I want to get a cat for myself. [nonspecific]

A plural noun phrase can be **collective** or **distributive**. Examples (1.26)-(1.28) are from Fauconnier [7].

- (1.26) The men gathered. [collective]
- (1.27) The men took off their hats. [distributive]
- (1.29) The men carried the couch. [ambiguous]

Sentence (1.29) is **ambiguous** because it can mean either (1.30) or (1.31).

- (1.30) Each man of the men carried the couch.
- (1.31) The team of men carried the couch.

Smith [33] has also noticed this distinction. This explains why (1.32)-(1.35) below are ambiguous.

- (1.32) John and Mary bought the new book by John Steinbeck.
- (1.33) Bricks and stones make strong walls.
- (1.34) George and Marmaduke have dogs.
- (1.35) Gerry likes ice cream and cake.

1.4 Pronouns

Pronouns are cross-classified by person, plural, gender, animate, reflexive, attributive possessive, and predicative possessive features among others. **First person pronouns** are given in (1.36).

(1.36) I, me, myself, my, mine, we, us, our, ours, ourselves.

Second person pronouns are given in (1.37).

(1.37) you, yourself, yourselves, your, yours

Third person pronouns are given in (1.38).

(1.38) she, he, it, they, her, him, them, herself, himself, itself, themselves, his, its, their, hers, theirs

Singular pronouns are given in (1.39).

(1.39) I, me, myself, my, mine, you, yourself, your, yours, she, he, it, her, him, herself, himself, itself, his, its, hers

Plural pronouns are given in (1.40).

(1.40) we, us, our, ours, ourselves, you, yourselves, your, yours, they, them, themselves, their, theirs

Pronouns with female gender are given in (1.41).

(1.41) she, her, herself, hers

Pronouns with male gender are given in (1.42).

(1.42) he, him, himself, his

Animate pronouns are given in (1.43).

(1.43) I, me, myself, mine, you, yourself, yourselves, your, yours, she, he, they, her, him, herself, himself, themselves, his, their, hers, theirs

Inanimate pronouns are given in (1.44).

(1.44) it, they, them, itself, themselves, its, theirs

Reflexive pronouns are given in (1.45).

(1.45) myself, yourself, yourselves, herself, himself, itself, themselves

Attributive possessive pronouns are given in (1.46).

(1.46) my, your, her, his, its, their

Predicative possessive pronouns are given in (1.47).

(1.47) mine, yours, hers, his, its, theirs

Besides the pronouns given above, we also have ones pronouns and reciprocal pronouns. **Ones pronouns** are given in (1.48).

(1.48) one, oneself, one's

Reciprocal pronouns are given in (1.49).

(1.49) each other, one another, each other's, one another's

1.5 Features

We use three kinds of **features** in this thesis. The symbol + indicates presence of a feature. The symbol – indicates absence of a feature. And the symbol ? indicates that the presence or absence of a feature is either unspecified or not applicable. In the coming chapters, we will speak of agreement of features. A ? feature agrees with any other feature. The only time two features do not agree is when we are comparing a + and a – feature. Using = to indicate agreement and \neq to indicate nonagreement, we have Figure 1.50.

Figure 1.50. Agreement and Nonagreement between Features

1.6 Parse Trees

Sentence **parse trees** are only drawn schematically in this thesis as extra detail is unnecessary. Parse trees shown more or less represent the surface structure of a sentence. Clause dominating nodes are labelled S and clause conjoining nodes are labelled C. In this thesis, genitives and adjectives are not treated as arising from transformations, but as occuring in the base component. Below, example (1.51) is from Huddleston [13] and example (1.52) is from Grosu [9].

(1.51) The man who lives next door said that he would mow my lawn.



(1.52) Somebody seduced Bill's sister, but no one will ever seduce Jack's and she knows it.



1.7 Clauses

Adverbial clauses are clauses beginning with an adverb. Some examples are (1.53)-(1.57) below.

- (1.53) after Fido made a mess on the carpet
- (1.54) before George kisses Betty
- (1.55) since John is an asshole
- (1.56) until Cathy behaves herself
- (1.57) although Lile flunked all his classes

Clauses complemented with <u>that</u> are **that clauses**. Example (1.58) is a that clause.

(1.58) that Snoopy is a cat

Clauses modified by the **For-To Transformation** are **infinitive clauses**. Example (1.59) is an infinitive clause.

(1.59) for Ruth to choose

Clauses modified by the **Possessive-Ing Transformation** are **genitive clauses**. Example (1.60) is a genitive clause.

(1.60) Mary's kissing Bob

Clauses modified by **WH-Fronting Transformation** but not the **Question Transformation** and which modify noun phrases are **relative clauses**. Examples (1.61)-(1.65) are relative clauses.

- (1.61) who ate five hamburgers
- (1.62) that has a leaky faucet
- (1.63) which doesn't run
- (1.64) whom he gave it to

(1.65) whose life isn't worth a postage stamp

Clauses without embedded subordinate clauses are simplex. In example (1.66) from Ross[29], the simplexes are (1.67)-(1.69). In example (1.70), from Huddleston [13], the simplexes are (1.71)-(1.73). In example (1.74), from Huddleston, the simplexes are (1.75)-(1.77).

(1.66) Realizing that he was unpopular didn't disturb Oscar.



- $\begin{array}{l} \phi \ precedes \ he \\ \phi \ commands \ he \\ \phi \ precedes \ Oscar \\ Oscar \ commands \ \phi \\ Oscar \ commands \ he \end{array}$
- (1.67) S didn't disturb Oscar
- (1.68) ϕ 's realizing that S
- (1.69) he was unpopular
 - (1.70) My neighbor who is pregnant said that she was very happy.





- (1.71) my neighbor said that S
- (1.72) who is pregnant
- (1.73) she was very happy

(1.74) The pilot who shot at it hit the Mig that chased him.



pilot precedes him pilot commands him it precedes the Mig Mig commands it

- (1.75) the pilot hit the Mig
- (1.76) who shot at it
- (1.77) that chased him

1.8 Precedes and Commands

The *precedes* and *commands* relations, first described by Langacker [19], are defined below in (1.78) and (1.79).

(1.78) precedes Relation

A node A precedes another node B if

- (a) neither A nor B *dominates* the other, and
- (b) A occurs before B (in preorder traversal)

(1.79) commands Relation

A node A *commands* another node B if

- (a) neither A nor B *dominates* the other, and
- (b) the S-node that most immediately dominates A also dominates B

Another relation that will be useful is the *is separate from* relation defined below in (1.80).

(1.80) is separate from Relation

- A node R is separate from another node B if
- (a) neither A nor B dominates the other, and
- (b) the lowest node in the tree dominating A and B is a C-node.

We will see that *precedes*, *commands*, and *is separate from* are useful in determining when pronominalization is or isn't possible.

In example (1.81), A precedes B, A commands B, and E commands A. We don't have A precedes A, B precedes A, B precedes B, A commands A, or B commands B.

(1.81)



A precedes B A commands B B commands A

In (1.82), A precedes B, A is separate from B, and B is separate from A. In (1.83), A precedes B and A commands B. In (1.84), A precedes B and B commands A. In (1.85), A precedes B, H is separate from B, and B is separate from A.

(1.82)



A precedes B A is separate from B A precedes B

(1.83)



A precedes B A commands B

(1.84)



A precedes B B commands A

(1.85)



A precedes B A is separate from B B is separate from A

Examples (1.86)-(1.89) are from Langacker [19].

(1.86) The mosquito which bit Algernon was killed by him. [him = Algernon]



Algernon *precedes* him him *precedes* Algernon (1.87) The mosquito which bit him was killed by Algernon. [him = Algernon]



him *precedes* Algernon Algernon *precedes* him

(1.88) Algernon killed the mosquito which bit him. [him = Algernon]



Algernon *precedes* him him *commands* Algernon

(1.89) He killed the mosquito which bit Algernon. $[he \neq Algernon]$



he *precedes* Algernon Algernon *commands* he

The Precedes and Commands Rule, essentially as stated by Langacker [19], is given in (1.90) below.

(1.90) Precedes and Commands Rule

A pronoun P may be used to pronominalize a noun phrase NP unless

(a) P precedes NP, and

(b) P commands NP or P is separate from NP

Note that the Precedes and Commands Rule explains the grammaticality and ungrammaticality of (1.86)-(1.89). These further examples from Ross [29] should drive the point home.

- (1.91) After John Adams woke up, he was hungry. [he = John Adams]
- (1.92) That Oscar was unpopular didn't disturb him. [him = Oscar]
- (1.93) For your brother to refuse to pay taxes would get him into trouble. [him = your brother]
- (1.94) Anna's complaining about Peter infuriated him. [him = Peter]
- (1.95) The possibility that Fred will be unpopular doesn't bother him. [him = Fred]



NP precedes P P commands NP

- (1.96) After he woke up, John Adams was hungry. [he = John Adams]
- (1.97) That he was unpopular didn't disturb Oscar. [he = Oscar]
- (1.98) For him to refuse to pay taxes would get your brother into trouble. [him = your brother]
- (1.99) Anna's complaining about him infuriated Peter. [him = Peter]
- (1.100) The possibility that he will be unpopular doesn't bother Fred. [him = Fred]



P precedes NP NP commands P

(1.101) John Adams was hungry after he woke up. [he = John Adams]

- (1.102) Oscar wasn't disturbed that he was unpopular. [he = Oscar]
- (1.103) It would get your brother into trouble for him to refuse to pay taxes. [him = your brother]
- (1.104) Peter was infuriated at Anna's complaining about him. [him = Peter]
- (1.105) Fred isn't bothered by the possibility that he will be unpopular. [he = Fred]



NP precedes P NP commands P

- (1.106) *He was hungry after John Adams woke up. [he = John Adams]
- (1.107) *He wasn't disturbed that Oscar was unpopular. [he = Oscar]
- (1.108) *It would get him into trouble for your brother to refuse to pay taxes. [him = your brother]
- (1.109) *He was infuriated at Anna's complaining about Peter. [he = Peter]
- (1.110) *He isn't bothered by the possibility that Fred will be unpopular. [he = Fred]



P precedes NP P commands NP

Examples (1.111) and (1.112) from Langacker [19] illustrate the Precedes and Commands Rule for **conjoined structures**.

(1.111) Penelope cursed Peter and slandered him. [him = Peter]



Peter precedes him Peter is separate from him him is separate from Peter

(1.112) *Penelope cursed him and slandered Peter. [him = Peter]



him precedes Peter him is separate from Peter Peter is separate from him

Examples (1.113) and (1.114) adapted from Chiba $\left[4\right]$ involve Equi-NP Deletion.

(1.113) The interest in visiting Las Vegas that Mary displayed is typical of gamblers.



 $\phi~precedes$ Mary

2 Resolution Module

2.1 Introduction

In the previous chapter we touched upon some basic notions such as the *precedes*, *commands*, and *is separate from* relations. We will see in the coming chapters how these concepts give rise to a very promising approach to the problem of pronoun resolution.

The algorithm we shall describe won't be complete in the sense that we will elaborate and refine it in later chapters and after we are done it will need elaboration and refinement, but it will be set in firm soil so that we have a foundation on which to build. Because personal and reflexive pronouns are easiest, these are the pronouns we shall consider first. But before we go any farther, let us take time out to indicate something of the environment and structure of the module that does resolving of pronouns in a natural language system, the Resolution module.

2.2 Environment

The center of a natural language system is the Language Processor module which is divided into five submodules. These are the Language Driver, Preprocessor, Parser, Semantic Processor, and Output Processor as indicated in Figure 2.1.



Figure 2.1. Submodules of the Language Processor

Briefly, from the point of view of the Language Processor, the following happens. A user types input at a terminal which is picked up by the Operating System of the natural language system. The Operating System maintains information about the user including the language version he is in as well as his state in that version. The user's state is known as his prefix. The Operating System, after picking up a user's input calls a Process Input routine of the Language Driver in the Language Processor. Once in the Language Driver, the first module to be called upon is the Preprocessor.

The Preprocessor in the Language Processor compresses blanks in the input string, straps right and left delimiters about it, recognizes and builds parsing graph arcs over identifiers and numbers, and looks the identifiers up in the lexicon. After calling the Preprocessor, the Language Driver calls the Parser.

The Parser in the Language Processor parses the output. of the Preprocessor using an algorithm such as the **Kay algorithm** and can handle any general rewrite rule grammar. Of course, since a sentence may be ambiguous, more than one system parse tree may be passed back by the Parser. If no good parsings are found, then the Syntax Diagnostics routine of the Syntax Diagnostics module of the natural language system is called. Otherwise, if there are good parsings, then the Language Driver calls the Semantic Processor on the output of the Parser.

The Semantic Processor is driven by the syntax of a system parse tree into making calls on semantic routines which can be postprocedures (called on their arguments after their arguments evaluate themselves), preprocedures (called on their arguments before their arguments evaluate themselves), and syntax procedures (called at syntax time during parsing before preprocedures and postprocedures are called during semantic processing). On return to the Language Driver, the Language Driver calls the Output Processor on the output of the Semantic Processor.

The Output Processor does some relatively menial processing such as removing duplicate lines from the output line list which will be sent back to the Operating System. The Output Processor is able to handle ambiguous output and removes diagnostic messages if at least one of the outputs is good.

On completion of the call on the Output Processor, the Language Driver returns to the Operating System and the Operating System displays the output line list on the user's terminal, at the same time updating its information on the user.

From the discussion of the *precedes*, *commands*, and *is separate from* relations in the previous chapter, we know that information about the syntax of the input sentence is critical to the resolving of pronouns in the input sentence. On the other hand, for semantic processing to carry out the processing it needs to carry out, the placing of information on the chaining of pronouns must already be placed in the system parse tree of the input sentence.

The logical conclusion of these two observations indicates that pronoun resolution takes place after parsing, but before semantic processing. This relationship of the Resolution module with the other modules of the Language Processor is indicated in Figure 2.2.



Figure 2.2. Resolution Module within the Language Processor

In practice, this formulation may not be quite correct because there can be other versions than English which will have nothing to do with the Pronoun Resolution module and so what we end up doing is making the Resolution module accessible via a semantic preprocedure which is associated with the parsing of the right delimiter of a sentence. So instead, what happens is that the first semantic preprocedure to be called will be the procedure which handles Pronoun Resolution.

2.3 Structure inside the Resolution Module

The Resolution module is partitioned into seven submodules besides a Global Declarations module. These are the Node Processor, Parser, Primary Utilities, Secondary Utilities, Table Processor, Table Interpreter, and Resolution Driver modules. The reader should not confuse the Parser of the Language Processor with the Parser of the Pronoun Resolution module which have entirely different functions. The relationship of these submodules of the Resolution module is indicated below in Figure 2.3.



Figure 2.3. Structure of the Resolution Module

Not shown is the Global Declarations module which does not have any procedures itself, but merely defines data structures. The Global Declarations submodule is accessible by all other submodules of the Resolution module.

3 Global Declarations

The Global Declarations module defines the data structures accessible to other modules within the pronoun resolution module. The Global Declarations module is shown below in Figure 3.1.

```
#globals.py
import sys
from typing import TextIO
from enum import Enum, IntEnum
from typing import Optional
class FeatureIndex(IntEnum): #Feature indices.
    PNF = 0 #Pronoun Feature
    FPF = 1 #First Person Feature
    SPF = 2 #Second Person Feature
    TPF = 3 #Third Person Feature
    PLF = 4 #Plural Feature
    GNF = 5 \#Gender Feature
    ANF = 6 #Animate Feature
    RPF = 7 #Reflexive Feature
    GEN = 8 #Genitive Feature
N_FEATURES = len(FeatureIndex) #Number of Features
class NodeId(Enum): #Identifies the type of node.
    C_NODE = 0 #Represents a C-node.
    S_NODE = 1 #Represents an S-node.
    N_NODE = 2 #Represents an N-node.
    E_NODE = 3 #Represents an E-node.
class Feature(Enum): #Linguistic feature values.
    PLUS = 0 \#Has this feature.
    MINUS = 1 #Doesn't have this feature.
    QUESTION = 2 #Might or might not have this feature.
Features = list[Feature] #list of Feature enums
```

```
Figure 3.1. Global Declarations Module (Part I)
```

```
class Node: #Base node class
    # Current tree
   _tree: Optional['Node'] = None
   def __init__(self):
        self.number: int = 0
        self.up_link: Optional['Node'] = None
        self.down link: Optional['Node'] = None
        self.left_link: Optional['Node'] = None
        self.right_link: Optional['Node'] = None
        self.thread_link: Optional['Node'] = None
        self.np_link: Optional['Node'] = None
        self.chain_link: Optional['Node'] = None
        self.col_link: Optional['Node'] = None
        self.ftr: Features = [Feature.QUESTION] * N_FEATURES
        self.id: NodeId = NodeId.C_NODE
        self.lit: str = ""
        self.end_col_link: Optional['Node'] = None
        self.pred_link: Optional['Node'] = None
        self.succ_link: Optional['Node'] = None
        self.sub: str = ' '
   @classmethod
   def tree(cls) -> Optional['Node']:
        return cls._tree
   @classmethod
   def set_tree(cls, new_tree: 'Node') -> None:
        cls. tree = new tree
```

Figure 3.1. Global Declarations Module (Part II)

Basically, our data structures are C-S-N trees, chaining tables, and the nodes they involve. It will help to get some feel for these data structures before we go on to other chapters.

3.1 Nodes

There are four kinds of **nodes**: C-nodes, S-nodes, N-nodes, and E-nodes. C-nodes, S-nodes, and N-nodes occur in C-S-N trees and correspond to conjoined structures, sentences, and noun phrases. E-nodes occur in chaining tables. The fields of the C-nodes, S-nodes, N-nodes, and E-nodes are as indicated in Figure 3.1.

3.2 C-S-N Trees

A C-S-N tree has three kinds of nodes: C-nodes, S-nodes, and N-nodes. Link fields which are relevant to C-S-N trees are up_link, down_link, left_link, right_link, thread_link, pred_link, and succ_link.
An example of a C-S-N tree is given in Figure 3.2.



Figure 3.2. C-S-N Tree

3.3 Chaining Tables

A chaining table contains N-nodes, E-nodes, and one S-node for keeping track of the chaining table. Link fields relevant to chaining tables are np_link, chain_link, col_link, end_col_link, pred_link, and succ_link. Chaining tables and C-S-N trees are connected through their N-nodes. An example of a chaining table is given in Figure 3.3.



Figure 3.3. Chaining Table

3.4 C-Nodes

A C-node has the following fields: up_link, down_link, left_link, right_link, thread_link, and number. C-nodes correspond to conjoined sentences and conjoined subordinate clauses.

3.5 S-Nodes

An **S-node** has exactly the same fields as a C-node and is only distinguished from a C-node by its NodeId. S-nodes correspond to sentences and subordinate clauses.

3.6 N-Nodes

An N-node has the following fields: lit, ftr, up_link, down_link, thread_link, np_link, chain_link, col_link, end_col_link, pred_link, succ_link, and number. N-nodes correspond to noun phrases without attached subordinate clause modifiers.

3.7 E-Nodes

An **E-node** has the following fields: sub, ftr, np_link, chain_link, and col_link. An E-node may be thought of as a copy of its np_link with a slightly more defined set of features.

3.8 lit Field

The lit field of an N-node is a string pointer to the string that the Nnode represents. The lit field is actually unnecessary in an N-node, but is convenient for displaying intermediate results. Function view_node_str of the Node Processor and some other procedures that display intermediate results use this field.

3.9 sub Field

The sub field of an E-node is a character representing the subscript of the E-node. The sub field of an E-node, like the lit field of an N-node, is an unnecessary field, but is convenient for displaying intermediate results.

3.10 ftr Field

The ftr field of an N-node or E-node is an array of Feature's representing the feature set of the N-node or E-node to which it corresponds. A Feature can be a PLUS, MINUS, or QUESTION as described in the previous chapter. The offsets PNF, FPF, SPF, TPF, PLF, GNF, ANF, and RPF are used to access elements of the ftr array. The accessed elements are pronoun feature, first person feature, second person feature, third person feature, plural feature, gender feature, animate feature, and reflexive feature. The number of Feature's is N_FEATURES. Figure 3.4 shows some examples of the settings of ftr for some typical noun phrases.

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| John | - | - | - | + | _ | _ | + | _ |
| flowers | - | - | - | + | + | ? | - | _ |
| he | + | - | - | + | _ | - | + | _ |
| them | + | - | - | + | + | ? | ? | - |
| I | + | + | - | - | - | ? | + | - |
| you | + | - | + | - | _ | ? | + | _ |
| her | + | - | - | + | - | + | + | - |
| myself | + | + | - | - | - | ? | + | + |
| herself | + | - | - | + | _ | + | + | + |
| itself | + | _ | - | + | _ | ? | - | + |

Figure 3.4. ftr Settings for Some Typical Noun Phrases

3.11 up_link Field

The up_link field of a C-node, S-node, or N-node links to the parent node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with up_link's shown is given in Figure 3.5.



Figure 3.5. C-S-N Tree with up_link's Shown

3.12 down_link Field

The down_link field of a C-node, S-node, or N-node links to the first child node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with down_link's shown is given in Figure 3.6.



Figure 3.6. C-S-N Tree with down_link's Shown

3.13 left_link Field

The left_link field of a C-node, S-node, or N-node links to the left brother node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with left_link's shown is given in Figure 3.7.



Figure 3.7. C-S-N Tree with left_link's Shown

3.14 right_link Field

The right_link field of a C-node, S-node, or N-node links to the right brother node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with right_link's shown is given in Figure 3.8.



Figure 3.8. C-S-N Tree with right_link's Shown

3.15 thread_link Field

The thread_link field of a C-node, S-node, or N-node links to the first node traversed after the C-node, S-node, or N-node in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with thread_link's shown is given in Figure 3.9.



Figure 3.9. C-S-N Tree with thread_link's Shown

3.16 number Field

C-node, S-node, or N-node have a number field which is the number that would be assigned to that node if the nodes of the C-S-N tree in which it occurs are numbered in a preorder traversal. An example of a C-S-N tree with number fields shown is given in Figure 3.10.



Figure 3.10. C-S-N Tree with number Fields Shown

3.17 np_link Field

For an E-node, the np_link is the N-node to which it is attached. Conceptually, we think of the E-node as being a copy of the N-node except for its subscript and different set of Feature's, chain_link, and col_link. The np_link is just a way of avoiding duplication of information. For an N-node, the np_link is always itself. An example of a chaining table with np_link's shown is given in Figure 3.11.



Figure 3.11. Chaining Table with up_link's Shown

3.18 chain_link Field

The chain_link of an E-node is another E-node representing the substitute to which the first E-node is attached. When chaining is obligatory, an N-node is chained to an N-node. An example of a chaining table with chain_link's shown is given in Figure 3.12.



Figure 3.12. Chaining Table with chain_link's Shown

3.19 col_link Field

The col_link field of an E-node or N-node links together the elements of a column in a table. An N-node is always on top of a column with E-nodes underneath. An example of a chaining table with col_link's shown is given in Figure 3.13.



Figure 3.13. Chaining Table with col_link's Shown

3.20 end_col_link Field

The end_col_link field of an N-node links to the end of the column of E-nodes lying under this N-node. An example of a chaining table with end_col_link's shown is given in Figure 3.14.



Figure 3.14. Chaining Table with end_col_link's Shown

3.21 pred_link Field

The pred_link field of an N-node links to the preceding N-node found in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with pred_link's shown is given in Figure 3.15.



Figure 3.15. C-S-N Tree with pred_link's Shown

An example of a chaining table with $\texttt{pred_link}\text{'s shown}$ is given in Figure 3.16



Figure 3.16. Chaining Table with pred_link's Shown

3.22 succ_link Field

The succ_link field of an N-node links to the succeeding N-node found in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with succ_link's shown is given in Figure 3.17.



Figure 3.17. C-S-N Tree with succ_link's Shown

An example of a chaining table with succ_link's shown is given in Figure 3.18.



Figure 3.18. Chaining Table with succ_link's Shown

4 Node Processor

The Node Processor module contains functions new_c_node, new_s_node, new_n_node, and new_e_node and has the skeleton shown below in Figure 4.1.

```
#node_proc.py
from globals import *
def new_c_node() -> Node:
def new_s_node() -> Node:
def new_n_node() -> Node:
def new_e_node() -> Node:
```

Figure 4.1. Skeleton of the Node Processor

new_c_node, new_s_node, new_n_node, and new_e_node generate, respectively, a new C-node, S-node, N-node, or E-node, with their fields initialized and are rather straightforward functions. These are shown below in Figures 4.2-4.5.

Function new_c_node returns a new C-node.

def new_c_node() -> Node:
 return new_node(NodeId.C_NODE)

Figure 4.2. Function new_c_node

Function new_s_node returns a new S-node.

```
def new_s_node() -> Node:
    return new_node(NodeId.S_NODE)
```

Figure 4.3. Function new_s_node

Function new_n_node returns a new N-node.

```
def new_n_node() -> Node:
    answer = new_node(NodeId.N_NODE)
    answer.lit = ""
    answer.ftr = [Feature.QUESTION] * N_FEATURES
    answer.end_col_link = None
    answer.pred_link = None
    answer.succ_link = None
    answer.np_link = answer
    return answer
```

Figure 4.4. Function new_n_node

Function new_e_node returns a new E-node.

```
def new_e_node() -> Node:
    answer = new_node(NodeId.E_NODE)
    answer.sub = ' '
    answer.ftr = [Feature.QUESTION] * N_FEATURES
    return answer
```

Figure 4.5. Function new_e_node

4.1 Function view_node_str

There is one output procedure in the Node Processor that has not been discussed above that we need to know about, because we will be looking at some of its output for short while. This is procedure view_node_str which takes as an argument a NodePointer and outputs it in readable form. Otherwise, procedure view_node_str does no processing of its own, and so we do not need to know the details of its inner workings. For us it is enough to be able to understand the output. Function view_node_str has the form indicated in Figure 4.6.

```
def view_node_str(node: Node) -> str:
    """Formatted string representation of node."""
    ...
```

Figure 4.6. Skeleton of Function view_node_str
Some typical output of procedure view_node_str is shown below in Figure 4.7 where a chaining table is listed. (Links from the chaining table to its associated C-S-N tree are also listed by procedure view_node_str.)

| | Nodes |
|----------------|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:5, th:3, nu:2) |
| 3 | (N, lit:June, ftr:[+-++-], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_b , |
| | pr:0, su:4, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-++-], np:3, ch:0, co:3 _b) |
| 3 _b | (E, sub:B, ftr:[+-++-], np:3, ch: 6_a , co: 0) |
| 4 | (N, lit:flowers, ftr:[++?], up:2, dn:0, |
| | lt:3, rt:0, th:5, np:4, ch:0, co: 4_a , ec: 4_b , |
| | pr:3, su:6, nu:4) |
| 4_{a} | (E, sub:A, ftr:[++?], np:4, ch:0, co:4 _b) |
| $4_{\rm b}$ | (E, sub:B, ftr:[++?], np:4, ch:7a, co:0) |
| 5 | (S, up:1, dn:6, lt:2, rt:0, th:6, nu:5) |
| 6 | (N, lit:she, ftr:[++-++-], up:5, dn:0, |
| | lt:0, rt:7, th:7, np:6, ch:0, co: $6_{ m a}$, ec: $6_{ m a}$, |
| | pr:4, su:7, nu:6) |
| 6a | (E, sub:A, ftr:[++-++-], np:6, ch:0, co:0) |
| 7 | (N, lit:them, ftr:[+++??-], up:5, dn:0, |
| | lt:6, rt:0, th:0, np:7, ch:0, co: $7_{ m a}$, ec: $7_{ m a}$, |
| | pr:6, su:0, nu:7) |
| 7_{a} | (E, sub:A, ftr:[+++??-], np:7, ch:0, co:0) |

Figure 4.7. Typical Output from Function view_node_str

(C = C-node, S = S-node, N = N-node, E = E-node, lit = lit field, sub = sub field, ftr = ftr field, up = up_link, dn = down_link, lt = left_link, rt = right_link, th = thread_link, nu = number, np = np_link, ch = chain_link, co = col_link, ec = end_col_link, pr = pred_link, and su = succ_link)

5 Parser

The Parser module defines function parse and has the form shown below in Figure 5.1.

```
#parse_proc.py
from lexicon import *
from node_proc import *
def parse(obj) -> Node:
```

Figure 5.1. Skeleton of the Parser

Function parse accepts as input a **system focus** representation and **system parse tree** that has been generated by a computer natural language system. The output of parse is a C-S-N tree incorporating the information contained in the system parse tree and system focus. The system focus represents the natural language system's focus of attention. This will be gone into in more detail in Chapter 13.

The representation of the system tree inputted to parse is system dependent, and so the details of parse are also system dependent. As the internals of parse are heavily dependent upon and rather involved for any system, we won't go into the details of parse for any particular system here. Hopefully, the reader may glean enough information from the multitude of examples presented in this thesis to get an idea of what parse does. In any case, lack of an actual algorithm for parse isn't so bad since the ideas presented in this thesis are really still in an early stage and it is enough to concentrate on them.

Even though the input to the Parser is not well defined, the output is. The Parser builds from the system parse tree it is given the corresponding C-S-N tree with all up_link's, down_link's, left_link's, right_link's, thread_link's, and number's set to what is expected. Consider example (5.2) below.

(5.2) June hates flowers, but she waters them anyway.



When procedure parse is called on the system parse tree representing (5.2), we get the following output in Figure 5.3.

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| June | - | - | - | + | _ | + | + | - |
| flowers | - | - | - | + | + | ? | - | - |
| she | + | - | - | + | _ | + | + | - |
| them | + | - | - | + | + | ? | ? | - |

| | Nodes |
|---|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:5, th:3, nu:2) |
| 3 | (N, lit:June, ftr:[+-++-], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_b , |
| | pr:0, su:4, nu:3) |
| 4 | (N, lit:flowers, ftr:[++?], up:2, dn:0, |
| | lt:3, rt:0, th:5, np:4, ch:0, co: 4_a , ec: 4_b , |
| | pr:3, su:6, nu:4) |
| 5 | (S, up:1, dn:6, lt:2, rt:0, th:6, nu:5) |
| 6 | (N, lit:she, ftr:[++-++-], up:5, dn:0, |
| | lt:0, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_a , |
| | pr:4, su:7, nu:6) |
| 7 | (N, lit:them, ftr:[+++??-], up:5, dn:0, |
| | lt:6, rt:0, th:0, np:7, ch:0, co:7 _a , ec:7 _a , |
| | pr:6, su:0, nu:7) |

Figure 5.3. Typical Output from parse

Listing of nodes in Figure 5.3 is done by procedure view_node_str of the Node Processor described in Chapter 4. The C-S-N parse tree is slightly more complicated when focusing is taken into account, but for the time being we will ignore its effects. We will discuss the effects of focusing on C-S-N parse trees in Chapter 13.

When Figure 5.3 is drawn as a tree, we get a structure like Figure 5.4.



Figure 5.4. Output from parse Drawn as a Tree

6 Primary Utilities

The Primary Utilities module defines the boolean functions precede, command, and separate corresponding to the *precedes*, *commands*, and *is separate from* relations discussed in Chapter 1. The skeleton of the primary utilities module is shown below in Figure 6.1.

```
#primary_uty.py
from globals import *
def precede(n1: Node, n2: Node) -> bool:
def command(n1: Node, n2: Node) -> bool:
def separate(n1: Node, n2: Node) -> bool:
```

Figure 6.1. Skeleton of the Primary Utilities

The precede, command, and separate functions do just what is expected. They are true if and only if the *precedes*, *commands*, and *is separate from* relations hold between their arguments. Along with function dominate which is used by separate, these functions are shown below in Figures 6.2-6.5.

Function precede is true if and only if n1 precedes n2.

```
def precede(n1: Node, n2: Node) -> bool:
    return n1.number < n2.number</pre>
```

Figure 6.2. Function precede

Function dominate is true if and only if n1 dominates n2.

```
def dominate(n1: Node, n2: Node) -> bool:
    if n1.number == n2.number:
        return True
    child = n1.down_link
    while child is not None:
        if dominate(child, n2):
            return True
        child = child.right_link
    return False
```

Figure 6.3. Function dominate

Function command is true if and only if n1 commands n2.

```
def command(n1: Node, n2: Node) -> bool:
    return dominate(n1.up_link, n2)
```

Figure 6.4. Function command

Function separate is true if and only if n1 is separate from n2.

```
def separate(n1: Node, n2: Node) -> bool:
    parent = n1.up_link
    while not dominate(parent, n2):
        parent = parent.up_link
    return parent.id == NodeId.C_NODE
```

Figure 6.5. Function separate

7 Secondary Utilities

The Secondary Utilities module defines functions sc, agr, and rnr. These stand for Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule. The skeleton of the Secondary Utilities module is shown below in Figure 7.1.

```
#secondary_uty.py
from primary_uty import *
def sc(n1: Node, n2: Node) -> bool:
def agr(n1: Node, n2: Node) -> bool:
def rnr(n1: Node, n2: Node) -> bool:
```

Figure 7.1. Skeleton of Secondary Utilities

7.1 Syntactic Conditions

As shown in Chapter 1, certain constraints such as the Precedes and Commands Rule apply in forward pronominalization. Function sc is true whenever these grosser syntactic constraints are met. In this thesis, we let sc be true when the Precedes and Commands Rule is satisfied, function sc is shown below in Figure 7.2.

```
def sc(n1: Node, n2: Node) -> bool:
    return not (precede(n1, n2) and (command(n1, n2) or separate(n1, n2)))
```

Figure 7.2. Function sc (Syntactic Conditions)

7.2 Agreement

Besides satisfying Syntactic Conditions, there has to be agreement between a node and its chaining node. First person, second person, third person, plural, gender, and animate features have to agree in order for one node to chain to another. Function agr is shown below in Figure 7.3.

```
def agr(n1: Node, n2: Node) -> bool:
    ftr1 = n1.ftr
    ftr2 = n2.ftr
    return (eq_feat(ftr1[FeatureIndex.FPF], ftr2[FeatureIndex.FPF]) and
        eq_feat(ftr1[FeatureIndex.SPF], ftr2[FeatureIndex.SPF]) and
        eq_feat(ftr1[FeatureIndex.TPF], ftr2[FeatureIndex.TPF]) and
        eq_feat(ftr1[FeatureIndex.PLF], ftr2[FeatureIndex.PLF]) and
        eq_feat(ftr1[FeatureIndex.GNF], ftr2[FeatureIndex.GNF]) and
        eq_feat(ftr1[FeatureIndex.ANF], ftr2[FeatureIndex.ANF]))
```

Figure 7.3. Function agr (Agreement)

7.3 Equal Features

Function agr uses function eq_feat. eq_feat tests if two Feature's are equal. As indicated in Chapter 1, Feature's are equal unless a PLUS and MINUS are compared. Function eq_feat is shown below in Figure 7.4.

```
def eq_feat(f1: Feature, f2: Feature) -> bool:
    if f1 == Feature.PLUS:
        return f2 != Feature.MINUS
    elif f1 == Feature.MINUS:
        return f2 != Feature.PLUS
    else: # f1 == Feature.QUESTION
        return True
```

Figure 7.4. Function eq_feat (Equal Features)

7.4 Reflexive Nonreflexive Rule

The distinction between reflexive pronouns and nonreflexive pronouns is that reflexive pronouns cannot chain to an N-node that is outside of the same simplex in which it occurs, while a nonreflexive pronoun can. This rule will have to be modified later for genitives, but for now we can suppose that a nonreflexive pronoun must chain to an N-node outside of the same simplex in which it is in. Shown in Figure 7.5 is function rnr which is true when the reflexive nonreflexive rule is satisfied.

```
def rnr(n1: Node, n2: Node) -> bool:
    ftr1 = n1.np_link.ftr
    ftr2 = n2.np_link.ftr
    if ftr2[FeatureIndex.GEN] == Feature.PLUS:
        return False
    elif ftr1[FeatureIndex.RPF] == Feature.PLUS:
        return (n1.up_link == n2.up_link)
            and (ftr1[FeatureIndex.GEN] == Feature.MINUS)
    elif ftr1[FeatureIndex.RPF] == Feature.MINUS:
        return (n1.up_link != n2.up_link)
            or (ftr1[FeatureIndex.GEN] != Feature.MINUS)
```

Figure 7.5. Function rnr (Reflexive Nonreflexive Rule)

8 Table Processor I

The Table Processor module defines function chaining which takes as input a C-S-N tree and returns its chaining table. The actions of function chaining in the Table Processor can only be understood by example, and this is what this chapter provides. In Chapter 9, we'll look at the actual algorithms and, in Chapter 10, we'll look at some actual output.

So, let us consider sentence (8.1) below.

(8.1) John wants to give June a present, but he isn't sure she'll like it.



The Parser builds from the system parse tree of (8.1) the corresponding C-S-N tree with six N-nodes which have the lit fields and ftr's indicated below in Figure 8.2.

| | Features | | | | | | | |
|---------|---------------------------------|---|---|---|---|---|---|---|
| | PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | |
| John | - | _ | - | + | _ | _ | + | - |
| ϕ | + | ? | ? | ? | ? | ? | ? | - |
| June | - | _ | - | + | _ | + | + | - |
| present | - | _ | - | + | _ | ? | - | - |
| he | + | - | - | + | _ | _ | + | - |
| she | + | - | - | + | - | + | + | - |
| it | + | - | - | + | _ | ? | - | - |

Figure 8.2. lit Fields and ftr's of the N-Nodes

The C-S-N tree itself has the form of Figure 8.3 below.



Figure 8.3. C-S-N Parse Tree

After parse is called, chaining is called. The first thing to happen is the initialization of the chaining table for the C-S-N tree. Below each N-node is suspended, by the col_link of the N-node, a new E-node with subscript A. Each new E-node has a np_link back to the N-node it is suspended from. As well, the Feature's of each new E-node are copied from the N-node it is suspended from. Attached to the first and last N-nodes is an S-node to make it easy to keep track of the first and last N-nodes in the chaining table. The chaining table, as it looks immediately after initialization, is shown below in Figure 8.4.



Figure 8.4. Chaining Table Immediately after Initialization

The chaining algorithm works by walking backwards across N-nodes in the top row and walking down columns of E-nodes. The chaining algorithm works on two N-nodes at a time. If the first is compatible with the second under Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule, then the E-nodes underneath the first N-node that agree with the second N-node are chain_link'ed to copies of the second N-node.

The last N-node in the table is \underline{it} , the chaining table begins with \underline{it} . \underline{it} can't chain to itself, so the second N-node in the description above becomes <u>she</u> and the chaining algorithm compares \underline{it} to <u>she</u>. Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(it, she) = True
agr(it, she) = False
```

The chaining algorithm now moves from <u>she</u> to <u>he</u> and compares <u>it</u> to <u>he</u>. Again Syntactic Conditions are satisfied, but Agreement isn't.

sc(it, he) = True
agr(it, he) = False

The chaining algorithm moves from <u>he</u> to <u>present</u> and compares <u>it</u> to <u>present</u>. This time, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule are satisfied.

```
sc(it, present) = True
agr(it, present) = True
rnr(it, present) = True
```

Since all three rules are satisfied, a chain from $\underline{it_a}$ to a copy of present may be created. This happens if $\underline{it_a}$ and present agree, and they do.

agr(ita, present) = True

The chaining algorithm makes a new E-node copy of present, present_b, and hangs it below present. The chain_link of present_b is set to ita and the semantic features of ita, but not the syntactic features, are copied into the semantic features of present_b. After chaining present_b to ita, the chaining table appears as shown in Figure 8.5.



Figure 8.5. Chaining Table after Chaining $present_b$ to <u>ita</u>

The chaining algorithm now moves from present to June. Syntactic Conditions are satisfied, but Agreement isn't.

sc(it, June) = True
agr(it, June) = False

The chaining algorithm now moves from <u>June</u> to ϕ . This time, all three rules, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule are satisfied.

sc(it, ϕ) = True agr(it, ϕ) = True rnr(it, ϕ) = True

Since all three rules are satisfied, E-nodes under <u>it</u> that agree with $\underline{\phi}$ chain to copies of $\underline{\phi}$. <u>it</u>_a is compared to $\underline{\phi}$, and it is seen that they agree.

agr(ita, ϕ) = True

The chaining algorithm makes a new E-node copy of $\underline{\phi}$, $\underline{\phi}_{\rm b}$, and hangs it below $\underline{\phi}$. The chain_link of $\underline{\phi}_{\rm b}$ is set to ita and the semantic features of ita are copied into the semantic features of $\underline{\phi}_{\rm b}$. After chaining $\underline{\phi}_{\rm b}$ to ita, the chaining table appears as shown in Figure 8.6.



Figure 8.6. Chaining Table After Chaining $\phi_{\mathbf{b}}$ to $\underline{\mathbf{it}}_{\mathbf{a}}$

The chaining algorithm now moves to <u>John</u> and compares <u>John</u> to <u>it</u>. Syntactic Conditions hold, but Agreement does not.

sc(it, John) = True
agr(it, John) = False

Having exhausted all possible combinations with \underline{it} , the chaining algorithm considers she.

The chaining algorithm tries comparing <u>she</u> to <u>it</u>, but Syntactic Conditions are not satisfied.

sc(she, it) = False

The chaining algorithm moves from \underline{it} to <u>she</u>, but <u>she</u> can't chain to <u>she</u>, so the chaining algorithm moves to <u>he</u>. This time Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(she, he) = True
agr(she, he) = False
```

The chaining algorithm now moves from \underline{he} to $\underline{present}$ where again Syntactic Conditions are satisfied, but Agreement isn't.

sc(she, present) = True
agr(she, present) = False

The chaining algorithm moves from $\underline{\text{present}}$ to $\underline{\text{June}}$. This time all three rules are satisfied.

sc(she, June) = True
agr(she, June) = True
rnr(she, June) = True

As all three rules are satisfied, E-nodes under <u>she</u> that agree with <u>June</u> chain to copies of <u>June</u>. she_a is compared to <u>June</u>, and it is seen that they agree.

```
agr(shea, June) = True
```

The chaining algorithm makes a new E-node copy of <u>June</u>, <u>June</u>, <u>June</u>, and hangs it below <u>June</u>. The chain_link of <u>June</u> is set to <u>shea</u> and the semantic features of <u>shea</u> are copied into the semantic features of <u>June</u>. After chaining June_b to <u>shea</u>, the chaining table appears as shown in Figure 8.7.



Figure 8.7. Chaining Table after Chaining June_b to shea

The chaining algorithm now moves from <u>June</u> to ϕ and compares <u>she</u> to ϕ . All three rules are satisfied.

sc(she, ϕ) = True agr(she, ϕ) = True rnr(she, ϕ) = True

Copies of ϕ are chain_link'ed to E-nodes under she that agree with ϕ . shea is compared to ϕ , and it is seen that they agree.

agr(shea, ϕ) = True

A new E-node copy of ϕ , $\phi_{\underline{C}}$, is made and hung below ϕ . The chain_link of $\phi_{\underline{C}}$ is set to shea and the semantic features of shea are copied into the semantic features of $\phi_{\underline{C}}$. After chaining $\phi_{\underline{C}}$ to shea, the chaining table appears as shown in Figure 8.8.



Figure 8.8. Chaining Table after Chaining $\phi_{\mathbf{C}}$ to shea

The chaining algorithm now moves from ϕ to John and compares she to John. It is seen that Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(she, John) = True
agr(she, John) = False
```

This completes the creation of chain_link's to E-nodes under <u>she</u>. The chaining algorithm now considers <u>he</u>.

he is compared to it, but it is seen that Syntactic Conditions aren't satisfied.

sc(he, it) = False

The chaining algorithm moves from \underline{it} to <u>she</u>, but again, Syntactic Conditions aren't satisfied.

sc(he, she) = False

The chaining algorithm moves from <u>she</u> to <u>he</u>, but <u>he</u> can't chain to <u>he</u>, so the chaining algorithm moves from <u>he</u> to <u>present</u>. This time Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(he, present) = True
agr(he, present) = False
```

The chaining algorithm moves from $\underline{\text{present}}$ to $\underline{\text{June}}$, and similar results happen.

```
sc(he, June) = True
agr(he, June) = False
```

Next, the chaining algorithm moves from <u>June</u> to ϕ , and, this time, all three rules, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule, are satisfied.

sc(he, ϕ) = True agr(he, ϕ) = True rnr(he, ϕ) = True

Copies of <u>he</u> are chain_link'ed to E-nodes under ϕ that agree with <u>he</u>. <u>hea</u> is compared to ϕ , and it is seen that they agree.

agr(hea,
$$\phi$$
) = True

A new E-node copy of $\underline{\phi}$, $\underline{\phi_d}$, is made and hung below $\underline{\phi}$. The chain_link of $\underline{\phi_d}$ is set to <u>hea</u> and the semantic features of <u>hea</u> are copied into the semantic features of $\underline{\phi_d}$. After chaining $\underline{\phi_d}$ to <u>hea</u>, the chaining table appears as shown in Figure 8.9.



Figure 8.9. Chaining Table after Chaining ϕ_d to <u>hea</u>

The chaining algorithm now moves from ϕ to <u>John</u> and <u>he</u> is compared to <u>John</u>. It is seen that all three rules are satisfied.

sc(he, John) = True
agr(he, John) = True
rnr(he, John) = True

So, copies of <u>he</u> are chain_link'ed to E-nodes under <u>John</u> that agree with <u>he</u>. hea is compared to <u>John</u>, and it is seen that they agree. $agr(he_a, John) = True$

A new E-node copy of <u>John</u>, <u>John</u>, is made and hung below <u>John</u>. The chain_link of John_b is set to <u>hea</u> and the semantic features of <u>hea</u> are copied into the semantic features of John_b. After chaining <u>John</u>_b to <u>hea</u>, the chaining table appears as shown in Figure 8.10.



Figure 8.10. Chaining Table after Chaining John_b to $\underline{he_a}$

Having completed the processing of <u>he</u>, the chaining algorithm considers <u>present</u> is not a pronoun though, so the chaining algorithm moves on to <u>June</u>. Similarly, <u>June</u> is not a pronoun, so the chaining algorithm now considers ϕ .

The chaining algorithm compares $\underline{\phi}$ to \underline{it} , and it is seen that Syntactic Conditions don't hold.

 $sc(\phi, it) = False$

The chaining algorithm moves from \underline{it} to \underline{she} , \underline{she} to \underline{he} , \underline{he} to $\underline{present}$, and present to June with little more success.

sc(ϕ , she) = False sc(ϕ , he) = False sc(ϕ , present) = False sc(ϕ , June) = False

The chaining algorithm moves from <u>June</u> to ϕ , but ϕ can't chain to ϕ . So now, the chaining algorithm moves from ϕ to <u>John</u>. This time, all three rules are satisfied.

```
sc(\phi, John) = True

agr(\phi, John) = True

rnr(\phi, John) = True
```

Copies of <u>John</u> are chain_link'ed to E-nodes under ϕ that agree with <u>John</u>. ϕ_a is compared to <u>John</u>, and it is seen that they agree.

 $agr(\phi_a, John) = True$

Thus, new copy of <u>John</u>, <u>John</u> is made. <u>John</u> is chain_link'ed to ϕ_{a} . The semantic features of ϕ_{a} are copied into <u>John</u>. After chaining <u>John</u> to ϕ_{a} , the chaining table appears as shown in Figure 8.11.



Figure 8.11. Chaining Table after Chaining John_c to ϕ_{a}

 $\phi_{\rm b}$ is compared to <u>John</u>, and it is seen that they don't agree.

 $agr(\phi_{b}, John) = False$

 $\underline{\phi_{b}}$ and <u>John</u> don't agree because when $\underline{\phi_{b}}$ was chain_link'ed to <u>ita</u>, the semantic features of <u>ita</u> were copied into the semantic features of $\underline{\phi_{b}}$. Hence, the information that <u>ita</u> was inanimate was copied into $\underline{\phi_{b}}$ preventing a ridiculous chain: <u>JohnX</u> is chained to $\underline{\phi_{b}}$ is chained to <u>ita</u>. $\underline{\phi_{c}}$, which was chained to <u>shea</u>, is compared to <u>John</u>, and it is seen that they don't agree.

 $agr(\phi_{C}, John) = False$

On the other hand, ϕ_{d} , which was chained to <u>hea</u>, does agree with <u>John</u>.

 $agr(\phi_d, John) = True$

Thus, new copy of <u>John</u>, <u>John</u> is made. <u>John</u> is chain_link'ed to $\phi_{\underline{d}}$. The semantic features of $\phi_{\underline{d}}$ are copied into the semantic features of John_d. After chaining John_d to $\phi_{\overline{d}}$, the chaining table appears as shown in Figure 8.12.



Figure 8.12. Chaining Table after Chaining John_d to ϕ_d

Having completed chaining to $\underline{\phi}$, the chaining algorithm moves to John. John is not a pronoun, so the chaining algorithm now stops as it has reached the end of the chaining table. This makes Figure 8.12, above, the finished chaining table.

9 Table Processor II

From Chapter 8 we know that the Table Processor module defines function chaining which takes as input a C-S-N tree and which returns as output the chaining table of the inputted C-S-N tree. In Chapter 8, we illustrated the kind of processing the Table Processor does by working through in detail a typical example. In this chapter, we will go into the particulars of the Table Processor algorithms. In Chapter 10, we'll look at some actual output.

The skeleton of the Table Processor module is shown below in Figure 9.1. The Table Processor module defines function chaining.

```
#table_proc.py
from node_proc import *
from parser import *
from secondary_uty import *
def chaining(nnodes: list[Node]) -> None:
```

Figure 9.1. Skeleton of the Table Processor

Function chaining is the algorithm we described by example in Chapter 8. chaining takes as input a C-S-N tree and returns the chaining table of the inputted C-S-N tree. Function chaining is shown below in Figure 9.2.

```
def chaining(nnodes: list[Node]) -> None:
    init_table(nnodes)
    for nl in reversed(nnodes):
        # For each N-node nl that is a pronoun, call
        # procedure chaining_n.
        if nl.ftr[FeatureIndex.PNF] == Feature.PLUS:
            chaining_n(nnodes, nl)
```

Figure 9.2. Function chaining

The first thing function chaining does is to call init_table which initializes the chaining table as described in the previous chapter. Function init_table is shown in Figure 9.3.

```
def init_table(nnodes: list[Node]) -> None:
    last = None
    for n in nnodes:
        n.col_link = new_e_node()
        n.col_link.ftr = n.ftr.copy()
        n.col_link.np_link = n
        n.col_link.sub = 'A'
        n.end_col_link = n.col_link
        n.pred_link = last
        if last is not None:
            last.succ_link = n
        last = n
```

Figure 9.3. Function init_table

Below each N-node is hung a new E-node with Feature's copied from the N-node and np_link back to the N-node. The col_link's and end_col_link's of the N-nodes are updated accordingly. pred_link's and succ_link's are set in init_table using the thread_link's which were established by the Parser. Finally, at the end of the procedure, table, a variable global inside the Table Processor, has its left_link and right_link set to the first and last N-node.

For each N-node that is a pronoun, function chaining calls procedure chaining_n. chaining_n calls refl_chaining or non_refl_chaining depending on whether or not the inputted N-node is reflexive or not. Function chaining_n is shown below in Figure 9.4.

```
def chaining_n(nnodes: list[Node], n1: Node) -> None:
    if n1.ftr[FeatureIndex.RPF] == Feature.PLUS:
        # Inputted pronoun N-node n1 is reflexive.
        refl_chaining(n1)
    elif n1.ftr[FeatureIndex.RPF] == Feature.MINUS:
        # Inputted pronoun N-node n1 isn't reflexive.
        non refl chaining(nnodes, n1)
```

Figure 9.4. Function chaining_n

Function non_refl_chaining handles nonreflexive pronouns. Function non_refl_chaining is shown below in Figure 9.5.

```
def non_refl_chaining(nnodes: list[Node], n1: Node) -> None:
    for n2 in reversed(nnodes):
        if n2 != n1:
            chaining_n_to_n(n1, n2)
```

Figure 9.5. Function non_refl_chaining

non_refl_chaining calls chaining_n_to_n on the inputted N-node with each N-node in the chaining table except itself. This takes care of creating all chains to E-nodes lying under the inputted N-node.

Function refl_chaining is very similar to non_refl_chaining and is shown below in Figure 9.6.

```
def refl_chaining(n1: Node) -> None:
    n2 = simplex_pred(n1)
    while n2 is not None:
        if n2 != n1:
            chaining_n_to_n(n1, n2)
        n2 = simplex_pred(n2)
```

Figure 9.6. Function refl_chaining

Since the N-node inputted to refl_chaining is reflexive, refl_chaining only calls chaining_n_to_n on the inputted N-node with each preceding N-node within the same simplex as the inputted N-node.

Function simplex_pred, which is used by procedure refl_chaining, simply returns the N-node that preceeds the inputted N-node in the same simplex. Function simplex_pred is shown below in Figure 9.7.

```
def simplex_pred(n1: Node) -> Node:
    answer = n1
    while True:
        answer = answer.left_link
        if answer is None or answer.id == NodeId.N_NODE:
            return answer
```

Figure 9.7. Function simplex_pred

Function chaining_n_to_n is called by procedures refl_chaining and non_refl_chaining and is shown below in Figure 9.8.

```
def chaining_n_to_n(n1: Node, n2: Node) -> None:
    if not sc(n1, n2) or not agr(n1, n2) or not rnr(n1, n2):
        return
    old_end_col_link = n1.end_col_link
    e1 = n1
    while e1 != old_end_col_link:
        e1 = e1.col_link
        if e1 is not None:
            chaining_e_to_n(e1, n2)
```

Figure 9.8. Function chaining_n_to_n

If Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule hold, then procedure chaining_e_to_n is called on each E-node lying underneath the first N-node.

Function chaining_e_to_n, which is called by procedure chaining_n_to_n, is shown below in Figure 9.9.

```
def chaining_e_to_n(e1: Node, n2: Node) -> None:
    if agr(e1, n2):
        new_chain(e1, n2)
```

Figure 9.9. Function chaining_e_to_n

If the inputted E-node agrees with the inputted N-node, then a new chain is created from a copy of the inputted N-node to the inputted E-node by calling procedure new_chain.

Function new_chain, which is called by chaining_e_to_n, is shown below in Figure 9.10.

```
def new_chain(e1: Node, n2: Node) -> None:
    n = new_e_node()
    n.np_link = n2
    n.chain_link = e1
    n.sub = chr(ord(n2.end_col_link.sub) + 1)
    # Replace n2 nonsyntactic QUESTION (?) features
    for i in range(N_FEATURES):
        if n2.ftr[i] == Feature.QUESTION and i != FeatureIndex.RPF:
            n.ftr[i] = e1.ftr[i]
        else:
            n.ftr[i] = n2.ftr[i]
        n2.end_col_link.col_link = n
        n2.end_col_link = n
```

Figure 9.10. Function new_chain

Function new_chain creates a copy of the inputted N-node and chain_link's it to the inputted E-node. Semantic Feature's are copied from the inputted E-node to the copy of the inputted N-node.

10 Table Processor III

The last two chapters have been devoted to describing the Table Processor. It is time for some real examples: The first example we present in this chapter was produced using procedure view_node_str of the Node Processor along with some intermittent write statements indicating when LIP are entering and exiting some of the more important routines and some or their results. We start off with the example first presented in Chapter 8.

(10.1) John wants to give June a present, but he isn't sure she'll like it.



Processing (10.1) with some intermediate output gives the listing shown below. This is somewhat verbose, but later examples will be cleaner and shorter, though less detailed output.

| | Features | | | | | | | |
|---------|---------------------------------|---|---|---|---|---|---|---|
| | PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | |
| John | - | - | - | + | - | - | + | - |
| ϕ | + | ? | ? | ? | ? | ? | ? | - |
| June | - | _ | - | + | - | + | + | - |
| present | - | _ | - | + | - | ? | - | - |
| he | + | - | - | + | - | - | + | - |
| she | + | - | - | + | - | + | + | - |
| it | + | - | - | + | - | ? | - | - |

| | Nodos |
|----------|---|
| | INDUES |
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:0, ec:0, |
| | pr:0, su:0, nu:3) |
| 4 | $(S un \cdot 2 dn \cdot 5]t \cdot 3 rt \cdot 0 th \cdot 5 nu \cdot 4)$ |
| 5 | $(N = 1)^{++,+} = (+222222)^{-1} = (N = 1)^{++,+} = (+222222)^{-1} = (N = 1)^{++,+} = (+222222)^{-1} = (N = 1)^{++,+} = (+22222)^{-1} = (N = 1)^{++} = (+22222)^{++} = (+22222)^{++} = (+22222)^{++} = (+222222)^{++} = (+22222)^{++} = (+222222)^{++} = (+22222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+2222222)^{++} = (+22222222)^{++} = (+22222222)^{++} = (+222222222)^{++} = (+222222222)^{++} = (+22222222)^{++} = (+222222222)^{++} = (+222222222)^{++} = (+2222222222)^{++} = (+222222222)^{++} = (+222222222)^{++} = (+2222222222)^{++} = (+2222222222)^{++} = (+22222222222)^{++} = (+222222222222)^{++} = (+222222222222222)^{++} = (+22222222222222222222222222222222222$ |
| 9 | (N, $\Pi \cup \varphi$, $\Pi \cup (\varphi)$ |
| | It:0, rt:6, tn:6, np:5, cn:0, co:0, ec:0, |
| | pr:0, su:0, nu:5) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co:0, ec:0, |
| | pr:0, su:0, nu:6) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | $1t \cdot 6$, $rt \cdot 0$, $th \cdot 8$, $np \cdot 7$, $ch \cdot 0$, $co \cdot 0$, $ec \cdot 0$. |
| | r:0 su:0 $ru:7$ |
| 0 | $(C_{1}, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,$ |
| 0 | (S, up:1, an:9, 11:2, rt:0, th:9, hu:8) |
| 9 | (N, lit:he, ftr:[++-], up:8, dn:0, |
| | lt:0, rt:10, th:10, np:9, ch:0, co:0, ec:0, |
| | pr:0, su:0, nu:9) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, co:0, ec:0, |
| | pr:0 su:0 pu:11) |
| 19 | $(N = 1i + i + f + r \cdot [+ - + -2] = u - 10 = d - 0$ |
| 14 | (w, 110, 10, 101, [1-1-1-2], up.10, all.0, 10, 11, 11, 11, 11, 11, 11, 11, 11, 1 |
| | IL:II, rL:U, TN:U, np:IZ, CN:U, CO:U, eC:U, |
| | pr:0, su:0, nu:12) |

chaining init_table

| | Nodes |
|----------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:3a, ec:3a, |
| | pr:0, su:5, nu:3) |
| 3 _a | (E, sub:A, ftr:[+-], np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m a}$, |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: $6_{ m a}$, ec: $6_{ m a}$, |
| _ | pr:5, su:7, nu:6) |
| 6 _a | (E, sub:A, ftr:[+-++-], np:6, ch:0, co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | lt:6, rt:0, th:8, np:7, ch:0, co: 7_a , ec: 7_a , |
| _ | pr:6, su:9, nu:7) |
| 7a | (E, sub:A, ftr:[+-?], np:7, ch:0, co:0) |
| 8 | (S, up:1, dn:9, 1t:2, rt:0, th:9, nu:8) |
| 9 | (N, lit:he, ftr:[++-], up:8, dn:0, |
| | It:0, rt:10, th:10, np:9, ch:0, co: 9_a , ec: 9_a , |
| 0 | pr:/, su:11, nu:9) |
| 9a | (E, Sub:A, $ICI:[+-++-+-]$, $Ip:9$, $CI:0$, $CO:0$) |
| 10 | (S, up:0, an:11, 10.9, 10.0, ch:11, ha:10) |
| 11 | (N, 110.5Ne, 101.[1-1-1+1], up.10, un.0, 11.0, 11.0] |
| | pr.9 eu.12 pu.11) |
| 11. | $(F \text{sub-A} \text{ftr} \cdot [+-+++-] \text{nn} \cdot 11 \text{cb} \cdot 0 \text{co} \cdot 0)$ |
| 12 | (N littit $ftr \cdot [+-++2] = 10 \cdot 10 dr \cdot 0$ |
| 14 | 1 ± 11 , $r \pm 0$, $t \pm 0$, $n \pm 12$, $c \pm 0$, $c \pm 12$, $c \pm 12$ |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
init_table: exiting
chaining_n(it)
    non_refl_chaining(it)
        chaining_n_to_n(it, she)
            sc(it, she) = True
            agr(it, she) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, he)
            sc(it, he) = True
            agr(it, he) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, present)
            sc(it, present) = True
            agr(it, present) = True
            rnr(it, present) = True
            chaining_e_to_n(ita, present)
                agr(ita, present) = True
                new_chain(ita, present)
                    new_chain: create presentb
                    new_chain: create present_b^{it_a}
```

| | Nodes |
|----------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_a , |
| | pr:0, su:5, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-], np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m a}$, |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: $6_{ m a}$, ec: $6_{ m a}$, |
| _ | pr:5, su:7, nu:6) |
| 6 _a | (E, sub:A, ftr:[+-++-], np:6, ch:0, co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | lt:6, rt:0, th:8, np:7, ch:0, co:7 _a , ec:7 _b , |
| _ | pr:6, su:9, nu:7) |
| $\frac{7a}{-}$ | (E, sub:A, ftr:[+-?], np:7, ch:0, co:7 _b) |
| $7\mathbf{b}$ | (E, sub:B, ftr:[+-?], np:7, ch:12a, co:0) |
| 8 | (S, up:1, dn:9, It:2, rt:0, th:9, nu:8) |
| 9 | (N, lit:ne, itr:[++-], up:8, dn:0, |
| | It:0, rt:10, tn:10, np:9, cn:0, co: 9_a , ec: 9_a , |
| 0 | $(F \text{ subs}) (F \text$ |
| 9a | (E, Sub: A, ICI: [T-T-T-], ID: 9, CII: 0, CO: 0) |
| 11 | (S, up.0, un.11, 10.9, 10.0, 0, 0) |
| 11 | 1+0 rt 12 th 12 np 11 ch 0 co 11 ec 11 |
| | pr.9 su.12 $pu.11$) |
| 11. | $(E \text{sub-A} \text{ftr} \cdot [+-+++-] \text{nn} \cdot 11 \text{ch} \cdot 0 \text{co} \cdot 0)$ |
| 12^{11a} | (N, lit:it, ftr:[+-++?], up:10, dn:0. |
| | 1t:11, rt:0, th:0, np:12, ch:0, co:12a, ec:12a. |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(it, June)

sc(it, June) = True

agr(it, June) = False

chaining_n_to_n: exiting

chaining_n_to_n(it, \phi)

sc(it, \phi) = True

agr(it, \phi) = True

rnr(it, \phi) = True

chaining_e_to_n(it_a, \phi)

agr(it_a, \phi) = True

new_chain: create \phi_b

new_chain: create \phi_b^{-}ita
```

| | Nodes |
|-----------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:3a, ec:3a, |
| | pr:0, su:5, nu:3) |
| 3 _a | (E, sub:A, ftr:[+-], np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m b}$, |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| ^{5}b | (E, sub:B, ftr:[++-?], np:5, ch:12a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_a , |
| | pr:5, su:7, nu:6) |
| 6 _a | (E, sub:A, ftr:[+++-], np:6, ch:0, co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | It:6, rt:0, th:8, np:/, ch:0, co: 7_{a} , ec: 7_{b} , |
| - | pr:6, su:9, nu:/) |
| /a | (E, sub:A, itr:[+-?], np:/, cn:0, co: (b)) |
| (b | (E, SUD:B, Itr: $[+-?]$, np: /, Cn: I_{2a} , CO: U) |
| 0 | (S, up:1, an:9, 11:2, r1:0, 1n:9, nu:8) |
| 9 | (N, 110.110, 101.1(1-1-1-1), up.o, un.0, 110.0) |
| | pr.7 su.11 $pu.9$ |
| 9. | $(E \text{sub-A} \text{ftr} \cdot [++-] \text{nn} \cdot 9 \text{ch} \cdot 0 \text{co} \cdot 0)$ |
| 10 | $(S_{1}, S_{2}, S_{2}, S_{1}, S_{2}, S_{1}, S_{2}, S_{2},$ |
| 11 | (N. lit:she. $ftr:[+-+++-]$, up:10, dn:0. |
| | 1t:0, rt:12, th:12, np:11, ch:0, co:11a, ec:11a, |
| | pr:9, su:12, nu:11) |
| 11 _a | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co:12a, ec:12a, |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(it, John)
            sc(it, John) = True
            agr(it, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(she)
    non_refl_chaining(she)
        chaining_n_to_n(she, it)
            sc(she, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, he)
            sc(she, he) = True
            agr(she, he) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, present)
            sc(she, present) = True
            agr(she, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, June)
            sc(she, June) = True
            agr(she, June) = True
            rnr(she, June) = True
            chaining_e_to_n(shea, June)
                agr(she_a, June) = True
                new_chain(shea, June)
                    new_chain: create Juneb
                    new_chain: create June<sub>b</sub>^shea
```

| | Nodes |
|----------------|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_a , |
| | pr:0, su:5, nu:3) |
| 3 _a | (E, sub:A, ftr:[+], np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: 5_a , ec: 5_b , |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| 5 _b | (E, sub:B, ftr:[++-?], np:5, ch:12a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| | pr:5, su:7, nu:6) |
| 6 _a | (E, sub:A, ftr:[+-++-], np:6, ch:0, co:6 _b) |
| 6 _b | (E, sub:B, ftr:[+-++-], np:6, ch:11 _a , co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | lt:6, rt:0, th:8, np:7, ch:0, co:7 _a , ec:7 _b , |
| | pr:6, su:9, nu:7) |
| 7_{a} | (E, sub:A, ftr:[+-?], np:7, ch:0, co:7 _b) |
| 7 _b | (E, sub:B, ftr:[+-?], np:7, ch:12a, co:0) |
| 8 | (S, up:1, dn:9, lt:2, rt:0, th:9, nu:8) |
| 9 | (N, lit:he, ftr:[++], up:8, dn:0, |
| | lt:0, rt:10, th:10, np:9, ch:0, co: 9_a , ec: 9_a , |
| | pr:7, su:11, nu:9) |
| 9 _a | (E, sub:A, ftr:[++-], np:9, ch:0, co:0) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, $co:1l_a$, $ec:1l_a$, |
| | pr:9, su:12, nu:11) |
| 11_{a} | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | It:11, rt:0, th:0, np:12, ch:0, co:1 2_a , ec:1 2_a , |
| 10 | pr:11, su:0, nu:12) |
| 12a | (E, sub:A, ftr: ++-? , np:12, ch:0, co:0) |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(she, \phi)

sc(she, \phi) = True

agr(she, \phi) = True

rnr(she, \phi) = True

chaining_e_to_n(shea, \phi)

agr(shea, \phi) = True

new_chain(shea, \phi)

new_chain: create \phi_c

new_chain: create \phi_c^shea
```

| | Nodes |
|-----------------|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_{a} , ec: 3_{a} , |
| | pr:0, su:5, nu:3) |
| 39 | (E. sub:A. ftr: $[++-]$, np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | $(N,]it:\phi, ftr:[+?????-], up:4, dn:0.$ |
| | $1t \cdot 0, rt \cdot 6, th \cdot 6, np \cdot 5, ch \cdot 0, co \cdot 5_{22}, ec \cdot 5_{23}$ |
| | pr:3 su:6 nu:5) |
| 5. | (E sub A ftr [+22222] no 5 cb 0 co 5) |
| 51 | (E sub B $ftr [+-+-2]$ np:5 $ch \cdot 12$ $co \cdot 5c$) |
| 50 | $(E \text{ sub} \cdot C \text{ ftr} [+-++++-] \text{ np} \cdot 5 \text{ ch} \cdot 11_a, \text{ co} \cdot 0)$ |
| 6 | (N) lit June $ftr [+++-]$ up $dn \cdot 0$ |
| U | 1 + 5 + 7 + 107 |
| | pr.5, su.7, pu.6) |
| 6. | $(F \text{sub-} \Lambda \text{ftr} \cdot [+++-] \text{nn-} h \text{ch-} h ch-$ |
| 61 | (E sub:B $ftr:[+++-]$ np:6 $cb:11_{a}$ $co:0)$ |
| | (N) literresent ftr: $[++2]$ un: $(A = 0)$ |
| • | 11.6 rt 0 th 8 np 7 ch 0 co 7_{0} ec 7_{1} |
| | r:6, su:9, nu:7) |
| 7. | $(F \text{sub-} \Lambda \text{ftr} \cdot [+-2] \text{nn} \cdot 7 \text{cb} \cdot 0 \text{co} \cdot 7_1)$ |
| 71 71 | (E sub: B ftr: $[+-2]$ np: 7 ch: 12 co: 0) |
| 'b 8 | (1, 3ub.b, 1cl.[1], np.7, cl.12a, co.0) |
| 9 | $(N = 1)^{+} (1 +$ |
| | 1 + 0 $r + 10$ $t + 10$ $n + 9$ $c + 0$ $c + 0$ |
| | pr.7 su.11 nu.9) |
| 95 | $(E_{1}, Sub: A_{2}, ftr: [++-], pp: 9, ch: 0, co: 0)$ |
| 10 | (S, up: 8, dn: 11, 1t: 9, rt: 0, th: 11, nu: 10) |
| 11 | (N. lit:she. ftr: $[++++-]$. up:10. dn:0. |
| | 1t:0, rt:12, th:12, np:11, ch:0, co:11, ec:11, |
| | pr:9, su:12, nu:11) |
| 11a | (E, sub:A, ftr: $[+-+++-]$, np:11, ch:0, co:0) |
| 12 | (N. lit:it. ftr: $[++-?]$, up:10. dn:0. |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_{9} , ec: 12_{9} , |
| | pr:11, su:0, nu:12) |
| 12 _a | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(she, John)
            sc(she, John) = True
            agr(she, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(he)
    non_refl_chaining(he)
        chaining_n_to_n(he, it)
            sc(he, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, she)
             sc(he, she) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, present)
            sc(he, present) = True
            agr(he, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, June)
            sc(he, June) = True
            agr(he, June) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, \phi)
            sc(he, \phi) = True
            agr(he, \phi) = True
            rnr(he, \phi) = True
            chaining_e_to_n(hea, \phi)
                 agr(hea, \phi) = True
                 new_chain(hea, \phi)
                     new_chain: create \phi_{\rm d}
                     new_chain: create \phi_d hea
```

| | Nodes |
|----------------|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_a , |
| | pr:0, su:5, nu:3) |
| 3 _a | (E, sub:A, ftr:[+], np:3, ch:0, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: 5_a , ec: 5_d , |
| | pr:3, su:6, nu:5) |
| 5 _a | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| 5 _b | (E, sub:B, ftr:[++-?], np:5, ch: 12_{a} , co: 5_{c}) |
| 5 _c | (E, sub:C, ftr:[++-++-], np:5, ch: 11_{a} , co: 5_{d}) |
| 54 | (E, sub:D, ftr:[++-], np:5, ch:9a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| | pr:5, su:7, nu:6) |
| 6a | (E, sub:A, ftr:[+-++-], np:6, ch:0, co:6 _b) |
| 6 _b | (E, sub:B, ftr:[+-++-], np:6, ch:11 _a , co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | lt:6, rt:0, th:8, np:7, ch:0, co:7 _a , ec:7 _b , |
| | pr:6, su:9, nu:7) |
| 7_{a} | (E, sub:A, ftr:[+-?], np:7, ch:0, co:7 _b) |
| 7 _b | (E, sub:B, ftr:[+-?], np:7, ch:12a, co:0) |
| 8 | (S, up:1, dn:9, lt:2, rt:0, th:9, nu:8) |
| 9 | (N, lit:he, ftr:[++-], up:8, dn:0, |
| | lt:0, rt:10, th:10, np:9, ch:0, co: 9_a , ec: 9_a , |
| | pr:7, su:11, nu:9) |
| 9 _a | (E, sub:A, ftr:[++], np:9, ch:0, co:0) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, co:11 $_a$, ec:11 $_a$, |
| | pr:9, su:12, nu:11) |
| 11_{a} | (E, sub:A, ftr:[++-+-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_a , ec: 12_a , |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

new_chain: exiting chaining_e_to_n: exiting chaining_n_to_n: exiting chaining_n_to_n(he, John) sc(he, John) = True agr(he, John) = True rnr(he, John) = True chaining_e_to_n(hea, John) agr(hea, John) = True new_chain(hea, John) new_chain: create John_b new_chain: create John_b^hea

| Nodes | |
|----------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:3 _a , ec:3 _b , |
| | pr:0, su:5, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-], np:3, ch:0, co:3 _b) |
| 3 _b | (E, sub:B, ftr:[+-], np:3, ch:9a, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m d}$, |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| ^{5}b | (E, sub:B, ftr:[++-?], np:5, ch: 12_a , co: 5_c) |
| 5c | (E, sub:C, ftr:[++-+-], np:5, $ch:11_a$, $co:5_d$) |
| ^{5}d | (E, sub:D, ftr:[++-], np:5, ch:9a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | It:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| 6 | pr:5, su:/, nu:6) |
| ba c | (E, sub:A, itr:[+-++-], np:6, ch:0, $co:o_b$) |
| °b | (E, SUD:B, Itr: $[+++-]$, np:6, Ch:II _a , CO:U) |
| | (N, III:present, III:[++?], up:4, an:0, |
| | 11:0, 11:0, 11:0, 11:0, 10:1, 00:1a, ec:1b, |
| 7. | $(F \text{sub-} \Lambda \text{ftr} \cdot [++2] \text{nn} \cdot 7 \text{cb} \cdot 0 \text{co} \cdot 7,)$ |
| 71 71 | (E sub B ftr: $[+-2]$ np: 7 ch: 12 co: 0) |
| 'b 8 | (S, up:1, dn:9, lt:2, rt:0, th:9, nu:8) |
| 9 | (N. lit:he. ftr: $[++-+-]$. up:8. dn:0. |
| | lt:0, rt:10, th:10, np:9, ch:0, co: 9_{a} , ec: 9_{a} , |
| | pr:7, su:11, nu:9) |
| 9a | (E, sub:A, ftr:[++-], np:9, ch:0, co:0) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, co:11a, ec:11a, |
| | pr:9, su:12, nu:11) |
| 11_{a} | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_a , ec: 12_a , |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)
    non refl chaining (\phi)
         chaining_n_to_n(\phi, it)
             sc(\phi, it) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, she)
             sc(\phi, she) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, he)
             sc(\phi, he) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, present)
             sc(\phi, present) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, June)
             sc(\phi, June) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = True
             rnr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create Johnc
                      new_chain: create John<sub>C</sub>^{\phi}a
```
| | Nodes |
|------------------|---|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co: 3_a , ec: 3_c , |
| | pr:0, su:5, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-], np:3, ch:0, co:3 _b) |
| 3 _b | (E, sub:B, ftr:[+], np:3, ch: 9_a , co: 3_c) |
| 3 _c | (E, sub:C, ftr:[+], np:3, ch:5a, co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m d}$, |
| | pr:3, su:6, nu:5) |
| 5_{a} | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| $^{5}\mathrm{b}$ | (E, sub:B, ftr:[++-?], np:5, ch: 12_a , co: 5_c) |
| $5_{\mathbf{C}}$ | (E, sub:C, ftr:[++-+-], np:5, ch: 11_a , co: 5_d) |
| 5 d | (E, sub:D, ftr:[++-], np:5, ch:9a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | lt:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| | pr:5, su:7, nu:6) |
| 6a | (E, sub:A, ftr:[+-++-], np:6, ch:0, co:6 _b) |
| ⁶ b | (E, sub:B, ftr:[+-++-], np:6, ch:ll _a , co:0) |
| 7 | (N, lit:present, ftr:[+-?], up:4, dn:0, |
| | It:6, rt:0, th:8, np:7, ch:0, co: 7_a , ec: 7_b , |
| - | pr:6, su:9, nu:7) |
| 7a | (E, sub:A, $ttr:[+-?]$, $np:7$, $ch:0$, $co:7_b$) |
| (b | (E, sub:B, itr: $[+-?]$, np: /, ch: I_{2a} , co: 0) |
| 8 | (S, up:1, an:9, It:2, rt:0, tn:9, nu:8) |
| 9 | (N, III: ne, III: [+-+-+-], up:0, an:0, |
| | pr:7 $qu:11$ $pu:9$ |
| 9. | $(F \text{sub-} \Lambda \text{ftr} \cdot [+-++-] \text{nn} \cdot 9 \text{cb} \cdot 0 \text{co} \cdot 0)$ |
| 10 | (E, up:8 dp:11 lt:9 rt:0 th:11 nu:10) |
| 11 | (N, lit:she, ftr:[+-+++-], up:10, dp:0. |
| | 1t:0, rt:12, th:12, np:11, ch:0, co:11a, ec:11a, |
| | pr:9, su:12, nu:11) |
| 11a | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_{a} , ec: 12_{a} , |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

```
\begin{array}{ll} {\rm new\_chain: exiting} \\ {\rm chaining\_e\_to\_n: exiting} \\ {\rm chaining\_e\_to\_n}(\phi_{\rm b}, \, {\rm John}) \\ {\rm agr}(\phi_{\rm b}, \, {\rm John}) = {\rm False} \\ {\rm chaining\_e\_to\_n: exiting} \\ {\rm chaining\_e\_to\_n}(\phi_{\rm C}, \, {\rm John}) \\ {\rm agr}(\phi_{\rm C}, \, {\rm John}) = {\rm False} \\ {\rm chaining\_e\_to\_n: exiting} \\ {\rm chaining\_e\_to\_n: exiting} \\ {\rm chaining\_e\_to\_n}(\phi_{\rm d}, \, {\rm John}) \\ {\rm agr}(\phi_{\rm d}, \, {\rm John}) = {\rm True} \\ {\rm new\_chain}(\phi_{\rm d}, \, {\rm John}) \\ {\rm new\_chain: create } {\rm John_d}^{-}\phi_{\rm d} \\ \end{array}
```

| | Nodes |
|------------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:3a, ec:3 _d , |
| | pr:0, su:5, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-], np:3, ch:0, co:3 _b) |
| 3_{b} | (E, sub:B, ftr:[+], np:3, ch: 9_a , co: 3_c) |
| $3_{\rm C}$ | (E, sub:C, ftr:[+-], np:3, ch: 5_a , co: 3_d) |
| ^{3}d | (E, sub:D, ftr:[+-], np:3, ch:5 _d , co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m d}$, |
| | pr:3, su:6, nu:5) |
| $\mathbf{5_{a}}$ | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| $^{5}\mathrm{b}$ | (E, sub:B, ftr:[++-?], np:5, $ch:12_a$, $co:5_c$) |
| $5_{\rm C}$ | (E, sub:C, ftr:[++-+-], np:5, $ch:11_a$, $co:5_d$) |
| ^{5}d | (E, sub:D, ftr:[++-], np:5, ch:9a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | It:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| | pr:5, su:7, nu:6) |
| 6a | (E, sub:A, $tr:[+++-]$, np:6, ch:0, co: b_b) |
| ⁶ b | (E, sub:B, $tr:[+++-]$, np:6, ch:II _a , co:0) |
| 7 | (N, lit:present, itr:[+-?], up:4, dn:0, |
| | It:6, rt:0, tn:8, np:/, cn:0, co: i_a , ec: i_b , |
| 7 | $p_{1:0}, s_{1:0}, n_{1:1}$ |
| 7a 7. | (E, Sub:R, It:: $[+-?]$, np:7, ch:0, co: n_b) (E sub:R ftr: $[++-?]$, np:7, ch:12, co: n_b) |
| 'b 8 | (E, Sub.D, ICI.[1 :], np.7, Ch.12a, CO.0) |
| 9 | $(N \ \text{lit} \cdot \text{he} \ \text{ftr} \cdot [+-++-] \ \text{up} \cdot 8 \ \text{dn} \cdot 0$ |
| U | 1 + 0 $r + 10$ $t + 10$ $n + 9$ $c + 0$ $c + 0$ $c + 0$ |
| | pr:7, su:11, nu:9) |
| 9a | (E, sub:A, ftr:[++-], np:9, ch:0, co:0) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, co:11 _a , ec:11 _a , |
| | pr:9, su:12, nu:11) |
| 11_{a} | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_a , ec: 12_a , |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

| | Nodes |
|------------------|--|
| 1 | (C, up:0, dn:2, lt:0, rt:0, th:2, nu:1) |
| 2 | (S, up:1, dn:3, lt:0, rt:8, th:3, nu:2) |
| 3 | (N, lit:John, ftr:[+], up:2, dn:0, |
| | lt:0, rt:4, th:4, np:3, ch:0, co:3a, ec:3 _d , |
| | pr:0, su:5, nu:3) |
| 3_{a} | (E, sub:A, ftr:[+-], np:3, ch:0, co:3 _b) |
| 3_{b} | (E, sub:B, ftr:[+], np:3, ch: 9_a , co: 3_c) |
| $3_{\rm C}$ | (E, sub:C, ftr:[+-], np:3, ch: 5_a , co: 3_d) |
| ^{3}d | (E, sub:D, ftr:[+-], np:3, ch:5 _d , co:0) |
| 4 | (S, up:2, dn:5, lt:3, rt:0, th:5, nu:4) |
| 5 | (N, lit: ϕ , ftr:[+?????-], up:4, dn:0, |
| | lt:0, rt:6, th:6, np:5, ch:0, co: $5_{ m a}$, ec: $5_{ m d}$, |
| | pr:3, su:6, nu:5) |
| $\mathbf{5_{a}}$ | (E, sub:A, ftr:[+?????-], np:5, ch:0, co:5 _b) |
| $^{5}\mathrm{b}$ | (E, sub:B, ftr:[++-?], np:5, $ch:12_a$, $co:5_c$) |
| $5_{\rm C}$ | (E, sub:C, ftr:[++-+-], np:5, $ch:11_a$, $co:5_d$) |
| ^{5}d | (E, sub:D, ftr:[++-], np:5, ch:9a, co:0) |
| 6 | (N, lit:June, ftr:[+-++-], up:4, dn:0, |
| | It:5, rt:7, th:7, np:6, ch:0, co: 6_a , ec: 6_b , |
| | pr:5, su:7, nu:6) |
| 6a | (E, sub:A, $tr:[+++-]$, np:6, ch:0, co: b_b) |
| ⁶ b | (E, sub:B, $tr:[+++-]$, np:6, ch:II _a , co:0) |
| 7 | (N, lit:present, itr:[+-?], up:4, dn:0, |
| | It:6, rt:0, tn:8, np:/, cn:0, co: i_a , ec: i_b , |
| 7 | $p_{1:0}, s_{1:0}, n_{1:1}$ |
| 7a 7. | (E, Sub:R, It:: $[+-?]$, np:7, ch:0, co: n_b) (E sub:R ftr: $[++-?]$, np:7, ch:12, co: n_b) |
| 'b 8 | (E, Sub.D, ICI.[1 :], np.7, Ch.12a, CO.0) |
| 9 | $(N \ \text{lit} \cdot \text{he} \ \text{ftr} \cdot [+-++-] \ \text{up} \cdot 8 \ \text{dn} \cdot 0$ |
| U | 1 + 0 $r + 10$ $t + 10$ $n + 9$ $c + 0$ $c + 0$ $c + 0$ |
| | pr:7, su:11, nu:9) |
| 9a | (E, sub:A, ftr:[++-], np:9, ch:0, co:0) |
| 10 | (S, up:8, dn:11, lt:9, rt:0, th:11, nu:10) |
| 11 | (N, lit:she, ftr:[++-++-], up:10, dn:0, |
| | lt:0, rt:12, th:12, np:11, ch:0, co:11 _a , ec:11 _a , |
| | pr:9, su:12, nu:11) |
| 11_{a} | (E, sub:A, ftr:[++-++-], np:11, ch:0, co:0) |
| 12 | (N, lit:it, ftr:[++-?], up:10, dn:0, |
| | lt:11, rt:0, th:0, np:12, ch:0, co: 12_a , ec: 12_a , |
| | pr:11, su:0, nu:12) |
| 12_{a} | (E, sub:A, ftr:[++-?], np:12, ch:0, co:0) |

The previous example should give enough details away to satisfy the reader's curiosity, but the form of the previous example is rather burdensome. From now on, we'll keep to a more concise, if less detailed, output. To indicate a chain_link between two E-nodes, we use the symbol ^.

Below are some more examples.

(10.2) Janet saw herself.



| Features | | | | | | | | | | |
|---------------------------------|---------------------|---|---|---|---|---|---|---|--|--|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | | | | |
| Janet | - | - | - | + | - | + | + | - | | |
| herself | herself + + - + + + | | | | | | | | | |

```
chaining
init_table
```

| Chaining | | | | | |
|----------|----------|--|--|--|--|
| Janet | herself | | | | |
| Janeta | herselfa | | | | |

```
init_table: exiting
chaining_n(herself)
  refl_chaining(herself)
    simplex_pred(herself)
    simplex_pred: Janet
    chaining_n_to_n(herself, Janet)
        sc(herself, Janet) = True
        agr(herself, Janet) = True
        rnr(herself, Janet) = True
        chaining_e_to_n(herself_a, Janet)
        agr(herself_a, Janet) = True
        new_chain(herself_a, Janet)
        new_chain: create Janet_b
        new_chain: create Janet_b^herself_a
```

| Chaining | | | | | | | |
|------------------------------|----------|--|--|--|--|--|--|
| Janet | herself | | | | | | |
| Janeta | herselfa | | | | | | |
| Janet _b ^herselfa | | | | | | | |

| Chaining | | | | | | | |
|------------------------------|----------|--|--|--|--|--|--|
| Janet | herself | | | | | | |
| Janeta | herselfa | | | | | | |
| Janet _b ^herselfa | | | | | | | |

 $(10.3)\,$ Janet saw her.



| Features | | | | | | | | | | |
|---------------------------------|---|---|---|---|---|---|---|---|--|--|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | | | | |
| Janet | - | - | - | + | - | + | + | - | | |
| her | + | _ | - | + | - | + | + | _ | | |

| Chaining | | | | | |
|-----------|------|--|--|--|--|
| Janet her | | | | | |
| Janeta | hera | | | | |

```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
    chaining_n_to_n(her, Janet)
        sc(her, Janet) = True
        agr(her, Janet) = True
        rnr(her, Janet) = True
        chaining_e_to_n(hera, Janet)
        agr(hera, Janet) = True
        new_chain(hera, Janet)
        new_chain: create Janetb
        new_chain: create Janetb^hera
```

| Chaining | | | | | | |
|--------------------------|------|--|--|--|--|--|
| Janet | her | | | | | |
| Janeta | hera | | | | | |
| Janet _b ^hera | | | | | | |

| Chaining | | | | | | | |
|--------------------------|------|--|--|--|--|--|--|
| Janet | her | | | | | | |
| Janeta | hera | | | | | | |
| Janet _b ^hera | | | | | | | |

(10.4) *Janet saw himself.



| Features | | | | | | | | | |
|----------|---------------------------------|---|---|---|---|---|---|---|--|
| | PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | | |
| Janet | - | _ | - | + | - | + | + | _ | |
| himself | + | - | - | + | - | - | + | + | |

| Chaining | | | | |
|---------------|----------|--|--|--|
| Janet himself | | | | |
| Janeta | himselfa | | | |

```
init_table: exiting
chaining_n(himself)
    refl_chaining(himself)
    simplex_pred(himself)
    simplex_pred: Janet
    chaining_n_to_n(himself, Janet)
        sc(himself, Janet) = True
        agr(himself, Janet) = False
        chaining_n_to_n: exiting
        simplex_pred(Janet)
        simplex_pred:
        refl_chaining: exiting
        chaining_n: exiting
        chaining_n: exiting
        chaining_: exiting
        chaining: exiting
        chaining: exiting
        chaining: exiting
        chaining: exiting
```

| Chaining | | | | |
|---------------|----------|--|--|--|
| Janet himself | | | | |
| Janeta | himselfa | | | |

Examples (10.5)-(10.11) are from Lees and Klima [22].

(10.5) The men threw a smokescreen around themselves.



the men threw a smokescreen around themselves

| Features | | | | | | | | |
|---------------------------------|---|---|---|---|---|---|---|---|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | | |
| men | - | _ | - | + | + | - | + | - |
| smokescreen | - | - | - | + | - | ? | - | - |
| themselves | + | - | - | + | + | ? | ? | + |

| Chaining | | | | | | |
|----------|--------------|-------------|--|--|--|--|
| men | smokescreen | themselves | | | | |
| mena | smokescreena | themselvesa | | | | |

```
init_table: exiting
chaining_n(themselves)
    refl_chaining(themselves)
        simplex_pred(themselves)
        simplex_pred: smokescreen
        chaining_n_to_n(themselves, smokescreen)
            sc(themselves, smokescreen) = True
            agr(themselves, smokescreen) = False
        chaining_n_to_n: exiting
        simplex_pred(smokescreen)
        simplex_pred: men
        chaining_n_to_n(themselves, men)
            sc(themselves, men) = True
            agr(themselves, men) = True
            rnr(themselves, men) = True
            chaining_e_to_n(themselvesa, men)
                agr(themselves_a, men) = True
                new_chain(themselvesa, men)
                    new_chain: create menb
                    new_chain: create men<sub>b</sub>^themselvesa
```

| Chaining | | | | | | |
|-------------------------------|--------------|-------------|--|--|--|--|
| men | smokescreen | themselves | | | | |
| mena | smokescreena | themselvesa | | | | |
| men _b ^themselvesa | | | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
simplex_pred(men)
simplex_pred:
refl_chaining: exiting
chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | |
|-------------------------------|--------------|-------------|--|--|--|--|
| men | smokescreen | themselves | | | | |
| mena | smokescreena | themselvesa | | | | |
| men _b ^themselvesa | | | | | | |

(10.6) The men found a smokescreen around them.



the men threw a smokescreen around them

| Features | | | | | | | | |
|---------------------------------|---|---|---|---|---|---|-----|---|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | RPF | |
| men | - | - | - | + | + | - | + | _ |
| smokescreen | - | - | - | + | - | ? | - | - |
| them | + | - | - | + | + | ? | ? | - |

| Chaining | | | | | |
|----------|--------------|-------|--|--|--|
| men | smokescreen | them | | | |
| mena | smokescreena | thema | | | |

```
init_table: exiting
chaining_n(them)
    non_refl_chaining(them)
    chaining_n_to_n(them, smokescreen) = True
        agr(them, smokescreen) = True
        agr(them, smokescreen) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, men)
        sc(them, men) = True
        agr(them, men) = True
        rnr(them, men) = False
        chaining_n_to_n: exiting
        non_refl_chaining: exiting
        chaining_n: exiting
        chaining: exiting
        chaining: exiting
```

| Chaining | | | | | |
|----------|--------------|-------|--|--|--|
| men | smokescreen | them | | | |
| mena | smokescreena | thema | | | |

(10.7) The men found a smokescreen to be around them.



| Features | | | | | | | | |
|---------------------------------|---|---|---|---|---|---|---|-----|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | | | RPF |
| men | - | - | - | + | + | - | + | - |
| smokescreen | - | - | - | + | - | ? | - | - |
| them | + | - | - | + | + | ? | ? | _ |

```
chaining
init_table
```

| Chaining | | | | | |
|----------|--------------|-------|--|--|--|
| men | smokescreen | them | | | |
| mena | smokescreena | thema | | | |

```
init_table: exiting
chaining_n(them)
    non_refl_chaining(them)
        chaining_n_to_n(them, smokescreen)
            sc(them, smokescreen) = True
            agr(them, smokescreen) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, men)
            sc(them, men) = True
            agr(them, men) = True
            rnr(them, men) = True
            chaining_e_to_n(thema, men)
                agr(them_a, men) = True
                new_chain(thema, men)
                    new_chain: create menb
                    new_chain: create menb^thema
```

| Chaining | | | | | | |
|-------------------------|--------------|-------|--|--|--|--|
| men | smokescreen | them | | | | |
| mena | smokescreena | thema | | | | |
| men _b ^thema | | | | | | |

| Chaining | | | | |
|-------------------------|--------------|-------|--|--|
| men | smokescreen | them | | |
| mena | smokescreena | thema | | |
| men _b ^thema | | | | |

 $\left(10.8\right)$ The men found a smokes creen and it was around them.



| Features | | | | | | | | |
|---------------------------------|---|---|---|---|---|-----|---|---|
| PNF FPF SPF TPF PLF GNF ANF RPF | | | | | | RPF | | |
| men | - | - | - | + | + | - | + | - |
| smokescreen | - | _ | - | + | - | ? | - | - |
| it | + | - | - | + | - | ? | - | - |
| them | + | _ | - | + | + | ? | ? | _ |

| Chaining | | | | |
|----------|--------------|-----|-------|--|
| men | smokescreen | it | them | |
| mena | smokescreena | ita | thema | |

```
init_table: exiting
chaining_n(them)
    non_refl_chaining(them)
        chaining_n_to_n(them, it)
            sc(them, it) = True
            agr(them, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, smokescreen)
            sc(them, smokescreen) = True
            agr(them, smokescreen) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, men)
            sc(them, men) = True
            agr(them, men) = True
            rnr(them, men) = True
            chaining_e_to_n(thema, men)
                agr(them_a, men) = True
                new_chain(thema, men)
                    new_chain: create menb
                    new_chain: create menb thema
```

| Chaining | | | | | |
|-------------------------|--------------|-----|-------|--|--|
| men | smokescreen | it | them | | |
| mena | smokescreena | ita | thema | | |
| men _b ^thema | | | | | |

```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
   non_refl_chaining: exiting
chaining_n: exiting
chaining_n(it)
    non refl chaining(it)
        chaining_n_to_n(it, them)
            sc(it, them) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, smokescreen)
            sc(it, smokescreen) = True
            agr(it, smokescreen) = True
            rnr(it, smokescreen) = True
            chaining_e_to_n(ita, smokescreen)
                agr(ita, smokescreen) = True
                new_chain(ita, smokescreen)
                    new_chain: create smokescreenb
                    new_chain: create smokescreenb^ita
```

| Chaining | | | | | |
|-------------------------|-------------------------------|------|-------|--|--|
| men | it | them | | | |
| mena | smokescreena | ita | thema | | |
| men _b ^thema | smokescreen _b ^ita | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(it, men)
sc(it, men) = True
agr(it, men) = False
chaining_n_to_n: exiting
non_refl_chaining: exiting
chaining_n: exiting
chaining: exiting
```

| Chaining | | | | |
|-------------------------|---|------|-------|--|
| men | it | them | | |
| mena | smokescreena | ita | thema | |
| men _b ^thema | men _b ^thema smokescreen _b ^ita | | | |

(10.9) I told John to protect himself.



| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| I | + | + | - | - | - | ? | + | - |
| John | _ | - | _ | + | - | - | + | _ |
| ϕ | + | ? | ? | ? | ? | ? | ? | _ |
| himself | + | - | - | + | - | - | + | + |

| Chaining | | | | |
|----------|-------|------------|----------|--|
| I | John | ϕ | himself | |
| Ia | Johna | ϕ_{a} | himselfa | |

```
init_table: exiting
chaining_n(himself)
refl_chaining(himself)
simplex_pred(himself)
simplex_pred: $\phi$
chaining_n_to_n(himself, $\phi$)
sc(himself, $\phi$) = True
agr(himself, $\phi$) = True
rnr(himself, $\phi$) = True
chaining_e_to_n(himself_a, $\phi$)
agr(himself_a, $\phi$) = True
new_chain(himself_a, $\phi$)
new_chain: create $\phi_b$
himself_a
```

| Chaining | | | | | |
|----------|-------|-------------------------|----------|--|--|
| I | John | ϕ | himself | | |
| Ia | Johna | $\phi_{ m b}$ ^himselfa | himselfa | | |

```
new_chain: exiting
              chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         simplex_pred(\phi)
         simplex_pred:
    refl_chaining: exiting
chaining_n: exiting
\texttt{chaining_n}(\phi)
    non_refl_chaining(\phi)
         chaining_n_to_n(\phi, himself)
              sc(\phi, himself) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
              sc(\phi, John) = True
              agr(\phi, John) = True
              rnr(\phi, John) = True
              chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = True
                  new_chain(\phi_{a}, John)
                       new_chain: create Johnb
                       new_chain: create John<sub>b</sub>^{\phi}a
```

| | Chaining | | | | | |
|----|--------------------------------|------------------------|----------|--|--|--|
| I | John | ϕ | himself | | | |
| Ia | Johna | ϕ_{a} | himselfa | | | |
| | John _b ^ ϕ_{a} | $\phi_{ m b}$ himselfa | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_e_to_n(\phi_b, John)
agr(\phi_b, John) = True
new_chain(\phi_b, John)
new_chain: create John<sub>c</sub>
new_chain: create John<sub>c</sub>^\phi_b
```

| | Chaining | | | | | |
|----|--------------------------------|--------------------------|----------|--|--|--|
| I | John | ϕ | himself | | | |
| Ia | Johna | ϕ_{a} | himselfa | | | |
| | $John_b^{\phi_a}$ | $\phi_{\rm b}$ ^himselfa | | | | |
| | $John_{\rm C}^{~}\phi_{\rm b}$ | | | | | |

```
\begin{array}{c} \texttt{new\_chain: exiting}\\ \texttt{chaining\_e\_to\_n: exiting}\\ \texttt{chaining\_n\_to\_n: exiting}\\ \texttt{chaining\_n\_to\_n}(\phi, \ \texttt{I})\\ \texttt{sc}(\phi, \ \texttt{I}) = \texttt{True}\\ \texttt{agr}(\phi, \ \texttt{I}) = \texttt{True}\\ \texttt{agr}(\phi, \ \texttt{I}) = \texttt{True}\\ \texttt{rnr}(\phi, \ \texttt{I}) = \texttt{True}\\ \texttt{chaining\_e\_to\_n}(\phi_a, \ \texttt{I})\\ \texttt{agr}(\phi_a, \ \texttt{I}) = \texttt{True}\\ \texttt{new\_chain}(\phi_a, \ \texttt{I})\\ \texttt{new\_chain: create } \ \texttt{Ib}\\ \texttt{new\_chain: create } \ \texttt{Ib}^{+}\phi_a \end{array}
```

| Chaining | | | | | |
|----------------|------------------------------------|--------------------------|----------|--|--|
| I | John | ϕ | himself | | |
| Ia | Johna | ϕ_{a} | himselfa | | |
| $I_b^{\phi_a}$ | John _b ^ $\phi_{\rm a}$ | $\phi_{\rm b}$ ^himselfa | | | |
| | ${\rm John_C}^{\circ}\phi_{\rm b}$ | | | | |

```
new_chain: exiting
                chaining_e_to_n: exiting
                chaining_e_to_n(\phi_b, I)
                    agr(\phi_b, I) = False
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, himself)
                sc(I, himself) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | |
|----------------|------------------------------------|--------------------------|----------|--|--|--|
| I | John | ϕ | himself | | | |
| Ia | Johna | ϕ_{a} | himselfa | | | |
| $I_b^{\phi_a}$ | $John_b^{\phi_a}$ | $\phi_{\rm b}$ ^himselfa | | | | |
| | ${\rm John_C}^{\circ}\phi_{\rm b}$ | | | | | |

(10.10) I told John to protect me.



| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| I | + | + | - | - | - | ? | + | - |
| John | - | - | - | + | - | - | + | - |
| ϕ | + | ? | ? | ? | ? | ? | ? | - |
| me | + | + | - | - | - | ? | + | - |

```
chaining
```

```
init_table
```

| Chaining | | | | | |
|----------|-------|------------|-----|--|--|
| I | me | | | | |
| Ia | Johna | ϕ_{a} | mea | | |

```
init_table: exiting
chaining_n(me)
    non_refl_chaining(me)
        chaining_n_to_n(me, \phi)
             sc(me, \phi) = True
             agr(me, \phi) = True
             rnr(me, \phi) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(me, John)
             sc(me, John) = True
             agr(me, John) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(me, I)
             sc(me, I) = True
             agr(me, I) = True
             rnr(me, I) = True
             chaining_e_to_n(mea, I)
                 agr(me_a, I) = True
                 new_chain(mea, I)
                     new_chain: create {\rm I}_{\rm b}
                     new_chain: create Ib^mea
```

| Chaining | | | | | |
|---------------------|-------|------------|-----|--|--|
| I | John | ϕ | me | | |
| Ia | Johna | ϕ_{a} | mea | | |
| I _b ^mea | | | | | |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)
    non_refl_chaining(\phi)
         chaining_n_to_n(\phi, me)
             sc(\phi, me) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = True
             rnr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create Johnb
                      new_chain: create John<sub>b</sub>^{\phi}a
```

| Chaining | | | | | | |
|---------------------|------------------------------|------------|-----|--|--|--|
| I | John | ϕ | me | | | |
| Ia | Johna | ϕ_{a} | mea | | | |
| I _b ^mea | John _b $^{\phi}a$ | | | | | |

```
\begin{array}{c} \text{new\_chain: exiting} \\ \text{chaining\_e\_to\_n: exiting} \\ \text{chaining\_n\_to\_n: exiting} \\ \text{chaining\_n\_to\_n}(\phi, \ \textbf{I}) \\ \text{sc}(\phi, \ \textbf{I}) = \text{True} \\ \text{agr}(\phi, \ \textbf{I}) = \text{True} \\ \text{rnr}(\phi, \ \textbf{I}) = \text{True} \\ \text{chaining\_e\_to\_n}(\phi_a, \ \textbf{I}) \\ \text{agr}(\phi_a, \ \textbf{I}) = \text{True} \\ \text{new\_chain}(\phi_a, \ \textbf{I}) \\ \text{new\_chain: create } \ \textbf{I}_c \\ \text{new\_chain: create } \ \textbf{I}_c \hat{\phi}_a \end{array}
```

| Chaining | | | | | | |
|-------------------------|-------------------|------------|-----|--|--|--|
| I | John | ϕ | me | | | |
| Ia | Johna | ϕ_{a} | mea | | | |
| I _b ^mea | $John_b^{\phi_a}$ | | | | | |
| $I_{c}^{\phi} \phi_{a}$ | | | | | | |

```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, me)
                sc(I, me) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | |
|-------------------------|---|------------|-----|--|--|
| I | John | ϕ | me | | |
| Ia | Johna | ϕ_{a} | mea | | |
| I _b ^mea | ${\tt John}_{\sf b}{}^{\uparrow}\phi_{\sf a}$ | | | | |
| $I_{c}^{\phi} \phi_{a}$ | | | | | |

(10.11) I told John to protect myself.



| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| I | + | + | - | - | _ | ? | + | - |
| John | - | - | - | + | - | - | + | - |
| ϕ | + | ? | ? | ? | ? | ? | ? | - |
| myself | + | + | - | - | _ | ? | + | + |

| Chaining | | | | |
|----------|-------|------------|---------|--|
| I | John | ϕ | myself | |
| Ia | Johna | ϕ_{a} | myselfa | |

```
init_table: exiting
chaining_n(myself)
    refl_chaining(myself)
    simplex_pred(myself)
    simplex_pred: \phi
    chaining_n_to_n(myself, \phi)
    sc(myself, \phi) = True
    agr(myself, \phi) = True
    rnr(myself, \phi) = True
    chaining_e_to_n(myself_a, \phi)
    agr(myself_a, \phi) = True
    new_chain(myself_a, \phi)
    new_chain: create \phi_b
    new_chain: create \phi_b^{-}myselfa
```

| Chaining | | | | | |
|----------|-------|-------------------------------------|---------|--|--|
| I | John | ϕ | myself | | |
| Ia | Johna | $\phi_{\rm b}$ ^myself _a | myselfa | | |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         simplex_pred(\phi)
         simplex_pred:
    refl_chaining: exiting
chaining n: exiting
chaining_n(\phi)
    non_refl_chaining(\phi)
         chaining_n_to_n(\phi, myself)
              sc(\phi, myself) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = True
             rnr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create Johnb
                       new_chain: create John<sub>b</sub>^{\phi}a
```

| | Chaining | | | | | | |
|----|--------------------------------|-------------------------|---------|--|--|--|--|
| I | John | ϕ | myself | | | | |
| Ia | Johna | ϕ_{a} | myselfa | | | | |
| | John _b ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | | |

```
\begin{array}{rl} & \text{new\_chain: exiting} \\ & \text{chaining\_e\_to\_n: exiting} \\ & \text{chaining\_e\_to\_n(\phi_b, \text{ John})} \\ & \text{agr}(\phi_b, \text{ John}) = \text{False} \\ & \text{chaining\_e\_to\_n: exiting} \\ & \text{chaining\_n\_to\_n: exiting} \\ & \text{chaining\_n\_to\_n(\phi, I)} \\ & \text{sc}(\phi, I) = \text{True} \\ & \text{agr}(\phi, I) = \text{True} \\ & \text{agr}(\phi, I) = \text{True} \\ & \text{chaining\_e\_to\_n(\phi_a, I)} \\ & \text{agr}(\phi_a, I) = \text{True} \\ & \text{new\_chain(\phi_a, I)} \\ & \text{new\_chain: create I}_b \\ & \text{new\_chain: create I}_b^{\phi_a} \end{array}
```

| Chaining | | | | | | |
|----------------|--|-------------------------|---------|--|--|--|
| I | John | ϕ | myself | | | |
| Ia | Johna | ϕ_{a} | myselfa | | | |
| $I_b^{\phi_a}$ | ${\tt John}_{\sf b}{}^{\circ}\phi_{\sf a}$ | $\phi_{\rm b}$ ^myselfa | | | | |

new_chain: exiting chaining_e_to_n: exiting chaining_e_to_n(ϕ_b , I) agr(ϕ_b , I) = True new_chain(ϕ_b , I) new_chain: create I_c new_chain: create I_c^ ϕ_b

| Chaining | | | | | | |
|------------------|--------------------------------|-------------------------|---------|--|--|--|
| I | John | ϕ | myself | | | |
| Ia | Johna | ϕ_{a} | myselfa | | | |
| $I_b^{\phi}a$ | John _b ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | |
| $I_{c} \phi_{b}$ | | | | | | |

```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, myself)
                sc(I, myself) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | |
|------------------|--------------------------------|-------------------------|---------|--|--|--|
| I | John | ϕ | myself | | | |
| Ia | Johna | ϕ_{a} | myselfa | | | |
| $I_b^{\phi_a}$ | John _b ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | |
| $I_{c} \phi_{b}$ | | | | | | |

11 Table Interpreter

The Table Interpreter module defines function interpret and has the form shown in Figure 11.1.

```
#table_interp;
from globals import *
def interpret(nnodes: list[Node]) -> list[list[list[Node]]]:
```

Figure 11.1. Skeleton of the Table Interpreter

Basically, after the chaining table is created, a number of chains are implicitly defined by the chaining table and it is the job of the Table Interpreter to mesh these chains back into copies of the system tree, returning all trees defined by legitimate interpretations.

A nonpronominal E-node with the E-nodes that are traced by walking down chain_link's until nil chain_link is reached constitute a **chain**. A set of chains defined by the chaining table which cover all the pronominal N-nodes and do not intersect constitute a legitimate **interpretation**.

Take the table given in Figure 11.2 as an example.

| Chaining | | | | | | | | |
|---------------------------------------|---------------------------------|-------------------------|---------------------------|-----|------|------|--|--|
| John | ϕ | June | present | he | she | it | | |
| Johna | ϕ_{a} | Junea | present _a | hea | shea | lita | | |
| John _b ^hea | $\phi_{ m b}$ îta | June _b ^shea | present _b ^ita | | | | | |
| ${\tt John_C}^{\uparrow}\phi_{\sf A}$ | $\phi_{\rm C}$ $^{\rm she_{a}}$ | | | | | | | |
| ${\tt John_d}^{}\phi_{d}$ | $\phi_{\rm d}$ hea | | | | | | | |

Figure 11.2. Typical Chaining Table

The chains present in Figure 11.2 are exactly (11.3)-(11.10) given below. Note that no chain begins with a pronoun. (Here, we are staying with the convention of Chapter 10 where a "^" symbol indicates a chain_link).

(11.3) Johna (11.4) John_b^{hea}

- (11.5) John_c ϕ_a
- (11.6) $John_d \phi_d hea$

(11.7) June_a

 $\begin{array}{c} (11.8) \ \mathrm{June_b}\ \hat{s}hea \\ (11.9) \ \mathrm{present_a} \\ (11.10) \ \mathrm{present_b}\ \hat{i}ta \end{array}$

The only interpretation derivable from Figure 11.2 is (11.11).

(11.11) $John_d \phi_d he_a June_b she_a present_h it_a$

Exactly how chaining information from a table to system parse tree will have to be system dependent, but we can imagine that noun phrases in a system parse tree are some kind of list elements which have linked to them, among other things, lists corresponding to their semantics. It is up to the Table Interpreter to set any chain_link's inside the semantics of the noun phrases of the system parse tree. The Semantic Processor module of the system should then be powerful enough to be able to handle the kind of coordination that chain_link's imply.

This strategy has a number of possibilities that simple methods of **coreferencing** are just not able to handle. Consider sentence (11.12) for example.

(11.12) Jack's house burned down, but he rebuilt it.

We can't really say that it **corefers** with <u>Jack's house</u> as <u>Jack's house</u> is some object that existed in the past and has stopped existing while <u>it</u> refers to some new object. This does not mean that <u>it</u> cannot chain from <u>Jack's house</u>, however, and indeed it should. The information the Semantic Processor module needs to give meaning to <u>it</u> is contained in <u>Jack's house</u>, and so there must be a chain_link from <u>Jack's house</u> to <u>it</u> in order for the Semantic Processor to give meaning to <u>it</u>.

A similar result holds for quantifiers. We see that (11.13) is not equivalent to (11.14).

(11.13) Every connoisseur loves his wine and cheese. \neq

(11.14) Every connoisseur loves every connoisseur's wine and cheese.

Quite clearly, <u>his</u> cannot be replaced by <u>every connoisseur</u> and preserve the meaning of the sentence. Instead, (11.13) has more the meaning given by (11.15).

(11.15) (For all x: x is a connoisseur)(x loves x's wine and cheese.)

A number of other examples are pointed out by Bresnan [3].

(11.16) All Italians think they are handsome. \neq

(11.17) All Italians think all Italians are handsome.

(11.18) Every Italian thinks he is handsome. \neq

(11.19) Every Italian thinks every Italian is handsome.

(11.20) Any Italian would die for his mother. \neq

(11.21) Any Italian would die for and Italian's mother.

(11.22) Every Italian thinks that he alone is handsome. \neq

(11.23) *Every Italian thinks that every Italian alone is handsome.

(11.24) One girl claimed that she herself could read Homer. \neq

(11.25) *One girl claimed that one girl herself could read Homer.

It appears that the proper interpretation for a pronoun chained to quantified noun phrase within the **scope of quantification** is for the pronoun to act as a **bound variable**.

When the pronoun is outside the scope of quantification, it is a different story. Consider (11.26) and (11.27) from Evans [6].

(11.26) John owns some sheep and Harry vaccinates them.

(11.27) Mary danced with many boys and they found her interesting.

This time the pronouns are chaining to quantified noun phrases, but do not themselves lie within the scope of quantification. Instead, they appear to refer to the **range of the quantification**.

Similar results hold for (11.28)-(11.31) from Sidner [32].

- (11.28) John lost a pen yesterday and Bill found one today.
- (11.29) John claimed to have found the solution to the problem, but Bill was sure he had found it.
- (11.30) John wants to catch a fish and eat it for supper.
- (11.31) No one would put the blame on himself.

The problems mentioned above are all rather tricky, but viewing them from the vantage point of chaining sheds more light on them than viewing them through some kind of coreference. The moral of the story seems to be that anaphora is not coreference.

Using the Table Interpreter now, we present some more examples.

(11.32) Sue told Sandy about herself.

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| Sue | - | - | - | + | _ | + | + | - |
| Sandy | - | - | - | + | - | + | + | - |
| herself | + | _ | _ | + | _ | + | + | + |

| Chaining | | | | | |
|-----------|--------|----------|--|--|--|
| Sue Sandy | | herself | | | |
| Suea | Sandya | herselfa | | | |

```
init_table: exiting
chaining_n(herself)
    refl_chaining(herself)
    simplex_pred(herself)
    simplex_pred: Sandy
    chaining_n_to_n(herself, Sandy)
        sc(herself, Sandy) = True
        agr(herself, Sandy) = True
        rnr(herself, Sandy) = True
        chaining_e_to_n(herself_a, Sandy)
        agr(herself_a, Sandy) = True
        new_chain(herself_a, Sandy)
        new_chain: create Sandyb
        new_chain: create Sandyb
        new_chain: create Sandyb
        new_chain: create Sandyb
        new_chain: create Sandyb
```

```
ChainingSueSandyherselfSueSandyaherselfaSandyb^herselfaherselfa
```

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
simplex_pred(Sandy)
simplex_pred: Sue
chaining_n_to_n(herself, Sue)
sc(herself, Sue) = True
agr(herself, Sue) = True
rnr(herself, Sue) = True
chaining_e_to_n(herself_a, Sue)
agr(herself_a, Sue) = True
new_chain(herself_a, Sue)
new_chain: create Sueb
new_chain: create Sueb^herself_a
```

| Chaining | | | | | |
|----------------------------|------------------------------|----------|--|--|--|
| Sue | Sandy | herself | | | |
| Suea | Sandya | herselfa | | | |
| Sue _b ^herselfa | Sandy _b ^herselfa | | | | |

| Chaining | | | | | |
|----------------------------|------------------------------|----------|--|--|--|
| Sue | Sandy | herself | | | |
| Suea | Sandya | herselfa | | | |
| Sue _b ^herselfa | Sandy _b ^herselfa | | | | |

| Interpretations |
|--|
| Sue _b ^herselfa |
| Sandy _b ^herself _a |

(11.33) *Jill kept talking about himself.

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| Jill | - | - | - | + | - | + | + | - |
| himself | + | - | - | + | - | - | + | + |

| Chaining | | | | |
|----------|----------|--|--|--|
| Jill | himself | | | |
| Jilla | himselfa | | | |

```
init_table: exiting
chaining_n(himself)
    refl_chaining(himself)
    simplex_pred(himself)
    simplex_pred: Jill
    chaining_n_to_n(himself, Jill)
        sc(himself, Jill) = True
        agr(himself, Jill) = False
        chaining_n_to_n: exiting
        simplex_pred(Jill)
        simplex_pred:
        refl_chaining: exiting
        chaining_n: exiting
        chaining: exiting
        chaining: exiting
```

| Chaining | | | | |
|----------|----------|--|--|--|
| Jill | himself | | | |
| Jilla | himselfa | | | |

| Interpretations | |
|-----------------|--|
| NONE | |

(11.34) Does Jack's making a pig of himself bother Bill?

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| Jack' s | - | _ | - | + | _ | _ | + | _ |
| pig | - | - | - | + | - | ? | + | - |
| himself | + | - | - | + | - | - | + | + |
| Bill | - | - | _ | + | - | _ | + | _ |

| Chaining | | | | | | |
|------------|------|----------|-------|--|--|--|
| Jack's pig | | himself | Bill | | | |
| Jack'sa | piga | himselfa | Billa | | | |

```
init_table: exiting
chaining_n(himself)
    refl_chaining(himself)
    simplex_pred(himself)
    simplex_pred: pig
    chaining_n_to_n(himself, pig)
        sc(himself, pig) = True
        agr(himself, pig) = True
        rnr(himself, pig) = True
        chaining_e_to_n(himself_a, pig)
        agr(himself_a, pig) = True
        new_chain(himself_a, pig)
        new_chain: create pigb
        new_chain: create pigb^himself_a
```

| Chaining | | | | | | | |
|----------|--|----------|-------|--|--|--|--|
| Jack's | pig | himself | Bill | | | | |
| Jack'sa | pig _a pig _b ^himselfa | himselfa | Billa | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
simplex_pred(pig)
simplex_pred: Jack's
chaining_n_to_n(himself, Jack's)
sc(himself, Jack's) = True
agr(himself, Jack's) = True
rnr(himself, Jack's) = False
chaining_n_to_n: exiting
simplex_pred(Jack's)
simplex_pred:
refl_chaining: exiting
chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | | |
|----------|--|----------|-------|--|--|--|--|
| Jack's | pig | himself | Bill | | | | |
| Jack'sa | pig _a pig _b ^himself _a | himselfa | Billa | | | | |

| Interpretations | | | | |
|----------------------------|--|--|--|--|
| pig _b ^himselfa | | | | |

| Features | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| John | - | - | - | + | - | - | + | - |
| ϕ | + | ? | ? | ? | ? | ? | ? | - |
| June | - | _ | - | + | - | + | + | - |
| present | - | _ | - | + | - | ? | - | - |
| he | + | - | - | + | - | _ | + | - |
| she | + | - | - | + | - | + | + | - |
| it | + | - | - | + | - | ? | - | - |

 $(11.35)\,$ John wants to give June a present, but he is a fraid she won't like it.

```
chaining
init_table
```

| Chaining | | | | | | | | |
|----------|------------|-------|----------|-----|------|-----|--|--|
| John | ϕ | June | present | he | she | it | | |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita | | |

```
init_table: exiting
chaining_n(it)
    non_refl_chaining(it)
        chaining_n_to_n(it, she)
            sc(it, she) = True
            agr(it, she) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, he)
            sc(it, he) = True
            agr(it, he) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, present)
            sc(it, present) = True
            agr(it, present) = True
            rnr(it, present) = True
            chaining_e_to_n(ita, present)
                 agr(it_a, present) = True
                 new_chain(ita, present)
                     new_chain: create present_b
                     new_chain: create present<sub>b</sub>^ita
```

| Chaining | | | | | | | | |
|----------|------------|-------|---------------------------|-----|------|-----|--|--|
| John | ϕ | June | present | he | she | it | | |
| Johna | ϕ_{a} | Junea | present _a | hea | shea | ita | | |
| | | | present _b ^ita | | | | | |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(it, June)

sc(it, June) = True

agr(it, June) = False

chaining_n_to_n: exiting

chaining_n_to_n(it, \phi)

sc(it, \phi) = True

agr(it, \phi) = True

rnr(it, \phi) = True

chaining_e_to_n(it_a, \phi)

agr(it_a, \phi) = True

new_chain(it_a, \phi)

new_chain: create \phi_b

new_chain: create \phi_b^{-}ita
```

| Chaining | | | | | | | | |
|----------|---------------------|-------|---------------------------|-----|------|-----|--|--|
| John | ϕ | June | present | he | she | it | | |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita | | |
| | $\phi_{\rm b}$ ^ita | | present _b ^ita | | | | | |

```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(it, John)
            sc(it, John) = True
            agr(it, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(she)
    non_refl_chaining(she)
        chaining_n_to_n(she, it)
            sc(she, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, he)
            sc(she, he) = True
            agr(she, he) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, present)
            sc(she, present) = True
            agr(she, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, June)
            sc(she, June) = True
            agr(she, June) = True
            rnr(she, June) = True
            chaining_e_to_n(shea, June)
                agr(shea, June) = True
                new_chain(shea, June)
                    new_chain: create Juneb
                    new_chain: create June<sub>b</sub>^shea
```

| Chaining | | | | | | | | | |
|----------|---------------------|-------------------------|---------------------------|-----|------|-----|--|--|--|
| John | ϕ | June | present | he | she | it | | | |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita | | | |
| | $\phi_{\rm b}$ ^ita | June _b ^shea | present _b ^ita | | | | | | |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(she, \phi)

sc(she, \phi) = True

agr(she, \phi) = True

rnr(she, \phi) = True

chaining_e_to_n(shea, \phi)

agr(shea, \phi) = True

new_chain(shea, \phi)

new_chain: create \phi_c

new_chain: create \phi_c ^ shea
```

| Chaining | | | | | | | | |
|-------------------|--------------------------------------|--|---|-----|------|-----|--|--|
| John | ϕ | June | present | he | she | it | | |
| John _a | $\phi_{ m b}$ îta $\phi_{ m c}$ shea | June _a June _b ^she _a | present _a present _b ^ita | hea | shea | ita | | |
```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(she, John)
            sc(she, John) = True
            agr(she, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(he)
    non_refl_chaining(he)
        chaining_n_to_n(he, it)
            sc(he, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, she)
            sc(he, she) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, present)
            sc(he, present) = True
            agr(he, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, June)
            sc(he, June) = True
            agr(he, June) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, \phi)
            sc(he, \phi) = True
            agr(he, \phi) = True
            rnr(he, \phi) = True
            chaining_e_to_n(hea, \phi)
                 agr(hea, \phi) = True
                 new_chain(hea, \phi)
                     new_chain: create \phi_{\rm d}
                     new_chain: create \phi_d hea
```

| | | \mathbf{Ch} | aining | | | |
|-------------------|--|--|---|-----|------|-----|
| John | ϕ | June | present | he | she | it |
| John _a | ϕ_a $\phi_b^it_a$ $\phi_c^she_a$ $\phi_d^he_a$ | June _a June _b ^she _a | present _a present _b ^ita | hea | shea | ita |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(he, John)
sc(he, John) = True
agr(he, John) = True
rnr(he, John) = True
chaining_e_to_n(hea, John)
agr(hea, John) = True
new_chain(hea, John)
new_chain: create John<sub>b</sub>
new_chain: create John<sub>b</sub>^hea
```

| | | Chair | ning | | | |
|------------------------|---|-------------------------|---------------------------|-----|------|-----|
| John | ϕ | June | present | he | she | it |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita |
| John _b ^hea | $\phi_{\rm b}$ îta $\phi_{\rm c}$ shea $\phi_{\rm d}$ hea | June _b ^shea | present _b ^ita | | | |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)
    non refl chaining (\phi)
         chaining_n_to_n(\phi, it)
             sc(\phi, it) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, she)
             sc(\phi, she) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, he)
             sc(\phi, he) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, present)
             sc(\phi, present) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, June)
             sc(\phi, June) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = True
             rnr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create Johnc
                      new_chain: create John<sub>C</sub>^{\phi}a
```

| | | Chair | ing | | | |
|---|--|--|---|-----|------|-----|
| John | ϕ | June | present | he | she | it |
| John _a John _b ^he _a John _C ^ ϕ_a | ϕ_a $\phi_b^it_a$ $\phi_c^she_a$ $\phi_d^he_a$ | June _a June _b ^she _a | present _a present _b ^ita | hea | shea | ita |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_e_to_n(\phi_b, John)

agr(\phi_b, John) = False

chaining_e_to_n: exiting

chaining_e_to_n: exiting

chaining_e_to_n: exiting

chaining_e_to_n: exiting

chaining_e_to_n(\phi_d, John)

agr(\phi_d, John) = True

new_chain(\phi_d, John)

new_chain: create Johnd

new_chain: create Johnd
```

| | | Chain | ning | | | |
|---------------------------|---------------------------------|-------------------------|---------------------------|-----|------|-----|
| John | ϕ | June | present | he | she | it |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita |
| John _b ^hea | $\phi_{ m b}$ îta | June _b ^shea | present _b ^ita | | | |
| $John_{C}^{\phi}\phi_{A}$ | $\phi_{\rm C}$ $^{\rm she_{a}}$ | | | | | |
| $John_d^{\phi}d$ | $\phi_{\rm d}$ hea | | | | | |

| | | Chair | ning | | | |
|---------------------------|---------------------|-------------------------|---------------------------|-----|------|-----|
| John | ϕ | June | present | he | she | it |
| Johna | ϕ_{a} | Junea | presenta | hea | shea | ita |
| John _b ^hea | $\phi_{ m b}$ îta | June _b ^shea | present _b ^ita | | | |
| $John_{C}^{\phi}\phi_{A}$ | $\phi_{\rm C}$ shea | | | | | |
| $John_d^{\phi}d$ | $\phi_{\rm d}$ hea | | | | | |

| Γ | | Interpretations | |
|---|---|-------------------------------------|---------------------------|
| Γ | ${\rm John}_{\rm d} \hat{\phi}_{\rm d} \hat{h} {\rm e}_{\rm d}$ | June _b ^she _a | present _b ^ita |

(11.36) Ernie doesn't like Bernie, because he is such an asshole.

| | | | Fea | \mathbf{tures} | | | | |
|---------|-----|-----|-----|------------------|-----|-----|-----|-----|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF |
| Ernie | - | _ | - | + | _ | - | + | _ |
| Bernie | - | _ | - | + | _ | - | + | _ |
| he | + | - | - | + | - | - | + | - |
| asshole | - | - | - | + | - | ? | + | - |

chaining init_table

| Chaining | | | | | |
|----------|---------|-----|----------|--|--|
| Ernie | Bernie | he | asshole | | |
| Erniea | Berniea | hea | assholea | | |

```
init_table: exiting
chaining_n(he)
    non_refl_chaining(he)
    chaining_n_to_n(he, asshole)
        sc(he, asshole) = False
    chaining_n_to_n: exiting
    chaining_n_to_n(he, Bernie)
        sc(he, Bernie) = True
        agr(he, Bernie) = True
        rnr(he, Bernie) = True
        chaining_e_to_n(he_a, Bernie)
        agr(he_a, Bernie) = True
        new_chain(he_a, Bernie)
        new_chain: create Bernieb
        new_chain: create Bernieb^he_a
```

| | Chainin | g | |
|--------|--------------------------|-----|----------|
| Ernie | Bernie | he | asshole |
| Erniea | Berniea | hea | assholea |
| | Bernie _b ^hea | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n (he, Ernie)
sc(he, Ernie) = True
agr(he, Ernie) = True
rnr(he, Ernie) = True
chaining_e_to_n (hea, Ernie)
agr(hea, Ernie) = True
new_chain(hea, Ernie)
new_chain: create Ernieb
new_chain: create Ernieb^hea
```

| | Chaining | | |
|-------------------------|--------------------------|-----|----------|
| Ernie | Bernie | he | asshole |
| Erniea | Berniea | hea | assholea |
| Ernie _b ^hea | Bernie _b ^hea | | |

| | Chaining | | |
|-------------------------|--------------------------|-----|----------|
| Ernie | Bernie | he | asshole |
| Erniea | Berniea | hea | assholea |
| Ernie _b ^hea | Bernie _b ^hea | | |

| Interpretations |
|--------------------------|
| Ernie _b ^hea |
| Bernie _b ^hea |

12 Genitives

Very little modification to what has been said so far is necessary to implement attributive possessive pronouns. Recall that the attributive possessive pronouns are those pronouns listed in (12.1). (12.1) my, our, your, her, his, its, their

Examining sentences like (12.2)-(12.5) reveals that reflexive pronouns don't chain to genitives within the same simplex. On the other hand, nonreflexive pronouns can.

- (12.2) Mary's father killed himself.
- (12.3) *Mary's father killed him.
- (12.4) *Mary's father killed herself.
- (12.5) Mary's father killed her.

The same conclusions also hold for **of-genitives**. Compare sentences (12.6)-(12.9) to (12.2)-(12.5).

- (12.6) The father of Mary killed himself.
- (12.7) *The father of Mary killed him.
- (12.8) *The father of Mary killed herself.
- (12.9) The father of Mary killed her.

The easiest way to handle genitives, apparently, is to introduce a new Feature, GEN, for genitive and to modify the Reflexive Nonreflexive Rule to handle genitives. The new form of the Reflexive Nonreflexive Rule is shown below in Figure 12.10.

```
function rnr(n1,n2:NodePointer):boolean;
    {Reflexive Nonreflexive Rule}
    var ftr1, ftr2:features;
begin
    n1:=n1^.np_link;
    n2:=n2^.np_link;
    ftr1:=n1^.ftr;
    ftr2:=n2^.ftr:
    if ftr2[GEN]=PLUS then rnr:=false
    else case ftr1[RPF] of
                  rnr:=(n1^.up_link=n2^.up_link)
        PLUS:
                        and (ftr1[GEN]==MINUS);
        QUESTION: {doesn't occur};
        MINUS: rnr:=(n1^.up_link<>n2^.up_link)
                        or (ftr1[GEN]<>MINUS);
    end;
end;
```

Figure 12.10. New Reflexive Nonreflexive Rule

The examples following illustrate the interpretation of attributive possessive pronouns and pronouns in the context of genitives.

(12.11) Mary's mother cooks only for herself.

| Features | | | | | | | | | | |
|----------|-----|-----|-----|-----|-----|-----|-----|-----|-----|--|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF | GEN | |
| Mary's | - | - | - | + | - | + | + | - | + | |
| mother | - | - | _ | + | - | + | + | - | _ | |
| herself | + | - | _ | + | - | + | + | + | - | |

(12.12) Mary's mother cooks only for her.

| Features | | | | | | | | | | |
|-------------------------------------|---|---|---|---|---|---|---|---|---|--|
| PNF FPF SPF TPF PLF GNF ANF RPF GEN | | | | | | | | | | |
| Mary's | - | - | _ | + | _ | + | + | - | + | |
| mother | - | - | _ | + | _ | + | + | - | _ | |
| her | + | - | _ | + | _ | + | + | - | ? | |

chaining init_table

| Chaining | | | | | |
|----------|---------|------|--|--|--|
| Mary's | mother | her | | | |
| Mary'sa | mothera | hera | | | |

```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
    chaining_n_to_n(her, mother)
        sc(her, mother) = True
        agr(her, mother) = True
        rnr(her, mother) = True
        chaining_e_to_n(hera, mother)
        agr(hera, mother) = True
        new_chain(hera, mother)
        new_chain: create motherb
        new_chain: create motherb^hera
```

| Chaining | | | | | | | |
|---------------------|---------------------------|------|--|--|--|--|--|
| Mary's | mother | her | | | | | |
| Mary's _a | mothera | hera | | | | | |
| | mother _b ^hera | | | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(her, Mary's)
sc(her, Mary's) = True
agr(her, Mary's) = True
rnr(her, Mary's) = False
chaining_n_to_n: exiting
non_refl_chaining: exiting
chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | |
|----------|---------------------------|------|--|--|--|--|
| Mary's | her | | | | | |
| Mary'sa | mothera | hera | | | | |
| | mother _b ^hera | | | | | |

| Interpretations | |
|---------------------------|--|
| mother _b ^hera | |

(12.13) Mary's mother cooks only for her mother.

| Features | | | | | | | | | | |
|---------------------|-----|-----|-----|-----|-----|-----|-----|-----|-----|--|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF | GEN | |
| Mary's | - | - | - | + | - | + | + | - | + | |
| mother1 | _ | - | - | + | - | + | + | - | - | |
| her | + | - | - | + | - | + | + | - | ? | |
| mother ₂ | - | - | - | + | - | + | + | - | - | |

chaining init_table

| Chaining | | | | | | | |
|---------------------|----------------------|------|----------------------|--|--|--|--|
| Mary's | mother1 | her | mother ₂ | | | | |
| Mary's _a | mother _{1a} | hera | mother _{2a} | | | | |

```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
    chaining_n_to_n(her, mother_2)
        sc(her, mother_2) = False
    chaining_n_to_n(her, mother_1)
        sc(her, mother_1) = True
        agr(her, mother_1) = True
        rnr(her, mother_1) = True
        rnr(her, mother_1) = True
        chaining_e_to_n(her_a, mother_1)
        agr(her_a, mother_1) = True
        new_chain(her_a, mother_1)
        new_chain: create mother_1b
        new_chain: create mother_1b^her_a
```

| Chaining | | | | | | | | |
|---------------------|--|------|----------------------|--|--|--|--|--|
| Mary's | mother ₁ | her | mother ₂ | | | | | |
| Mary's _a | mother _{1a} mother _{1b} ^hera | hera | mother _{2a} | | | | | |

```
new_chain: exiting
chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(her, Mary's)
sc(her, Mary's) = True
agr(her, Mary's) = True
rnr(her, Mary's) = False
chaining_n_to_n: exiting
non_refl_chaining: exiting
chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | | | |
|----------|----------------------------|------|----------------------|--|--|--|--|--|
| Mary's | mother ₁ | her | mother ₂ | | | | | |
| Mary'sa | mother _{la} | hera | mother _{2a} | | | | | |
| | mother _{lb} ^hera | | | | | | | |

| Interpretations |
|----------------------------|
| mother _{lb} ^hera |

13 Focusing

Extrasentential anaphora and **ellipsis** is possible through the maintenance of a **focus** of conversation. This maintenance is known as **focusing** and has been described at length by Grosz [10] and Sidner [32]. By focus of conversation, we mean the common view of the participants of conversation of what their conversation is about. Focusing is useful because it allows the participants of conversation to avoid redundant repetition of old material. Assuming focusing is desirable in a computer natural language system, how do we implement it?

Grosz has examined task dialogues in which an expert helps an apprentice to assemble a mechanical air compressor. She finds it convenient to represent the focus of conversation as a set of overlapping **focus spaces**, where each focus space is a collection of objects. One focus space is active and the others are open. When a focus space is no longer needed, it is closed. One of Grosz's assumptions is that goals and subgoals are definable and recognizable in a task dialogue system with the consequence that in any conversation there is an open focus space hierarchy with the active focus space at the bottom of the hierarchy.

Sidner has approached the problem of focusing from a different perspective by analyzing monologues. For Sidner, focus is kept track of through a **discourse focus**, **actor focus**, **potential discourse foci**, **potential actor foci**, **discourse focus stack**, and **actor focus stack**. Sidner's work, which came after Grosz's, is very commendable for the algorithms she presents, although most of these are fairly sketchy.

In our approach, we will treat the focus of conversation as a collection of nonpronominal N-nodes. Among the N-nodes that we would ordinairly expect to always be in focus are the <u>I</u> and <u>you</u> of a conversation. To get a handle on the focused N-nodes, we dominate them by an S-node just as if they all had occurred in one simplex. So, for example, if I₀ and <u>you</u>₀ are the nonpronominal N-nodes currently in focus, then the current focus representation is given by a structure like Figure 13.1.



Figure 13.1. Typical Focus Representation

When it comes time to analyze a sentence, the current focus representation is attached to the C-S-N parse tree of the sentence via a C-node which *dominates* them both. This makes the focused N-nodes available to the N-nodes of the C-S-N parse tree for chaining.

As an example, suppose $\underline{I_0}$ and $\underline{you_0}$ are in focus and the current input sentence is (13.2) from Grinder [8].

(13.2) It was difficult to sketch myself.



The C-S-N parse tree of (13.2) will have a form like that indicated below in Figure 13.3.



Figure 13.3. C-S-N Parse Tree of (13.2)

As the parse tree now stands in Figure 2, <u>myself</u> may chain to ϕ , but ϕ does not have an N-node to chain to. Thus, there are no legitimate interpretations without focusing. With focusing, the Parser attaches the current focus representation containing <u>I₀</u> and <u>you</u>₀ to the C-S-N parse tree by a C-node obtaining the new C-S-N parse tree shown in Figure 13.4.



Figure 13.4. C-S-N Parse Tree with Focusing

Now, <u>myself</u> can chain from $\underline{\phi}$ and $\underline{\phi}$ can chain from $\underline{I_0}$ giving us a legitimate interpretation of the C-S-N tree.

This kind of strategy explains a number of other examples from Grinder. Grinder lists (13.5a)-(13.11a) as grammatical.

(13.5a) It was difficult for me to sketch myself.

- (13.6a) It was difficult for you to sketch yourself.
- (13.7a) It was difficult for him to sketch himself.
- (13.8a) It was difficult for her to sketch herself.
- (13.9a) It was difficult for us to sketch ourselves.
- (13.10a) It was difficult for you to sketch yourselves.
- (13.11a) It was difficult for them to sketch themselves.

After Equi-NP Deletion, Grinder lists only (13.5b), (13.6b), (13.9b), and (13.10b) as grammatical.

- (13.5b) It was difficult to sketch myself.
- (13.6b) It was difficult to sketch yourself.
- (13.7b) *It was difficult to sketch himself.
- (13.8b) *It was difficult to sketch herself.
- (13.9b) It was difficult to sketch ourselves.
- (13.10b) It was difficult to sketch yourselves.
- (13.11b) *It was difficult to sketch themselves.

The probable reason this comes about is that we are used to thinking of \underline{I} , you singular, us, and you plural as always being in focus, while referents for <u>he</u>, <u>she</u>, and <u>they</u> are ordinairly not in focus. Needless to say, if referents for <u>he</u>, <u>she</u>, or <u>they</u> are in focus, the situation changes completely. This is shown by (13.12)-(13.14).

- (13.12a) Nurse Bob Breezy gave up drawing.
- (13.12b) [Bob] It was difficult to sketch himself.
- (13.13a) Astronaut Linda Smith gave up drawing.
- (13.13b) [Linda] It was difficult to sketch herself.
- (13.14a) The bank embezzlers gave up drawing.
- (13.14b) [the bank embezzlers] It was difficult to sketch themselves.

To indicate that various N-nodes are in focus, we bracket them at the beginning of a sentence. Thus (13.15)-(13.17) are not interpretable while (13.18)-(13.20) are.

- (13.15) *It was difficult to sketch himself.
- (13.16) *It was difficult to sketch herself.
- (13.17) *It was difficult to sketch themselves.
- (13.18) [Bob] It was difficult to sketch himself.
- (13.19) [Linda] It was difficult to sketch herself.
- (13.20) [the bank embezzlers] It was difficult to sketch themselves.

The following examples involve resolution through focusing.

(13.21) It was difficult to sketch myself.

| Features | | | | | | | | | | |
|------------------|-----|-----|-----|-----|-----|-----|-----|-----|-----|--|
| | PNF | FPF | SPF | TPF | PLF | GNF | ANF | RPF | GEN | |
| IO | + | + | - | - | - | ? | + | - | - | |
| you ₀ | + | - | + | - | ? | ? | + | - | - | |
| ϕ | + | ? | ? | ? | ? | ? | ? | - | - | |
| myself | + | + | - | - | - | ? | + | + | - | |

```
chaining
init_table
```

| Chaining | | | | | | | |
|-----------------|-------------------|------------|---------|--|--|--|--|
| I ₀ | you0 | ϕ | myself | | | | |
| I _{0a} | you _{0a} | ϕ_{a} | myselfa | | | | |

```
init_table: exiting
chaining_n(myself)
    refl_chaining(myself)
    simplex_pred(myself)
    simplex_pred: \phi
    chaining_n_to_n(myself, \phi)
        sc(myself, \phi) = True
        agr(myself, \phi) = True
        rnr(myself, \phi) = True
        chaining_e_to_n(myself_a, \phi)
        agr(myself_a, \phi) = True
        new_chain(myself_a, \phi)
        new_chain: create \phi_b
        new_chain: create \phi_b^{-}myselfa
```

| Chaining | | | | | | |
|-------------------|-----------------------|---|--|--|--|--|
| you ₀ | ϕ | myself | | | | |
| you _{0a} | ϕ_a | myselfa | | | | |
| - | you 0 you0a | Chainingyou ϕ you ϕ_a ϕ_b myself | | | | |

```
new_chain: exiting
              chaining_e_to_n: exiting
          chaining_n_to_n: exiting
          simplex_pred(\phi)
          simplex_pred:
     refl_chaining: exiting
chaining n: exiting
chaining_n(\phi)
    non_refl_chaining(\phi)
          chaining_n_to_n(\phi, myself)
               sc(\phi, myself) = False
          chaining_n_to_n: exiting
          chaining_n_to_n(\phi, you<sub>0</sub>)
               sc(\phi, you_0) = True
              agr(\phi, you_0) = True
              rnr(\phi, you_0) = True
              chaining_e_to_n(\phi_a, you<sub>0</sub>)
                   agr(\phi_a, you_0) = True
                   new_chain(\phi_a, you<sub>0</sub>)
                        new_chain: create you<sub>0b</sub>
                         new_chain: create you<sub>0b</sub>\phi_a
```

| | Chaining | | | | | | | | |
|-----------------|--------------------------------|-------------------------|---------|--|--|--|--|--|--|
| IO | you ₀ | ϕ | myself | | | | | | |
| I _{0a} | you _{0a} | ϕ_{a} | myselfa | | | | | | |
| | you _{0b} ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | | | | |

```
\begin{array}{c} \text{new\_chain: exiting}\\ \text{chaining\_e\_to\_n: exiting}\\ \text{chaining\_e\_to\_n: exiting}\\ \text{chaining\_e\_to\_n: exiting}\\ \text{chaining\_n\_to\_n: exiting}\\ \text{chaining\_n\_to\_n: exiting}\\ \text{chaining\_n\_to\_n(\phi, I_0)}\\ \text{sc}(\phi, I_0) = \text{True}\\ \text{agr}(\phi, I_0) = \text{True}\\ \text{rnr}(\phi, I_0) = \text{True}\\ \text{chaining\_e\_to\_n(\phi_a, I_0)}\\ \text{agr}(\phi_a, I_0) = \text{True}\\ \text{new\_chain: create } I_{0b}\\ \text{new\_chain: create } I_{0b}^{-}\phi_a \end{array}
```

| Chaining | | | | | | |
|------------------|--------------------------------|-------------------------|---------|--|--|--|
| IO | you0 | ϕ | myself | | | |
| I _{0a} | you _{0a} | ϕ_{a} | myselfa | | | |
| $I_{0b}^{\phi}a$ | you _{Ob} ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | |

new_chain: exiting chaining_e_to_n: exiting chaining_e_to_n(ϕ_b , I_0) agr(ϕ_b , I_0) = True new_chain(ϕ_b , I_0) new_chain: create I_{0c} new_chain: create $I_{0c}^{-}\phi_b$

| Chaining | | | | | | | | |
|------------------|--------------------------|-------------------------|---------|--|--|--|--|--|
| IO | you0 | ϕ | myself | | | | | |
| I _{0a} | you _{0a} | ϕ_{a} | myselfa | | | | | |
| $I_{0b}^{\phi}a$ | you $_{0b}$ ^ ϕ_{a} | $\phi_{\rm b}$ ^myselfa | | | | | | |
| $I_{0c} \phi_b$ | | | | | | | | |

```
new_chain: exiting
                  chaining_e_to_n: exiting
              chaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(you∩)
         non_refl_chaining(you<sub>0</sub>)
              chaining_n_to_n(you<sub>0</sub>, myself)
                  sc(you_0, myself) = False
              chaining_n_to_n: exiting
              chaining_n_to_n(you_0, \phi)
                  sc(you_0, \phi) = False
              chaining_n_to_n: exiting
              chaining_n_to_n(you<sub>0</sub>, I<sub>0</sub>)
                   sc(you_0, I_0) = True
                  agr(you_0, I_0) = False
              chaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I_0)
         non_refl_chaining(I<sub>0</sub>)
              chaining_n_to_n(I<sub>0</sub>, myself)
                   sc(I_0, myself) = False
              chaining_n_to_n: exiting
              chaining_n_to_n(I<sub>0</sub>, \phi)
                  sc(I_0, \phi) = False
              chaining_n_to_n: exiting
              chaining_n_to_n(I<sub>0</sub>, you<sub>0</sub>)
                   sc(I_0, you_0) = False
              chaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | | | | |
|------------------|-------------------------------|-------------------------|---------------------|--|--|--|--|--|
| IO | you0 | ϕ | myself | | | | | |
| I _{0a} | you _{0a} | ϕ_{a} | myself _a | | | | | |
| $I_{0b}^{\phi}a$ | you $_{0b}{}^{\circ}\phi_{a}$ | $\phi_{\rm b}$ ^myselfa | | | | | | |
| $I_{0c}^{\phi}b$ | | | | | | | | |

| Interpretations | | | | | |
|---------------------------|--|--|--|--|--|
| $I_{0c}^{\phi_b}$ myselfa | | | | | |

(13.22) [toy] Give me that!

| | Features | | | | | | | | | | |
|--------|-------------------------------------|---|---|---|---|---|---|---|---|--|--|
| | PNF FPF SPF TPF PLF GNF ANF RPF GEN | | | | | | | | | | |
| IO | + | + | - | _ | _ | ? | + | - | - | | |
| you0 | + | _ | + | _ | ? | ? | + | - | - | | |
| toy | - | - | - | + | - | ? | ? | - | - | | |
| ϕ | + | ? | ? | ? | ? | ? | ? | - | - | | |
| me | + | + | - | - | - | ? | + | - | - | | |
| that | + | - | - | + | - | ? | ? | - | - | | |

```
chaining
init_table
```

| Chaining | | | | | | | |
|-----------------|-------------------|------|------------|-----|-------|--|--|
| IO | you ₀ | toy | ϕ | me | that | | |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata | | |

```
init_table: exiting
chaining_n(that)
    non_refl_chaining(that)
        chaining_n_to_n(that, me)
            sc(that, me) = True
            agr(that, me) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(that, \phi)
            sc(that, \phi) = True
            agr(that, \phi) = True
            rnr(that, \phi) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(that, toy)
            sc(that, toy) = True
            agr(that, toy) = True
            rnr(that, toy) = True
            chaining_e_to_n(thata, toy)
                agr(that_a, toy) = True
                new_chain(thata, toy)
                     new_chain: create toyb
                     new_chain: create toy_b^thata
```

| Chaining | | | | | | | |
|-----------------|-------------------|-------------------------|------------|-----|-------|--|--|
| I ₀ | you ₀ | toy | ϕ | me | that | | |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata | | |
| | | toy _b "thata | | | | | |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         chaining_n_to_n(that, you_0)
             sc(that, you_0) = True
             agr(that, you_0) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(that, I_0)
             sc(that, I_0) = True
             agr(that, I_0) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(me)
    non_refl_chaining(me)
        chaining_n_to_n(me, that)
             sc(me, that) = False
        chaining_n_to_n: exiting
         chaining_n_to_n(me, \phi)
             sc(me, \phi) = True
             agr(me, \phi) = True
             rnr(me, \phi) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(me, toy)
             sc(me, toy) = True
             agr(me, toy) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(me, you<sub>0</sub>)
             sc(me, you_0) = True
             agr(me, you_0) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(me, I<sub>0</sub>)
             sc(me, I_0) = True
             agr(me, I_0) = True
             rnr(me, I_0) = True
             chaining_e_to_n(mea, I_0)
                 agr(me_a, I_0) = True
                 new_chain(mea, I_0)
                      new_chain: create I<sub>0b</sub>
                      new_chain: create I<sub>Ob</sub>^mea
```

| Chaining | | | | | | | |
|----------------------|-------------------|-------------------------|------------|-----|-------|--|--|
| IO | you ₀ | toy | ϕ | me | that | | |
| IOa | you _{0a} | toya | ϕ_{a} | mea | thata | | |
| I _{Ob} ^mea | | toy _b ^thata | | | | | |

```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)
    non_refl_chaining(\phi)
         chaining_n_to_n(\phi, that)
             sc(\phi, that) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, me)
             sc(\phi, me) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, toy)
             sc(\phi, toy) = True
             agr(\phi, toy) = True
             rnr(\phi, toy) = True
             chaining_e_to_n(\phi_a, toy)
                  agr(\phi_a, toy) = True
                  new_chain(\phi_a, toy)
                      new_chain: create toyc
                      new_chain: create toy_c^\phi_{\rm A}
```

| Chaining | | | | | | | |
|----------------------|-------------------|-------------------------|------------|-----|-------|--|--|
| IO | you ₀ | toy | ϕ | me | that | | |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata | | |
| I _{Ob} ^mea | | toy _b ^thata | | | | | |
| | | toyc ϕ_a | | | | | |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(\phi, you_0)

sc(\phi, you_0) = True

agr(\phi, you_0) = True

rnr(\phi, you_0) = True

chaining_e_to_n(\phi_a, you_0)

agr(\phi_a, you_0) = True

new_chain(\phi_a, you_0)

new_chain: create you_0b

new_chain: create you_0b
```

| Chaining | | | | | |
|----------------------|--------------------------------|-------------------------|------------|-----|-------|
| IO | you ₀ | toy | ϕ | me | that |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata |
| I _{Ob} ^mea | you _{0b} ^ ϕ_{a} | toy _b ^thata | | | |
| | | $toy_{C}^{\phi}a$ | | | |

```
new_chain: exiting

chaining_e_to_n: exiting

chaining_n_to_n: exiting

chaining_n_to_n(\phi, I<sub>0</sub>)

sc(\phi, I<sub>0</sub>) = True

agr(\phi, I<sub>0</sub>) = True

rnr(\phi, I<sub>0</sub>) = True

chaining_e_to_n(\phi_a, I<sub>0</sub>)

agr(\phi_a, I<sub>0</sub>) = True

new_chain(\phi_a, I<sub>0</sub>)

new_chain: create I<sub>0c</sub>

new_chain: create I<sub>0c</sub>^\phi_a
```

| Chaining | | | | | |
|----------------------|--------------------------------|-------------------------|------------|-----|-------|
| I ₀ | you ₀ | toy | ϕ | me | that |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata |
| I _{Ob} ^mea | you _{0b} ^ ϕ_{a} | toy _b ^thata | | | |
| $I_{0c}^{\phi}a$ | | $toy_{C}^{\phi}a$ | | | |

```
new_chain: exiting
                 chaining_e_to_n: exiting
             chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(you∩)
        non refl chaining(you)
             chaining_n_to_n(you0, that)
                 sc(you_0, that) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(you_, me)
                  sc(you_0, me) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(you_0, \phi)
                  sc(you_0, \phi) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(you<sub>0</sub>, toy)
                  sc(you_0, toy) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(you_, I_)
                 sc(you_0, I_0) = True
                  agr(you_0, I_0) = False
             chaining_n_to_n: exiting
        non refl chaining: exiting
    chaining_n: exiting
    chaining_n(I_0)
        non_refl_chaining(I<sub>0</sub>)
             chaining_n_to_n(I_0, that)
                  sc(I_0, that) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(I<sub>0</sub>, me)
                 sc(I_0, me) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(I<sub>0</sub>, \phi)
                  sc(I_0, \phi) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(I<sub>0</sub>, toy)
                  sc(I_0, toy) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(I<sub>0</sub>, you<sub>0</sub>)
                  sc(I_0, you_0) = False
             chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```

| Chaining | | | | | |
|----------------------|--------------------------------|-------------------------|------------|-----|-------|
| I _O | you ₀ | toy | ϕ | me | that |
| I _{0a} | you _{0a} | toya | ϕ_{a} | mea | thata |
| I _{Ob} ^mea | you _{0b} ^ ϕ_{a} | toy _b ^thata | | | |
| $I_{0c}^{\phi}a$ | | toyc ϕ_a | | | |

| Interpretations | | | | | |
|--------------------------------|-------------------------------------|--|--|--|--|
| you _{Ob} ^ ϕ_{a} | toy _b ^that _a | | | | |
| I _{Ob} ^mea you | $0b^{\phi}a$ toy _b thata | | | | |
| I _{OC} ^¢a | toy _b ^that _a | | | | |

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