# Pronouns, Second Edition (Python Version)

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## Contents





## <span id="page-2-0"></span>Preface

Pronouns, Second Edition is a 2024 LaTeX-formatted version of the author's original 1980 Caltech M.S. thesis, Pronouns [\[27\]](#page-132-0). The content of Pronouns, Second Edition is substantially the same as the original Pronouns with the following principal differences:

- OCR'd content of *Pronouns* converted to modern LaTeX style.
- Misspellings, minor grammar points, numberings, and minor technical points are corrected.
- Capitalization changes.
- Figure captions shortened.
- PEP 8 Coding Style identifier spellings.
- Tikz figures replace partially hand-drawn TXT figures.
- LaTeX tabular tables replace TXT tables.
- LaTeX References replace TXT References.
- LaTeX Index added.
- Preface added.

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Postscript and PDF versions of the original Pronouns are also available on the author's PLANETQUANTUM.COM website [\[26\]](#page-132-1).

The author's M.S. thesis and undergraduate adviser was Frederick B. Thompson [\[34\]](#page-132-2).

## <span id="page-3-0"></span>Introduction

Certain substitutions and abbreviations occur in English which are not well understood yet that we would like to understand better so that we may implement them in computer natural language systems intended for man-machine communication. These include pronouns and other function words like those below in Figure 0.1 acting both in isolation and with each other.



#### Figure 0.1. Pronouns and Other Function Words

As well, we have noun phrases modified by demonstratives, Head Deletion, and Equi-NP Deletion. Bloomfield [\[2\]](#page-131-1) defined substitution as a replacement operation.

A substitute is a linguistic form or grammatical feature which, under certain conventional circumstances, replaces any one of a class of linguistic forms. Thus, in English, the substitute I replaces any singular-number substantive expression, provided that this substantive expression denotes the speaker of the utterance in which the substitute is used.

In this thesis we will be concerned with pronouns. Possibly because this will be the only chance we get, we should note the wide variety of substitution mechanisms in general. Examples  $(0.2)-(0.11)$  are from Sag [\[30\]](#page-132-3).

#### (0.2) Do It Anaphor

Jerry won't prove that theorem; Alice will do it.  $[do it = prove that theorem]$ 

#### (0.3) Sentential It Anaphor

I believe that she means business and you'd better believe it too.  $[it = that she means business]$ 

#### (0.4) Null Complement

They asked me to leave but I refused  $\phi.$  $[\phi = \text{to leave}]$ 

#### (0.5) Ones Pronominalization

Betsy has a blue car, and Randy has a red one.  $[one = car]$ 

#### (0.6) Verb Phrase Deletion

Joan wouldn't eat a Quarter Pounder, but Annie would  $\phi$ .  $\phi =$  eat a Quarter Pounder

## (0.7) Sluicing

Someone has drunk my entire six-pack of Schlitz Light, but I don't know who  $\phi$ .

 $\phi =$  has drunk my entire six-pack of Schlitz-Light

## (0.8) Stripping

Gwendolyn snorts cocaine, but  $\phi_1$  not  $\phi_2$  in her own apartment.  $[\phi_1 = \text{Gwendolyn (does)}, \phi_2 = \text{snort cocaine}]$ 

### $(0.9)$  Gapping

Erichman duped Haldeman and Nixon  $\phi$  Mitchell.  $[\phi = \text{duped}]$ 

## (0.10) Conjunction Reduction

Mitchell lied to the committee and  $\phi$  was sentenced last year.  $[\phi = \text{ Mitchell}]$ 

#### (0.11) So Anaphor

Mitchell said he was innocent and Nixon said so too.  $[so = he was innocent]$ 

To this list we can add pronominalizations. Examples  $(0.12)-(0.14)$  are from Lees and Klima [\[22\]](#page-132-4).

## (0.12) Reflexive Pronominalization

Mary's father supported himself.  $[\text{himself} = \text{Mary's father}]$ 

#### (0.13) Pronominalization

Mary's father supported her.  $[her = Mary]$ 

#### (0.14) Reciprocal Pronominalization

John and Mary kissed each other. [each other = John and Mary]

And we might add (0.15) and (0.16) as well.

## (0.15) Head Deletion Joan's cat purrs but Mary's  $\phi$  doesn't.  $[\phi = \text{cat}]$

#### (0.16) Equi-NP Deletion

John is a fraid of  $\phi$  cutting himself.  $\phi =$  John's

Clearly, this list starts to grow very large with addition or refinement and it is probably safe to say that many volumes could be written on substitution processes without putting it to bed. This thesis is about pronouns and chaining of pronouns, and so is much narrower in scope. But this is not much comfort if the goals are not clearly in sight. We are just as lost in the middle of Lake Michigan as we are in the middle of the Pacific Ocean if we don't have a horizon to steer us by.

Part of the problem with investigations of anaphora today is that there is no horizon to steer by. Even though work on anaphora continues in an intelligent way, little progress is being made towards a really comprehensive theory. Instead we have a lot of scattered and independent results.

One goal of this thesis, besides talking about pronouns, is to seek out an algorithmic framework on which to build theory. Accordingly, various data structures such as nodes, C-S-N trees, and chaining tables are created for this purpose. Hopefully, the reader will recognize these data structures as too simplistic and will be moved to improve upon them. This thesis is, by no means at all, a solution to pronouns. At best, it may be a small compass in the middle of Lake Michigan, but this is our approach.

## <span id="page-5-0"></span>1 Fundamentals

## <span id="page-5-1"></span>1.1 Introduction

This chapter describes notation and basic ideas that will be used throughout this thesis. Hopefully, most of the notation described in this chapter is already familiar to the reader, but if not, then this chapter should be self-contained enough to be understandable by a reader with less experience.

## <span id="page-5-2"></span>1.2 Sentences

Sentences are numbered and are kept separate from the text of discussion for ease of reference. For example, (1.1) is from Huddleston [\[13\]](#page-131-2) and is an example of a Bach Peters sentence.

(1.1) The boy who was fooling her kissed the girl who loved him.

**Ungrammatical sentences** are prefixed with an asterisk  $(*)$  and sentences of questionable grammaticality are prefixed with a question mark  $(?)$ . Here,  $(1.3)$  is from Chomsky [\[5\]](#page-131-3).

- (1.2) \*John killed herself.
- (1.3) ?Colorless green ideas sleep furiously.

Subscripts are used to indicate identity between constituents, meaning roughly that they mean the same thing or denote the same referent. More properly, we may think of constituents having the same subscript as being chained together. Below, (1.4) and (1.5) are from Bresnan [\[3\]](#page-131-4).

- (1.4) Some studentsI think theyI are smarter than theyI are.
- (1.5) \*Some studentsI think some studentsI are smarter than some studentsI are.

Sometimes we enclose information in brackets at the beginning or end of a sentence. This same notation is also sometimes used as an alternative to subscripts in identifying constituents. Here, (1.6) is from Bresnan [\[3\]](#page-131-4), (1.7) and (1.8) are from Roberts [\[28\]](#page-132-5) and (1.9) is from Bloom and Hayes [\[1\]](#page-131-5).

- (1.6) My uncle has never ridden a camel but his brother has, although it was  $lame.$  [it = camel]
- (1.7) Men are mortal. [All men are mortal]
- (1.8) Men are waiting. [Some men are waiting]
- (1.9) [Seeing a picture of John Smith] That's John Smith.

A deletion site is indicated by a  $\phi$ . Example (1.10) is from Hockett [\[12\]](#page-131-6).

(1.10) I like the fresh candy better than the stale  $\phi$ . [ $\phi = \text{candy}$ ]

Deletion sites arising from transformations like Equi-NP Deletion are treated similar to pronouns in this paper. Although there are many different kinds of deletion sites with distinct properties, we won't pay attention to this distinction in this thesis.

The symbol  $=$  is used between sentences to indicate that they are equivalent, while the symbol  $\neq$  is used between sentences to indicate that they are not equivalent. Below,  $(1.11)-(1.14)$  are from Ross  $[29]$ .

- $(1.11)$  If John can, he will do it.  $=$
- (1.12) If he can, John will do it.
- $(1.13)$  John will do it if he can.  $\neq$
- (1.14) He will do it if John can.

## <span id="page-7-0"></span>1.3 Noun Phrases

Quantified noun phrases are noun phrases modified by quantifiers. Examples (1.15)-(1.18) are quantified noun phrases.

- (1.15) all female astronauts
- (1.16) at least 10 sexual perverts
- (1.17) many notorious criminals
- (1.18) nearly a dozen Unicorns

Genitives are possessive noun phrases. Examples  $(1.19)-(1.22)$  are genitives.

- (1.19) Uncle Iggy's
- (1.20) my cobra's
- (1.21) the Nazi war criminal's
- (1.22) the alien creatures'

A noun phrase can be generic, specific, or nonspecific, indicated respectively by  $(1.23)-(1.25)$  from Kuno  $[16]$ .

- (1.23) A cat is a malicious animal. [generic]
- (1.24) I have a cat at home, but hate it. [specific]
- (1.25) I want to get a cat for myself. [nonspecific]

A plural noun phrase can be collective or distributive. Examples (1.26)- (1.28) are from Fauconnier [\[7\]](#page-131-8).

- (1.26) The men gathered. [collective]
- (1.27) The men took off their hats. [distributive]
- (1.29) The men carried the couch. [ambiguous]

Sentence  $(1.29)$  is **ambiguous** because it can mean either  $(1.30)$  or  $(1.31)$ .

- (1.30) Each man of the men carried the couch.
- (1.31) The team of men carried the couch.

Smith [\[33\]](#page-132-7) has also noticed this distinction. This explains why (1.32)-(1.35) below are ambiguous.

- (1.32) John and Mary bought the new book by John Steinbeck.
- (1.33) Bricks and stones make strong walls.
- (1.34) George and Marmaduke have dogs.
- (1.35) Gerry likes ice cream and cake.

## <span id="page-8-0"></span>1.4 Pronouns

Pronouns are cross-classified by person, plural, gender, animate, reflexive, attributive possessive, and predicative possessive features among others. First person pronouns are given in  $(1.36)$ .

(1.36) I, me, myself, my, mine, we, us, our, ours, ourselves.

Second person pronouns are given in  $(1.37)$ .

(1.37) you, yourself, yourselves, your, yours

Third person pronouns are given in (1.38).

(1.38) she, he, it, they, her, him, them, herself, himself, itself, themselves, his, its, their, hers, theirs

Singular pronouns are given in (1.39).

(1.39) I, me, myself, my, mine, you, yourself, your, yours, she, he, it, her, him, herself, himself, itself, his, its, hers

Plural pronouns are given in (1.40).

(1.40) we, us, our, ours, ourselves, you, yourselves, your, yours, they, them, themselves, their, theirs

Pronouns with female gender are given in (1.41).

(1.41) she, her, herself, hers

Pronouns with male gender are given in (1.42).

(1.42) he, him, himself, his

Animate pronouns are given in (1.43).

(1.43) I, me, myself, mine, you, yourself, yourselves, your, yours, she, he, they, her, him, herself, himself, themselves, his, their, hers, theirs

Inanimate pronouns are given in (1.44).

(1.44) it, they, them, itself, themselves, its, theirs

Reflexive pronouns are given in (1.45).

(1.45) myself, yourself, yourselves, herself, himself, itself, themselves

Attributive possessive pronouns are given in (1.46).

(1.46) my, your, her, his, its, their

Predicative possessive pronouns are given in  $(1.47)$ .

(1.47) mine, yours, hers, his, its, theirs

Besides the pronouns given above, we also have ones pronouns and reciprocal pronouns. Ones pronouns are given in (1.48).

(1.48) one, oneself, one's

Reciprocal pronouns are given in (1.49).

(1.49) each other, one another, each other's, one another's

#### <span id="page-9-0"></span>1.5 Features

We use three kinds of features in this thesis. The symbol **+** indicates presence of a feature. The symbol **-** indicates absence of a feature. And the symbol **?** indicates that the presence or absence of a feature is either unspecified or not applicable. In the coming chapters, we will speak of agreement of features. A ? feature agrees with any other feature. The only time two features do not agree is when we are comparing  $a +$  and  $a -$  feature. Using  $=$  to indicate agreement and  $\neq$  to indicate nonagreement, we have Figure 1.50.

 $+ = +$   $+ = ?$   $+ \neq -$ <br>  $? = +$   $? = ?$   $? = ? = +$  ? = ?  $- \neq + - = ? - =$ 

#### Figure 1.50. Agreement and Nonagreement between Features

## <span id="page-9-1"></span>1.6 Parse Trees

Sentence parse trees are only drawn schematically in this thesis as extra detail is unnecessary. Parse trees shown more or less represent the surface structure of a sentence. Clause dominating nodes are labelled S and clause conjoining nodes are labelled C. In this thesis, genitives and adjectives are not treated as arising from transformations, but as occuring in the base component. Below, example  $(1.51)$  is from Huddleston [\[13\]](#page-131-2) and example  $(1.52)$  is from Grosu [\[9\]](#page-131-9).

(1.51) The man who lives next door said that he would mow my lawn.



(1.52) Somebody seduced Bill's sister, but no one will ever seduce Jack's and she knows it.



## <span id="page-10-0"></span>1.7 Clauses

Adverbial clauses are clauses beginning with an adverb. Some examples are (1.53)-(1.57) below.

- (1.53) after Fido made a mess on the carpet
- (1.54) before George kisses Betty
- (1.55) since John is an asshole
- (1.56) until Cathy behaves herself
- (1.57) although Lile flunked all his classes

Clauses complemented with that are **that clauses**. Example  $(1.58)$  is a that clause.

(1.58) that Snoopy is a cat

Clauses modified by the For-To Transformation are infinitive clauses. Example (1.59) is an infinitive clause.

(1.59) for Ruth to choose

Clauses modified by the Possessive-Ing Transformation are genitive clauses. Example (1.60) is a genitive clause.

(1.60) Mary's kissing Bob

Clauses modified by WH-Fronting Transformation but not the Question Transformation and which modify noun phrases are relative clauses. Examples (1.61)-(1.65) are relative clauses.

- (1.61) who ate five hamburgers
- (1.62) that has a leaky faucet
- (1.63) which doesn't run
- (1.64) whom he gave it to

(1.65) whose life isn't worth a postage stamp

Clauses without embedded subordinate clauses are simplex. In example (1.66) from Ross[\[29\]](#page-132-6), the simplexes are (1.67)-(1.69). In example (1.70), from Huddleston [\[13\]](#page-131-2), the simplexes are  $(1.71)-(1.73)$ . In example  $(1.74)$ , from Huddleston, the simplexes are (1.75)-(1.77).

(1.66) Realizing that he was unpopular didn't disturb Oscar.



- ϕ precedes he  $\phi$  commands he ϕ precedes Oscar Oscar *commands*  $\phi$ Oscar commands he
- (1.67) S didn't disturb Oscar
- $(1.68)$   $\phi$ 's realizing that S
- (1.69) he was unpopular
	- (1.70) My neighbor who is pregnant said that she was very happy.



neighbor precedes she neighbor commands she

- (1.71) my neighbor said that S
- (1.72) who is pregnant
- (1.73) she was very happy

(1.74) The pilot who shot at it hit the Mig that chased him.



pilot precedes him pilot commands him it precedes the Mig Mig commands it

- (1.75) the pilot hit the Mig
- (1.76) who shot at it
- (1.77) that chased him

## <span id="page-12-0"></span>1.8 Precedes and Commands

The **precedes** and **commands** relations, first described by Langacker [\[19\]](#page-132-8), are defined below in  $(1.78)$  and  $(1.79)$ .

#### (1.78) precedes Relation

A node A precedes another node B if

- (a) neither A nor B dominates the other, and
- (b) A occurs before B (in preorder traversal)

#### (1.79) commands Relation

A node A commands another node B if

- (a) neither A nor B dominates the other, and
- (b) the S-node that most immediately dominates A also dominates B

Another relation that will be useful is the **is separate from** relation defined below in  $(1.80)$ .

#### (1.80) is separate from Relation

- A node R is separate from another node B if
- (a) neither A nor B dominates the other, and
- (b) the lowest node in the tree dominating A and B is a C-node.

We will see that *precedes*, *commands*, and *is separate from* are useful in determining when pronominalization is or isn't possible.

In example  $(1.81)$ , A precedes B, A commands B, and E commands A. We don't have A precedes A, B precedes A, B precedes B, A commands A, or B commands B.

(1.81)



A precedes B A commands B B commands A

In (1.82), A precedes B, A is separate from B, and B is separate from A. In (1.83), A precedes B and A commands B. In (1.84), A precedes B and B commands A. In (1.85), A precedes B, H is separate from B, and B is separate from A.

(1.82)



A precedes B A is separate from B A precedes B

(1.83)



A precedes B A commands B

(1.84)



A precedes B B commands A

(1.85)



A precedes B A is separate from B B is separate from A

Examples (1.86)-(1.89) are from Langacker [\[19\]](#page-132-8).

(1.86) The mosquito which bit Algernon was killed by him. [him = Algernon]



Algernon precedes him him precedes Algernon  $(1.87)$  The mosquito which bit him was killed by Algernon. [him = Algernon]



him precedes Algernon Algernon precedes him

(1.88) Algernon killed the mosquito which bit him. [him = Algernon]



Algernon precedes him him commands Algernon

(1.89) He killed the mosquito which bit Algernon. [he  $\neq$  Algernon]



he precedes Algernon Algernon commands he

The Precedes and Commands Rule, essentially as stated by Langacker [\[19\]](#page-132-8), is given in (1.90) below.

#### (1.90) Precedes and Commands Rule

A pronoun P may be used to pronominalize a noun phrase NP unless

(a) P precedes NP, and

(b) P commands NP or P is separate from NP

Note that the Precedes and Commands Rule explains the grammaticality and ungrammaticality of (1.86)-(1.89). These further examples from Ross [\[29\]](#page-132-6) should drive the point home.

- (1.91) After John Adams woke up, he was hungry. [he = John Adams]
- $(1.92)$  That Oscar was unpopular didn't disturb him. [him = Oscar]
- (1.93) For your brother to refuse to pay taxes would get him into trouble. [him = your brother]
- $(1.94)$  Anna's complaining about Peter infuriated him. [him = Peter]
- $(1.95)$  The possibility that Fred will be unpopular doesn't bother him. [him = Fred]



NP precedes P P commands NP

- (1.96) After he woke up, John Adams was hungry. [he = John Adams]
- $(1.97)$  That he was unpopular didn't disturb Oscar. [he = Oscar]
- (1.98) For him to refuse to pay taxes would get your brother into trouble. [him = your brother]
- (1.99) Anna's complaining about him infuriated Peter. [him = Peter]
- $(1.100)$  The possibility that he will be unpopular doesn't bother Fred. [him = Fred]



P precedes NP NP commands P

 $(1.101)$  John Adams was hungry after he woke up.  $[he = John Adams]$ 

- $(1.102)$  Oscar wasn't disturbed that he was unpopular. [he = Oscar]
- (1.103) It would get your brother into trouble for him to refuse to pay taxes. [him = your brother]
- (1.104) Peter was infuriated at Anna's complaining about him. [him = Peter]
- $(1.105)$  Fred isn't bothered by the possibility that he will be unpopular. [he = Fred]



NP precedes P NP commands P

- (1.106) \*He was hungry after John Adams woke up. [he = John Adams]
- $(1.107)$  \*He wasn't disturbed that Oscar was unpopular. [he = Oscar]
- (1.108) \*It would get him into trouble for your brother to refuse to pay taxes. [him = your brother]
- $(1.109)$  \*He was infuriated at Anna's complaining about Peter. [he = Peter]
- $(1.110)$  \*He isn't bothered by the possibility that Fred will be unpopular. [he = Fred]



P precedes NP P commands NP

Examples (1.111) and (1.112) from Langacker [\[19\]](#page-132-8) illustrate the Precedes and Commands Rule for conjoined structures.

(1.111) Penelope cursed Peter and slandered him. [him = Peter]



Peter precedes him Peter is separate from him him is separate from Peter

 $(1.112)$  \*Penelope cursed him and slandered Peter. [him = Peter]



him precedes Peter him is separate from Peter Peter is separate from him

Examples (1.113) and (1.114) adapted from Chiba [\[4\]](#page-131-10) involve Equi-NP Deletion.

(1.113) The interest in visiting Las Vegas that Mary displayed is typical of gamblers.



 $\phi$  precedes Mary

## <span id="page-19-0"></span>2 Resolution Module

## <span id="page-19-1"></span>2.1 Introduction

In the previous chapter we touched upon some basic notions such as the precedes, commands, and is separate from relations. We will see in the coming chapters how these concepts give rise to a very promising approach to the problem of pronoun resolution.

The algorithm we shall describe won't be complete in the sense that we will elaborate and refine it in later chapters and after we are done it will need elaboration and refinement, but it will be set in firm soil so that we have a foundation on which to build. Because personal and reflexive pronouns are easiest, these are the pronouns we shall consider first. But before we go any farther, let us take time out to indicate something of the environment and structure of the module that does resolving of pronouns in a natural language system, the Resolution module.

## <span id="page-19-2"></span>2.2 Environment

The center of a natural language system is the Language Processor module which is divided into five submodules. These are the Language Driver, Preprocesor, Parser, Semantic Processor, and Output Processor as indicated in Figure 2.1.



#### Figure 2.1. Submodules of the Language Processor

Briefly, from the point of view of the Language Processor, the following happens. A user types input at a terminal which is picked up by the Operating System of the natural language system. The Operating System maintains information about the user including the language version he is in as well as his state in that version. The user's state is known as his prefix. The Operating System, after picking up a user's input calls a Process Input routine of the Language Driver in the Language Processor. Once in the Language Driver, the first module to be called upon is the Preprocessor.

The Preprocessor in the Language Processor compresses blanks in the input string, straps right and left delimiters about it, recognizes and builds parsing

graph arcs over identifiers and numbers, and looks the identifiers up in the lexicon. After calling the Preprocessor, the Language Driver calls the Parser.

The Parser in the Language Processor parses the output. of the Preprocessor using an algorithm such as the Kay algorithm and can handle any general rewrite rule grammar. Of course, since a sentence may be ambiguous, more than one system parse tree may be passed back by the Parser. If no good parsings are found, then the Syntax Diagnostics routine of the Syntax Diagnostics module of the natural language system is called. Otherwise, if there are good parsings, then the Language Driver calls the Semantic Processor on the output of the Parser.

The Semantic Processor is driven by the syntax of a system parse tree into making calls on semantic routines which can be postprocedures (called on their arguments after their arguments evaluate themselves), preprocedures (called on their arguments before their arguments evaluate themselves), and syntax procedures (called at syntax time during parsing before preprocedures and postprocedures are called during semantic processing). On return to the Language Driver, the Language Driver calls the Output Processor on the output of the Semantic Processor.

The Output Processor does some relatively menial processing such as removing duplicate lines from the output line list which will be sent back to the Operating System. The Output Processor is able to handle ambiguous output and removes diagnostic messages if at least one of the outputs is good.

On completion of the call on the Output Processor, the Language Driver returns to the Operating System and the Operating System displays the output line list on the user's terminal, at the same time updating its information on the user.

From the discussion of the *precedes, commands,* and *is separate from* relations in the previous chapter, we know that information about the syntax of the input sentence is critical to the resolving of pronouns in the input sentence. On the other hand, for semantic processing to carry out the processing it needs to carry out, the placing of information on the chaining of pronouns must already be placed in the system parse tree of the input sentence.

The logical conclusion of these two observations indicates that pronoun resolution takes place after parsing, but before semantic processing. This relationship of the Resolution module with the other modules of the Language Processor is indicated in Figure 2.2.



#### Figure 2.2. Resolution Module within the Language Processor

In practice, this formulation may not be quite correct because there can be other versions than English which will have nothing to do with the Pronoun Resolution module and so what we end up doing is making the Resolution module accessible via a semantic preprocedure which is associated with the parsing of the right delimiter of a sentence. So instead, what happens is that the first semantic preprocedure to be called will be the procedure which handles Pronoun Resolution.

## <span id="page-21-0"></span>2.3 Structure inside the Resolution Module

The Resolution module is partitioned into seven submodules besides a Global Declarations module. These are the Node Processor, Parser, Primary Utilities, Secondary Utilities, Table Processor, Table Interpreter, and Resolution Driver modules. The reader should not confuse the Parser of the Language Processor with the Parser of the Pronoun Resolution module which have entirely different functions. The relationship of these submodules of the Resolution module is indicated below in Figure 2.3.



Figure 2.3. Structure of the Resolution Module

Not shown is the Global Declarations module which does not have any procedures itself, but merely defines data structures. The Global Declarations submodule is accessible by all other submodules of the Resolution module.

## <span id="page-22-0"></span>3 Global Declarations

The Global Declarations module defines the data structures accessible to other modules within the pronoun resolution module. The Global Declarations module is shown below in Figure 3.1.

```
#globals.py
import sys
from typing import TextIO
from enum import Enum, IntEnum
from typing import Optional
class FeatureIndex(IntEnum): #Feature indices.
    PNF = 0 #Pronoun Feature
    FPF = 1 #First Person Feature
    SPF = 2 #Second Person Feature
    TPF = 3 #Third Person Feature
    PLF = 4 #Plural Feature
    GNF = 5 #Gender Feature
   ANF = 6 #Animate Feature
   RPF = 7 #Reflexive Feature
    GEN = 8 #Genitive Feature
N_FEATURES = len(FeatureIndex) #Number of Features
class NodeId(Enum): #Identifies the type of node.
    C_NODE = 0 #Represents a C-node.
    S_NODE = 1 #Represents an S-node.
    N_NODE = 2 #Represents an N-node.
    E_NODE = 3 #Represents an E-node.
class Feature(Enum): #Linguistic feature values.
    PLUS = 0 #Has this feature.
    MINUS = 1 #Doesn't have this feature.
    QUESTION = 2 #Might or might not have this feature.
Features = list[Feature] #list of Feature enums
```

```
Figure 3.1. Global Declarations Module (Part I)
```

```
class Node: #Base node class
    # Current tree
   _tree: Optional['Node'] = None
   def __init__(self):
        self.number: int = 0
        self.up_link: Optional['Node'] = None
        self.down link: Optional['Node'] = None
        self.left_link: Optional['Node'] = None
        self.right_link: Optional['Node'] = None
        self.thread_link: Optional['Node'] = None
        self.np_link: Optional['Node'] = None
        self.chain_link: Optional['Node'] = None
        self.col_link: Optional['Node'] = None
        self.ftr: Features = [Feature.QUESTION] * N_FEATURES
        self.id: NodeId = NodeId.C_NODE
        self.lit: str = ""
        self.end_col_link: Optional['Node'] = None
        self.pred_link: Optional['Node'] = None
        self.succ_link: Optional['Node'] = None
        self.sub: str = '@classmethod
   def tree(cls) -> Optional['Node']:
        return cls._tree
    @classmethod
   def set_tree(cls, new_tree: 'Node') -> None:
        cls. tree = new tree
```
#### Figure 3.1. Global Declarations Module (Part II)

Basically, our data structures are C-S-N trees, chaining tables, and the nodes they involve. It will help to get some feel for these data structures before we go on to other chapters.

#### <span id="page-24-0"></span>3.1 Nodes

There are four kinds of nodes: C-nodes, S-nodes, N-nodes, and E-nodes. Cnodes, S-nodes, and N-nodes occur in C-S-N trees and correspond to conjoined structures, sentences, and noun phrases. E-nodes occur in chaining tables. The fields of the C-nodes, S-nodes, N-nodes, and E-nodes are as indicated in Figure 3.1.

## <span id="page-24-1"></span>3.2 C-S-N Trees

A C-S-N tree has three kinds of nodes: C-nodes, S-nodes, and N-nodes. Link fields which are relevant to C-S-N trees are up\_link, down\_link,

left\_link, right\_link, thread\_link, pred\_link, and succ\_link. An example of a C-S-N tree is given in Figure 3.2.



Figure 3.2. C-S-N Tree

## <span id="page-25-0"></span>3.3 Chaining Tables

A chaining table contains N-nodes, E-nodes, and one S-node for keeping track of the chaining table. Link fields relevant to chaining tables are np\_link, chain\_link, col\_link, end\_col\_link, pred\_link, and succ\_link. Chaining tables and C-S-N trees are connected through their N-nodes. An example of a chaining table is given in Figure 3.3.



Figure 3.3. Chaining Table

### <span id="page-25-1"></span>3.4 C-Nodes

A C-node has the following fields: up\_link, down\_link, left\_link, right\_link, thread\_link, and number. C-nodes correspond to conjoined sentences and conjoined subordinate clauses.

## <span id="page-26-0"></span>3.5 S-Nodes

An S-node has exactly the same fields as a C-node and is only distinguished from a C-node by its NodeId. S-nodes correspond to sentences and subordinate clauses.

#### <span id="page-26-1"></span>3.6 N-Nodes

An N-node has the following fields: lit, ftr, up\_link, down\_link, thread\_link, np\_link, chain\_link, col\_link, end\_col\_link, pred\_link, succ\_link, and number. N-nodes correspond to noun phrases without attached subordinate clause modifiers.

## <span id="page-26-2"></span>3.7 E-Nodes

An E-node has the following fields: sub, ftr, np\_link, chain\_link, and col\_link. An E-node may be thought of as a copy of its np\_link with a slightly more defined set of features.

## <span id="page-26-3"></span>3.8 **lit** Field

The lit field of an N-node is a string pointer to the string that the Nnode represents. The lit field is actually unnecessary in an N-node, but is convenient for displaying intermediate results. Function view\_node\_str of the Node Processor and some other procedures that display intermediate results use this field.

## <span id="page-26-4"></span>3.9 **sub** Field

The sub field of an E-node is a character representing the subscript of the E-node. The sub field of an E-node, like the lit field of an N-node, is an unnecessary field, but is convenient for displaying intermediate results.

## <span id="page-26-5"></span>3.10 **ftr** Field

The ftr field of an N-node or E-node is an array of Feature's representing the feature set of the N-node or E-node to which it corresponds. A Feature can be a PLUS, MINUS, or QUESTION as described in the previous chapter. The offsets PNF, FPF, SPF, TPF, PLF, GNF, ANF, and RPF are used to access elements of the ftr array. The accessed elements are pronoun feature, first person feature, second person feature, third person feature, plural feature, gender feature, animate feature, and reflexive feature. The number of Feature's is N\_FEATURES. Figure 3.4 shows some examples of the settings of ftr for some typical noun phrases.

<b>Features</b>								
	PNF	FPF	<b>SPF</b>	TPF	PLF	<b>GNF</b>	<b>ANF</b>	<b>RPF</b>
John				$^{+}$			$^{+}$	
flowers				$+$	$^{+}$	っ		
he	$+$			$+$			$^{+}$	
them	$+$			$^{+}$	$^{+}$	っ	っ	
I	$+$	$^{+}$				っ	$^{+}$	
you	$^{+}$		$+$			2	$^{+}$	
her	$+$			$+$		$+$	$^{+}$	
myself	$^{+}$	$^{+}$				っ	$^{+}$	
herself	$+$			$^{+}$		$+$	$^{+}$	
itself	$+$			$^{+}$		っ		

Figure 3.4. **ftr** Settings for Some Typical Noun Phrases

## <span id="page-27-0"></span>3.11 **up link** Field

The up\_link field of a C-node, S-node, or N-node links to the parent node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with up\_link's shown is given in Figure 3.5.



Figure 3.5. C-S-N Tree with **up\_link**'s Shown

## <span id="page-27-1"></span>3.12 **down link** Field

The down\_link field of a C-node, S-node, or N-node links to the first child node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with down\_link's shown is given in Figure 3.6.



Figure 3.6. C-S-N Tree with **down\_link**'s Shown

### <span id="page-28-0"></span>3.13 **left link** Field

The left\_link field of a C-node, S-node, or N-node links to the left brother node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with left\_link's shown is given in Figure 3.7.



Figure 3.7. C-S-N Tree with **left\_link**'s Shown

### <span id="page-28-1"></span>3.14 **right link** Field

The right link field of a C-node, S-node, or N-node links to the right brother node of the C-node, S-node, or N-node in the C-S-N tree in which it occurs. An example of a C-S-N tree with right\_link's shown is given in Figure 3.8.



Figure 3.8. C-S-N Tree with **right\_link**'s Shown

### <span id="page-29-0"></span>3.15 **thread link** Field

The thread\_link field of a C-node, S-node, or N-node links to the first node traversed after the C-node, S-node, or N-node in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with thread\_link's shown is given in Figure 3.9.



Figure 3.9. C-S-N Tree with **thread\_link**'s Shown

## <span id="page-29-1"></span>3.16 **number** Field

C-node, S-node, or N-node have a number field which is the number that would be assigned to that node if the nodes of the C-S-N tree in which it occurs are numbered in a preorder traversal. An example of a C-S-N tree with number fields shown is given in Figure 3.10.



Figure 3.10. C-S-N Tree with **number** Fields Shown

### <span id="page-30-0"></span>3.17 **np link** Field

For an E-node, the np\_link is the N-node to which it is attached. Conceptually, we think of the E-node as being a copy of the N-node except for its subscript and different set of Feature's, chain\_link, and col\_link. The np\_link is just a way of avoiding duplication of information. For an N-node, the np\_link is always itself. An example of a chaining table with np\_link's shown is given in Figure 3.11.



Figure 3.11. Chaining Table with **up\_link**'s Shown

#### <span id="page-30-1"></span>3.18 **chain link** Field

The chain link of an E-node is another E-node representing the substitute to which the first E-node is attached. When chaining is obligatory, an N-node is chained to an N-node. An example of a chaining table with chain\_link's shown is given in Figure 3.12.



Figure 3.12. Chaining Table with **chain\_link**'s Shown

## <span id="page-31-0"></span>3.19 **col link** Field

The col\_link field of an E-node or N-node links together the elements of a column in a table. An N-node is always on top of a column with E-nodes underneath. An example of a chaining table with col\_link's shown is given in Figure 3.13.



Figure 3.13. Chaining Table with **col\_link**'s Shown

## <span id="page-31-1"></span>3.20 **end col link** Field

The end\_col\_link field of an N-node links to the end of the column of E-nodes lying under this N-node. An example of a chaining table with end\_col\_link's shown is given in Figure 3.14.



 $S_1$ 

Figure 3.14. Chaining Table with **end\_col\_link**'s Shown

## <span id="page-32-0"></span>3.21 **pred link** Field

The pred\_link field of an N-node links to the preceding N-node found in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with pred\_link's shown is given in Figure 3.15.



Figure 3.15. C-S-N Tree with **pred\_link**'s Shown

An example of a chaining table with pred\_link's shown is given in Figure 3.16



Figure 3.16. Chaining Table with **pred\_link**'s Shown

## <span id="page-33-0"></span>3.22 **succ link** Field

The succ\_link field of an N-node links to the succeeding N-node found in a preorder traversal of the C-S-N tree in which it occurs. An example of a C-S-N tree with succ\_link's shown is given in Figure 3.17.



Figure 3.17. C-S-N Tree with **succ\_link**'s Shown

An example of a chaining table with succ\_link's shown is given in Figure 3.18.



Figure 3.18. Chaining Table with **succ\_link**'s Shown

## <span id="page-34-0"></span>4 Node Processor

The Node Processor module contains functions new\_c\_node, new\_s\_node, new\_n\_node, and new\_e\_node and has the skeleton shown below in Figure 4.1.

```
#node_proc.py
from globals import *
def new_c_node() -> Node:
def new_s_node() -> Node:
def new_n_node() -> Node:
def new_e_node() -> Node:
```
#### Figure 4.1. Skeleton of the Node Processor

new\_c\_node, new\_s\_node, new\_n\_node, and new\_e\_node generate, respectively, a new C-node, S-node, N-node, or E-node, with their fields initialized and are rather straightforward functions. These are shown below in Figures 4.2-4.5.

Function new\_c\_node returns a new C-node.

def new\_c\_node() -> Node: return new\_node(NodeId.C\_NODE)

#### Figure 4.2. Function **new\_c\_node**

Function new\_s\_node returns a new S-node.

```
def new s node() -> Node:
    return new_node(NodeId.S_NODE)
```
#### Figure 4.3. Function **new\_s\_node**

Function new\_n\_node returns a new N-node.

```
def new_n_node() -> Node:
   answer = new node(NodeId.N_NODE)
   answer.lit = ""
   answer.ftr = [Feature.QUESTION] * N_FEATURESanswer.end_col_link = None
   answer.pred_link = None
   answer.succ_link = None
    answer.np_link = answer
    return answer
```
#### Figure 4.4. Function **new\_n\_node**

Function new\_e\_node returns a new E-node.

```
def new_e_node() -> Node:
    answer = new\_node(Noded.A.E_NODE)answer.sub = ' '
    answer.ftr = [Feature.QUESTION] * N_FEATURESreturn answer
```
Figure 4.5. Function **new\_e\_node**

#### <span id="page-35-0"></span>4.1 Function **view node str**

There is one output procedure in the Node Processor that has not been discussed above that we need to know about, because we will be looking at some of its output for short while. This is procedure view\_node\_str which takes as an argument a NodePointer and outputs it in readable form. Otherwise, procedure view\_node\_str does no processing of its own, and so we do not need to know the details of its inner workings. For us it is enough to be able to understand the output. Function view\_node\_str has the form indicated in Figure 4.6.

```
def view_node_str(node: Node) -> str:
    """Formatted string representation of node."""
    ...
```
Figure 4.6. Skeleton of Function **view\_node\_str**
Some typical output of procedure view\_node\_str is shown below in Figure 4.7 where a chaining table is listed. (Links from the chaining table to its associated C-S-N tree are also listed by procedure view\_node\_str.)

<b>Nodes</b>									
1	(C, up:0, dn:2, lt:0, rt:0, th:2, nu:1)								
$\bf{2}$	$(S, up:1, dn:3, It:0, rt:5, th:3, nu:2)$								
$\bf{3}$	(N, lit:June, ftr: [---+-++-], up:2, dn:0,								
	1t:0, rt:4, th:4, np:3, ch:0, co:3a, ec:3 <sub>b</sub> ,								
	$pr:0$ , su:4, nu:3)								
3a	(E, sub:A, ftr: $[- - + + + -]$ , np:3, ch:0, co:3 <sub>b</sub> )								
$\mathbf{3_{h}}$	(E, sub:B, ftr: $[- - + - + + -]$ , np:3, ch: $6a$ , co:0)								
4	(N, lit: flowers, ftr: $[- - + + ? - -]$ , up: 2, dn: 0,								
	1t:3, rt:0, th:5, np:4, ch:0, co:4a, ec:4b,								
	$pr:3$ , su:6, nu:4)								
4a	(E, sub:A, ftr: $[- - + + ? - -]$ , np:4, ch:0, co:4 <sub>b</sub> )								
4 <sub>b</sub>	(E, sub:B, ftr: $[- - + + ? - -]$ , np:4, ch:7 <sub>a</sub> , co:0)								
5	$(S, up:1, dn:6, lt:2, rt:0, th:6, nu:5)$								
6	(N, lit:she, ftr: $[+ - + + + -]$ , up:5, dn:0,								
	1t:0, rt:7, th:7, np:6, ch:0, co: $6a$ , ec: $6a$ ,								
	$pr:4$ , su: 7, nu: 6)								
6a	$(E, sub:A, ftr:[+--+-++-], np:6, ch:0, co:0)$								
7	(N, lit:them, ftr: $[+ - + + ? ? -]$ , up:5, dn:0,								
	1t:6, rt:0, th:0, np:7, ch:0, co:7a, ec:7a,								
	$pr:6$ , su:0, nu:7)								
7a	$(E, sub:A, ftr:[+ - + + ? ? -], np: 7, ch: 0, co: 0)$								

Figure 4.7. Typical Output from Function **view\_node\_str**

 $(C = C\text{-node}, S = S\text{-node}, N = N\text{-node}, E = E\text{-node}, \text{lit} = \text{lit field}, \text{sub} = \text{sub}$ field,  $ftr = ftr$  field,  $up = up\_link$ ,  $dn = down\_link$ ,  $lt = left\_link$ ,  $rt = right\_link, th = thread\_link, nu = number, np = np\_link, ch$ = chain\_link, co = col\_link, ec = end\_col\_link, pr = pred\_link, and  $su = succ\_link)$ 

# 5 Parser

The Parser module defines function parse and has the form shown below in Figure 5.1.

```
#parse_proc.py
from lexicon import *
from node_proc import *
def parse(obj) -> Node:
```
Figure 5.1. Skeleton of the Parser

Function parse accepts as input a system focus representation and system parse tree that has been generated by a computer natural language system. The output of parse is a C-S-N tree incorporating the information contained in the system parse tree and system focus. The system focus represents the natural language system's focus of attention. This will be gone into in more detail in Chapter 13.

The representation of the system tree inputted to parse is system dependent, and so the details of parse are also system dependent. As the internals of parse are heavily dependent upon and rather involved for any system, we won't go into the details of parse for any particular system here. Hopefully, the reader may glean enough information from the multitude of examples presented in this thesis to get an idea of what parse does. In any case, lack of an actual algorithm for parse isn't so bad since the ideas presented in this thesis are really still in an early stage and it is enough to concentrate on them.

Even though the input to the Parser is not well defined, the output is. The Parser builds from the system parse tree it is given the corresponding C-S-N tree with all up\_link's, down\_link's, left\_link's, right\_link's, thread\_link's, and number's set to what is expected. Consider example (5.2) below.

(5.2) June hates flowers, but she waters them anyway.



When procedure parse is called on the system parse tree representing  $(5.2)$ , we get the following output in Figure 5.3.





## Figure 5.3. Typical Output from **parse**

Listing of nodes in Figure 5.3 is done by procedure view\_node\_str of the Node Processor described in Chapter 4. The C-S-N parse tree is slightly more complicated when focusing is taken into account, but for the time being we will ignore its effects. We will discuss the effects of focusing on C-S-N parse trees in Chapter 13.

When Figure 5.3 is drawn as a tree, we get a structure like Figure 5.4.



Figure 5.4. Output from **parse** Drawn as a Tree

# 6 Primary Utilities

The Primary Utilities module defines the boolean functions precede, command, and separate corresponding to the precedes, commands, and is separate from relations discussed in Chapter 1. The skeleton of the primary utilitites module is shown below in Figure 6.1.

```
#primary_uty.py
from globals import *
def precede(n1: Node, n2: Node) -> bool:
def command(n1: Node, n2: Node) -> bool:
def separate(n1: Node, n2: Node) -> bool:
```
#### Figure 6.1. Skeleton of the Primary Utilities

The precede, command, and separate functions do just what is expected. They are true if and only if the precedes, commands, and is separate from relations hold between their arguments. Along with function dominate which is used by separate, these functions are shown below in Figures 6.2-6.5.

Function precede is true if and only if n1 precedes n2.

```
def precede(n1: Node, n2: Node) -> bool:
   return n1.number < n2.number
```
### Figure 6.2. Function **precede**

Function dominate is true if and only if n1 dominates n2.

```
def dominate(n1: Node, n2: Node) -> bool:
    if n1.number == n2.number:
        return True
    child = n1.down_link
    while child is not None:
        if dominate(child, n2):
            return True
        child = child.right_link
    return False
```
### Figure 6.3. Function **dominate**

Function command is true if and only if n1 commands n2.

def command(n1: Node, n2: Node) -> bool: return dominate(n1.up\_link, n2)

#### Figure 6.4. Function **command**

Function separate is true if and only if n1 is separate from n2.

```
def separate(n1: Node, n2: Node) -> bool:
   parent = n1.up_link
   while not dominate(parent, n2):
       parent = parent.up_link
    return parent.id == NodeId.C_NODE
```
Figure 6.5. Function **separate**

# 7 Secondary Utilities

The Secondary Utiltities module defines functions sc, agr, and rnr. These stand for Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule. The skeleton of the Secondary Utilities module is shown below in Figure 7.1.

```
#secondary_uty.py
from primary_uty import *
def sc(n1: Node, n2: Node) -> bool:
def agr(n1: Node, n2: Node) -> bool:
def rnr(n1: Node, n2: Node) -> bool:
```
Figure 7.1. Skeleton of Secondary Utilities

### 7.1 Syntactic Conditions

As shown in Chapter 1, certain constraints such as the Precedes and Commands Rule apply in forward pronominalization. Function sc is true whenever these grosser syntactic constraints are met. In this thesis, we let sc be true when the Precedes and Commands Rule is satisfied, function sc is shown below in Figure 7.2.

```
def sc(n1: Node, n2: Node) -> bool:
   return not (precede(n1, n2) and (command(n1, n2) or separate(n1, n2)))
```
Figure 7.2. Function **sc** (Syntactic Conditions)

### 7.2 Agreement

Besides satisfying Syntactic Conditions, there has to be agreement between a node and its chaining node. First person, second person, third person, plural, gender, and animate features have to agree in order for one node to chain to another. Function agr is shown below in Figure 7.3.

```
def agr(n1: Node, n2: Node) -> bool:
   ftr1 = n1.ftr
   ftr2 = n2.ftr
   return (eq_feat(ftr1[FeatureIndex.FPF], ftr2[FeatureIndex.FPF]) and
           eq_feat(ftr1[FeatureIndex.SPF], ftr2[FeatureIndex.SPF]) and
            eq_feat(ftr1[FeatureIndex.TPF], ftr2[FeatureIndex.TPF]) and
            eq_feat(ftr1[FeatureIndex.PLF], ftr2[FeatureIndex.PLF]) and
            eq_feat(ftr1[FeatureIndex.GNF], ftr2[FeatureIndex.GNF]) and
            eq_feat(ftr1[FeatureIndex.ANF], ftr2[FeatureIndex.ANF]))
```
Figure 7.3. Function **agr** (Agreement)

### 7.3 Equal Features

Function agr uses function eq\_feat. eq\_feat tests if two Feature's are equal. As indicated in Chapter 1, Feature's are equal unless a PLUS and MINUS are compared. Function eq\_feat is shown below in Figure 7.4.

```
def eq_feat(f1: Feature, f2: Feature) -> bool:
    if f1 == Feature.PLUS:
        return f2 != Feature.MINUS
    elif f1 == Feature.MINUS:
        return f2 != Feature.PLUS
    else: # f1 == Feature.OUESTION
        return True
```
Figure 7.4. Function **eq\_feat** (Equal Features)

### 7.4 Reflexive Nonreflexive Rule

The distinction between reflexive pronouns and nonreflexive pronouns is that reflexive pronouns cannot chain to an N-node that is outside of the same simplex in which it occurs, while a nonreflexive pronoun can. This rule will have to be modified later for genitives, but for now we can suppose that a nonreflexive pronoun must chain to an N-node outside of the same simplex in which it is in. Shown in Figure 7.5 is function rnr which is true when the reflexive nonreflexive rule is satisfied.

```
def rnr(n1: Node, n2: Node) -> bool:
    ftr1 = n1(np\_link.ftr)ftr2 = n2.np_link.ftr
    if ftr2[FeatureIndex.GEN] == Feature.PLUS:
        return False
    elif ftr1[FeatureIndex.RPF] == Feature.PLUS:
        return (n1.up link == n2.up link)
                and (ftr1[FeatureIndex.GEN] == Feature.MINUS)
    elif ftr1[FeatureIndex.RPF] == Feature.MINUS:
        return (n1.up_link != n2.up_link)
                or (ftr1[FeatureIndex.GEN] != Feature.MINUS)
```
Figure 7.5. Function **rnr** (Reflexive Nonreflexive Rule)

## 8 Table Processor I

The Table Processor module defines function chaining which takes as input a C-S-N tree and returns its chaining table. The actions of function chaining in the Table Processor can only be understood by example, and this is what this chapter provides. In Chapter 9, we'll look at the actual algorithms and, in Chapter 10, we'll look at some actual output.

So, let us consider sentence (8.1) below.

(8.1) John wants to give June a present, but he isn't sure she'll like it.



The Parser builds from the system parse tree of (8.1) the corresponding C-S-N tree with six N-nodes which have the lit fields and ftr's indicated below in Figure 8.2.

Features										
	PNF	FPF	<b>SPF</b>	TPF	PLF	<b>GNF</b>	<b>ANF</b>	<b>RPF</b>		
John				$^+$			$^+$			
Φ	$^{+}$	?	?	$\mathcal{P}$	?	ာ	?			
June										
present						?				
he	$^{+}$						$^+$			
she	$^{+}$						$^{+}$			
it						っ				

Figure 8.2. **lit** Fields and **ftr**'s of the N-Nodes

The C-S-N tree itself has the form of Figure 8.3 below.



Figure 8.3. C-S-N Parse Tree

After parse is called, chaining is called. The first thing to happen is the initialization of the chaining table for the C-S-N tree. Below each N-node is suspended, by the col\_link of the N-node, a new E-node with subscript A. Each new E-node has a np\_link back to the N-node it is suspended from. As well, the Feature's of each new E-node are copied from the N-node it is suspended from. Attached to the first and last N-nodes is an S-node to make it easy to keep track of the first and last N-nodes in the chaining table. The chaining table, as it looks immediately after initialization, is shown below in Figure 8.4.



Figure 8.4. Chaining Table Immediately after Initialization

The chaining algorithm works by walking backwards across N-nodes in the top row and walking down columns of E-nodes. The chaining algorithm works on two N-nodes at a time. If the first is compatible with the second under Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule, then the E-nodes underneath the first N-node that agree with the second N-node are chain\_link'ed to copies of the second N-node.

The last N-node in the table is  $\underline{\text{it}}$ , the chaining table begins with  $\underline{\text{it}}$ . it can't chain to itself, so the second N-node in the description above becomes she and the chaining algorithm compares it to she. Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(it, she) = Trueagr(it, she) = False
```
The chaining algorithm now moves from she to he and compares it to he. Again Syntactic Conditions are satisfied, but Agreement isn't.

 $sc(it, he) = True$  $aqr(it, he) = False$ 

The chaining algorithm moves from he to present and compares it to present. This time, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule are satisfied.

```
sc(it, present) = Trueagr(it, present) = Truernr(it, present) = True
```
Since all three rules are satisfied, a chain from  $it_a$  to a copy of present may be created. This happens if  $it_a$  and present agree, and they do.

 $agr(it<sub>a</sub>, present) = True$ 

The chaining algorithm makes a new E-node copy of  $\overline{\mathrm{present}}$ ,  $\overline{\mathrm{present}}$ <sub>b</sub>, and hangs it below <u>present</u>. The chain\_link of  $\frac{\text{present}}{\text{b}}$  is set to  $\frac{\text{it}}{\text{it}}$  and the semantic features of  $it_{a}$ , but not the syntactic features, are copied into the semantic features of  $\underline{\text{present}}_{\text{b}}$ . After chaining  $\underline{\text{present}}_{\text{b}}$  to  $\underline{\text{ita}}$ , the chaining table appears as shown in Figure 8.5.



Figure 8.5. Chaining Table after Chaining  $\overline{\mathrm{present}}_{\mathrm{b}}$  to  $\overline{\mathrm{ita}}$ 

The chaining algorithm now moves from present to June. Syntactic Conditions are satisfied, but Agreement isn't.

 $sc(it, June) = True$  $aqr(it, June) = False$ 

The chaining algorithm now moves from June to  $\phi$ . This time, all three rules, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule are satisfied.

sc(it,  $\phi$ ) = True  $aqr(it, \phi) = True$ rnr(it,  $\phi$ ) = True

Since all three rules are satisfied, E-nodes under it that agree with  $\phi$  chain to copies of  $\phi$ . it<sub>a</sub> is compared to  $\phi$ , and it is seen that they agree.

agr(it<sub>a</sub>,  $\phi$ ) = True

The chaining algorithm makes a new E-node copy of  $\underline{\phi}$ ,  $\underline{\phi}_b$ , and hangs it below  $\underline{\phi}$ . The chain\_link of  $\underline{\phi_b}$  is set to  $\underline{\text{ita}}$  and the semantic features of  $\frac{\text{it}_{a}}{\text{at}_{a}}$  are copied into the semantic features of  $\phi_{b}$ . After chaining  $\phi_{b}$  to  $\frac{\text{it}_{a}}{\text{at}_{a}}$ , the chaining table appears as shown in Figure 8.6.



Figure 8.6. Chaining Table After Chaining  $\phi_{\mathbf{b}}$  to  $\underline{\text{it}}_{\mathbf{a}}$ 

The chaining algorithm now moves to John and compares John to it. Syntactic Conditions hold, but Agreement does not.

 $sc(it, John) = True$  $aqr(it, John) = False$ 

Having exhausted all possible combinations with it, the chaining algorithm considers she.

The chaining algorithm tries comparing she to it, but Syntactic Conditions are not satisfied.

 $sc(she, it) = False$ 

The chaining algorithm moves from  $\underline{it}$  to she, but she can't chain to she, so the chaining algorithm moves to  $he$ . This time Syntactic Conditions are</u> satisfied, but Agreement isn't.

sc(she, he) = True  $agr(she, he) = False$ 

The chaining algorithm now moves from he to present where again Syntactic Conditions are satisfied, but Agreement isn't.

sc(she, present) = True agr(she, present) = False

The chaining algorithm moves from present to <u>June</u>. This time all three rules are satisfied.

sc(she, June) = True agr(she, June) = True rnr(she, June) = True As all three rules are satisfied, E-nodes under she that agree with June chain to copies of June. shea is compared to June, and it is seen that they agree.

 $aqr(she_a, June) = True$ 

The chaining algorithm makes a new E-node copy of <u>June,</u> June<sub>b</sub>, and hangs it below <u>June</u>. The chain\_link of  $\underline{\text{June}}_{\text{b}}$  is set to  $\underline{\text{she}}_{\text{a}}$  and the semantic features of  $\underline{\text{shea}}$  are copied into the semantic features of  $\underline{\text{June}}_{\text{b}}$ . After chaining  $\underline{\text{June}}_{\text{b}}$  to  $\underline{\text{shea}}$ , the chaining table appears as shown in Figure 8.7.



Figure 8.7. Chaining Table after Chaining  $\underline{\mathrm{June}}_{\mathrm{b}}$  to  $\underline{\mathrm{shea}}$ 

The chaining algorithm now moves from  $June to  $\phi$  and compares she to$ </u>  $\phi.$  All three rules are satisfied.

sc(she,  $\phi$ ) = True agr(she,  $\phi$ ) = True rnr(she,  $\phi$ ) = True

Copies of  $\phi$  are chain\_link'ed to E-nodes under she that agree with  $\phi$ . she<sub>a</sub> is compared to  $\phi$ , and it is seen that they agree.

agr(shea,  $\phi$ ) = True

A new E-node copy of  $\phi$ ,  $\phi$ <sub>C</sub>, is made and hung below  $\phi$ . The chain\_link of  $\phi_{\rm C}$  is set to  $\frac{\text{she}_a}{\text{ene}_a}$  and the semantic features of  $\frac{\text{she}_a}{\text{ene}_a}$  are copied into the semantic features of  $\phi_{\rm C}$ . After chaining  $\phi_{\rm C}$  to she<sub>a</sub>, the chaining table appears as shown in Figure  $8.\overline{8}$ .



Figure 8.8. Chaining Table after Chaining  $\phi_c$  to shea

The chaining algorithm now moves from  $\phi$  to John and compares she to John. It is seen that Syntactic Conditions are satisfied, but Agreement isn't.

sc(she, John) = True  $aqr(she, John) = False$ 

This completes the creation of chain\_link's to E-nodes under she. The chaining algorithm now considers he.

he is compared to it, but it is seen that Syntactic Conditions aren't satisfied.

 $sc(he, it) = False$ 

The chaining algorithm moves from it to she, but again, Syntactic Conditions aren't satisfied.

 $sc(he, she) = False$ 

The chaining algorithm moves from she to he, but he can't chain to he, so the chaining algorithm moves from he to present. This time Syntactic Conditions are satisfied, but Agreement isn't.

```
sc(he, present) = True
aqr(he, present) = False
```
The chaining algorithm moves from present to <u>June</u>, and similar results happen.

```
sc(he, June) = True
agr(he, June) = False
```
Next, the chaining algorithm moves from  $June to  $\phi$ , and, this time, all$ </u> three rules, Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule, are satisfied.

sc(he,  $\phi$ ) = True agr(he,  $\phi$ ) = True rnr(he,  $\phi$ ) = True

Copies of <u>he</u> are chain\_link'ed to E-nodes under  $\phi$  that agree with <u>he</u>. hea is compared to  $\phi$ , and it is seen that they agree.

$$
agr(he_a, \phi) = True
$$

A new E-node copy of  $\underline{\phi}, \, \underline{\phi_d},$  is made and hung below  $\underline{\phi}$ . The chain\_link of  $\phi_{\underline{d}}$  is set to  $\underline{he}_{\underline{a}}$  and the semantic features of  $\underline{he}_{\underline{a}}$  are copied into the semantic features of  $\phi_{\bf d}$ . After chaining  $\phi_{\bf d}$  to  $\underline{\text{hea}}$ , the chaining table appears as shown in Figure 8.9.



Figure 8.9. Chaining Table after Chaining  $\phi_{\bf d}$  to  ${\rm he}_{\bf d}$ 

The chaining algorithm now moves from  $\phi$  to John and he is compared to John. It is seen that all three rules are satisfied.

sc(he, John) = True agr(he, John) = True rnr(he, John) = True

So, copies of <u>he</u> are chain\_link'ed to E-nodes under John that agree with he. hea is compared to John, and it is seen that they agree.

 $aqr(he_a, John) = True$ 

A new E-node copy of <u>John</u>,  $\frac{\text{John}}{\text{John}}$ , is made and hung below <u>John</u>. The chain\_link of  $\underline{\text{John}}_b$  is set to  $\underline{\text{he}_a}$  and the semantic features of  $\underline{\text{he}_a}$  are copied into the semantic features of  $\underline{\rm{John}}_{\underline{b}}$ . After chaining  $\underline{\rm{John}}_{\underline{b}}$  to  $\underline{\rm{hea}}$ , the chaining table appears as shown in Figure 8.10.



Figure 8.10. Chaining Table after Chaining  $\underline{\mathrm{John}}_{\underline{\mathrm{b}}}$  to  $\underline{\mathrm{hea}}$ 

Having completed the processing of he, the chaining algorithm considers present. present is not a pronoun though, so the chaining algorithm moves on to June. Similarly, June is not a pronoun, so the chaining algorithm now considers  $\phi$ .

The chaining algorithm compares  $\phi$  to it, and it is seen that Syntactic Conditions don't hold.

sc( $\phi$ , it) = False

The chaining algorithm moves from  $\underline{it}$  to  $\underline{she}$ ,  $\underline{she}$  to  $\underline{he}$ ,  $\underline{he}$  to present, and present to June with little more success.

sc( $\phi$ , she) = False sc( $\phi$ , he) = False sc( $\phi$ , present) = False sc( $\phi$ , June) = False

The chaining algorithm moves from <u>June</u> to  $\phi$ , but  $\phi$  can't chain to  $\phi$ . So now, the chaining algorithm moves from  $\phi$  to <u>John</u>. This time, all three rules are satisfied.

```
sc(\phi, John) = True
agr(\phi, John) = Truernr(\phi, John) = True
```
Copies of <u>John</u> are chain\_link'ed to E-nodes under  $\phi$  that agree with <u>John</u>.  $\phi_{a}$  is compared to <u>John</u>, and it is seen that they agree.

 $aqr(\phi_a, John) = True$ 

Thus, new copy of <u>John, John<sub>c</sub></u> is made. <u>John<sub>c</sub></u> is chain\_link'ed to  $\phi_a$ . The semantic features of  $\phi_a$  are copied into John<sub>c</sub>. After chaining John<sub>c</sub> to  $\phi_a$ , the chaining table appears as shown in Figure 8.11.



Figure 8.11. Chaining Table after Chaining John<sub>c</sub> to  $\phi_{a}$ 

 $\phi_{\underline{b}}$  is compared to <u>John,</u> and it is seen that they don't agree.

 $agr(\phi_b, John) = False$ 

 $\phi_{\underline{b}}$  and <u>John</u> don't agree because when  $\phi_{\underline{b}}$  was chain\_link'ed to  $\underline{it}_{\underline{a}}$ , the semantic features of  $\underline{\text{it}}_{\underline{\text{a}}}$  were copied into the semantic features of  $\underline{\phi}_{\underline{\text{b}}}$ . Hence, the information that  $\underline{it_{a}}$  was inanimate was copied into  $\phi_{\underline{b}}$  preventing a ridiculous chain: <u>JohnX</u> is chained to  $\phi_{\underline{b}}$  is chained to  $\underline{it}_{\underline{a}}$ .  $\underline{\phi_{\underline{c}}}$ , which was chained to  $\underline{she}_{\underline{a}}$ , is compared to John, and it is seen that they don't agree.

 $agr(\phi_C, John) = False$ 

On the other hand,  $\phi_d$ , which was chained to  $\underline{he_a}$ , does agree with John.

agr( $\phi_d$ , John) = True

Thus, new copy of <u>John, John<sub>d</sub></u> is made.  $\underline{\text{John}}_d$  is chain\_link'ed to  $\underline{\phi_d}$ . The semantic features of  $\phi_d$  are copied into the semantic features of  $\underline{\text{John}}_d$ . After chaining  $\underline{\text{John}}_d$  to  $\underline{\phi_d}$ , the chaining table appears as shown in Figure 8.12.



Figure 8.12. Chaining Table after Chaining  $\underline{\mathrm{John}}_{\mathrm{d}}$  to  $\phi_{\mathrm{d}}$ 

Having completed chaining to  $\phi$ , the chaining algorithm moves to John. John is not a pronoun, so the chaining algorithm now stops as it has reached the end of the chaining table. This makes Figure 8.12, above, the finished chaining table.

## 9 Table Processor II

From Chapter 8 we know that the Table Processor module defines function chaining which takes as input a C-S-N tree and which returns as output the chaining table of the inputted C-S-N tree. In Chapter 8, we illustrated the kind of processing the Table Processor does by working through in detail a typical example. In this chapter, we will go into the particulars of the Table Processor algorithms. In Chapter 10, we'll look at some actual output.

The skeleton of the Table Processor module is shown below in Figure 9.1. The Table Processor module defines function chaining.

```
#table_proc.py
from node_proc import *
from parser import *
from secondary_uty import *
def chaining(nnodes: list[Node]) -> None:
```
### Figure 9.1. Skeleton of the Table Processor

Function chaining is the algorithm we described by example in Chapter 8. chaining takes as input a C-S-N tree and returns the chaining table of the inputted C-S-N tree. Function chaining is shown below in Figure 9.2.

```
def chaining(nnodes: list[Node]) -> None:
   init_table(nnodes)
    for n1 in reversed(nnodes):
        # For each N-node n1 that is a pronoun, call
        # procedure chaining_n.
        if n1.ftr[FeatureIndex.PNF] == Feature.PLUS:
            chaining_n(nnodes, n1)
```
### Figure 9.2. Function **chaining**

The first thing function chaining does is to call init\_table which initializes the chaining table as described in the previous chapter. Function init\_table is shown in Figure 9.3.

```
def init_table(nnodes: list[Node]) -> None:
    last = None
    for n in nnodes:
        n.col_link = new_e_node()
        n.col_link.ftr = n.ftr.copy()
        n.col_link.np_link = n
        n.col_link.sub = 'A'
        n.end_col_link = n.col_link
        n.pred_link = last
        if last is not None:
            last.succ_link = n
        last = n
```
### Figure 9.3. Function **init\_table**

Below each N-node is hung a new E-node with Feature's copied from the N-node and np\_link back to the N-node. The col\_link's and end\_col\_link's of the N-nodes are updated accordingly. pred\_link's and succ\_link's are set in init\_table using the thread\_link's which were established by the Parser. Finally, at the end of the procedure,

table, a variable global inside the Table Processor, has its left\_link and right\_link set to the first and last N-node.

For each N-node that is a pronoun, function chaining calls procedure chaining\_n. chaining\_n calls refl\_chaining or non\_refl\_chaining depending on whether or not the inputted N-node is reflexive or not. Function chaining\_n is shown below in Figure 9.4.

```
def chaining_n(nnodes: list[Node], n1: Node) -> None:
    if n1.ftr[FeatureIndex.RPF] == Feature.PLUS:
        # Inputted pronoun N-node n1 is reflexive.
        refl_chaining(n1)
    elif n1.ftr[FeatureIndex.RPF] == Feature.MINUS:
        # Inputted pronoun N-node n1 isn't reflexive.
        non refl chaining(nnodes, n1)
```
### Figure 9.4. Function **chaining\_n**

Function non\_refl\_chaining handles nonreflexive pronouns. Function non\_refl\_chaining is shown below in Figure 9.5.

```
def non_refl_chaining(nnodes: list[Node], n1: Node) -> None:
    for n2 in reversed(nnodes):
        if n2 := n1:
            chaining_n_to_n(n1, n2)
```
### Figure 9.5. Function **non\_refl\_chaining**

non\_refl\_chaining calls chaining\_n\_to\_n on the inputted N-node with each N-node in the chaining table except itself. This takes care of creating all chains to E-nodes lying under the inputted N-node.

Function refl\_chaining is very similar to non\_refl\_chaining and is shown below in Figure 9.6.

```
def refl_chaining(n1: Node) -> None:
   n2 = simplex_pred(n1)
    while n2 is not None:
        if n2 != n1:
            chaining_n_to_n(n1, n2)
        n2 = simplex_pred(n2)
```
#### Figure 9.6. Function **refl\_chaining**

Since the N-node inputted to refl\_chaining is reflexive, refl\_chaining only calls chaining\_n\_to\_n on the inputted N-node with each preceding N-node within the same simplex as the inputted N-node.

Function simplex\_pred, which is used by procedure refl\_chaining, simply returns the N-node that preceeds the inputted N-node in the same simplex. Function simplex\_pred is shown below in Figure 9.7.

```
def simplex_pred(n1: Node) -> Node:
   answer = n1while True:
        answer = answer.left_link
        if answer is None or answer.id == NodeId.N_NODE:
            return answer
```
### Figure 9.7. Function **simplex\_pred**

Function chaining\_n\_to\_n is called by procedures refl\_chaining and non\_refl\_chaining and is shown below in Figure 9.8.

```
def chaining_n_to_n(n1: Node, n2: Node) -> None:
   if not sc(n1, n2) or not agr(n1, n2) or not rnr(n1, n2):
        return
   old_end_col_link = n1.end_col_link
   e1 = n1while e1 != old_end_col_link:
        e1 = e1.col link
        if e1 is not None:
            chaining_e_to_n(e1, n2)
```
### Figure 9.8. Function **chaining\_n\_to\_n**

If Syntactic Conditions, Agreement, and the Reflexive Nonreflexive Rule hold, then procedure chaining\_e\_to\_n is called on each E-node lying underneath the first N-node.

Function chaining e to n, which is called by procedure chaining\_n\_to\_n, is shown below in Figure 9.9.

```
def chaining_e_to_n(e1: Node, n2: Node) -> None:
    if agr(e1, n2):
        new_chain(e1, n2)
```
#### Figure 9.9. Function **chaining\_e\_to\_n**

If the inputted E-node agrees with the inputted N-node, then a new chain is created from a copy of the inputted N-node to the inputted E-node by calling procedure new\_chain.

Function new\_chain, which is called by chaining\_e\_to\_n, is shown below in Figure 9.10.

```
def new chain(e1: Node, n2: Node) -> None:
   n = new_e node()
    n.np_1ink = n2n.chain_link = e1
    n.sub = chr(ord(n2.end\_col\_link.sub) + 1)# Replace n2 nonsyntactic QUESTION (?) features
    for i in range(N_FEATURES):
        if n2.ftr[i] == Feature.QUESTION and i != FeatureIndex.RPF:
            n. ftr[i] = e1. ftr[i]else:
            n.ftr[i] = n2.ftr[i]n2.end_col_link.col_link = n
   n2.end_col_link = n
```
### Figure 9.10. Function **new\_chain**

Function new\_chain creates a copy of the inputted N-node and chain\_link's it to the inputted E-node. Semantic Feature's are copied from the inputted E-node to the copy of the inputted N-node.

## 10 Table Processor III

The last two chapters have been devoted to describing the Table Processor. It is time for some real examples: The first example we present in this chapter was produced using procedure view node str of the Node Processor along with some intermittent write statements indicating when LIP are entering and exiting some of the more important routines and some or their results. We start off with the example first presented in Chapter 8.

(10.1) John wants to give June a present, but he isn't sure she'll like it.



Processing (10.1) with some intermediate output gives the listing shown below. This is somewhat verbose, but later examples will be cleaner and shorter, though less detailed output.





chaining init\_table



```
init_table: exiting
chaining_n(it)
    non_refl_chaining(it)
        chaining_n_to_n(it, she)
            sc(it, she) = Trueagr(it, she) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(it, he)
            sc(it, he) = Trueagr(it, he) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(it, present)
            sc(it, present) = True
            agr(it, present) = True
            rnr(it, present) = True
            chaining_e_to_n(ita, present)
                agr(ita, present) = Truenew_chain(ita, present)
                    new_chain: create presentb
                    new_chain: create present<sub>h</sub>^ita
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(it, June)
    sc(it, June) = True
    agr(it, June) = False
chaining_n_to_n: exiting
chaining_n_to_n(it, \phi)
    sc(it, \phi) = True
    agr(it, \phi) = True
    rnr(it, \phi) = True
    chaining_e_to_n(ita, \phi)
         agr(it<sub>a</sub>, \phi) = True
         new_chain(ita, \phi)
              new_chain: create \phi_{\rm b}new_chain: create \phi_b<sup>2</sup>ita
```


```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(it, John)
            sc(it, John) = True
            agr(it, John) = Falsechaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(she)
    non_refl_chaining(she)
        chaining_n_to_n(she, it)
            sc(she, it) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(she, he)
            sc(she, he) = True
            agr(she, he) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(she, present)
            sc(she, present) = True
            agr(she, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, June)
            sc(she, June) = True
            agr(she, June) = True
            rnr(she, June) = True
            chaining_e_to_n(shea, June)
                agr(she<sub>a</sub>, June) = Truenew_chain(shea, June)
                    new_chain: create Juneb
                    new_chain: create June<sub>h</sub>^shea
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(she, \phi)
    sc(she, \phi) = True
    agr(she, \phi) = True
    rnr(she, \phi) = True
    chaining_e_to_n(shea, \phi)
         agr(shea, \phi) = True
         new_chain(shea, \phi)
             new_chain: create \phi_Cnew_chain: create \phi_c<sup>o</sup>shea
```


```
new_chain: exiting
             chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(she, John)
             sc(she, John) = True
             agr(she, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(he)
    non_refl_chaining(he)
        chaining_n_to_n(he, it)
             sc(he, it) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, she)
             sc(he, she) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, present)
             sc(he, present) = True
             agr(he, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, June)
             sc(he, June) = True
             aqr(he, June) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, \phi)
             sc(he, \phi) = True
             agr(he, \phi) = True
             rnr(he, \phi) = True
             chaining_e_to_n(hea, \phi)
                 agr(he<sub>a</sub>, \phi) = True
                 new_chain(hea, \phi)
                     new_chain: create \phi_dnew_chain: create \phi_d<sup>^</sup>hea
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_n\_to\_n: exiting chaining\_n\_to\_n(he, John) sc(he, John) = True agr(he, John) = True rnr(he, John) = True chaining\_e\_to\_n(hea, John) agr(hea, John) = True new\_chain(hea, John) new\_chain: create Johnb  $new\_chain: create John<sub>b</sub><sup>~</sup>he<sub>a</sub>$ 



```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non refl chaining(\phi)
         chaining_n_to_n(\phi, it)
             sc(\phi, it) = Falsechaining_n_to_n: exiting
         chaining_n_to_n(\phi, she)
             sc(\phi, she) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, he)
             sc(\phi, he) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, present)
             sc(\phi, present) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, June)
             sc(\phi, June) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = Truernr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phia, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create Johnc
                      new_chain: create John<sub>c</sub>^\phi<sub>a</sub>
```


```
new_chain: exiting
chaining_e_to_n: exiting
chaining_e_to_n(\phi_b, John)
    agr(\phi<sub>b</sub>, John) = False
chaining_e_to_n: exiting
chaining_e_to_n(\phi_C, John)
    agr(\phi_C, John) = Falsechaining_e_to_n: exiting
chaining_e_to_n(\phi_d, John)
     agr(\phi_d, John) = True
     new_chain(\phi_{\rm d}, John)
         new_chain: create Johnd
         new_chain: create John<sub>d</sub><sup>^</sup>\phi<sub>d</sub>
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_n\_to\_n: exiting non\_refl\_chaining: exiting chaining\_n: exiting chaining: exiting



The previous example should give enough details away to satisfy the reader's curiosity, but the form of the previous example is rather burdensome. From now on, we'll keep to a more concise, if less detailed, output. To indicate a chain\_link between two E-nodes, we use the symbol ˆ.

Below are some more examples.

(10.2) Janet saw herself.







```
init_table: exiting
chaining_n(herself)
    refl_chaining(herself)
        simplex_pred(herself)
        simplex_pred: Janet
        chaining_n_to_n(herself, Janet)
            sc(herself, Janet) = True
            agr(herself, Janet) = True
            rnr(herself, Janet) = True
            chaining_e_to_n(herselfa, Janet)
                agr(herselfa, Janet) = Truenew_chain(herselfa, Janet)
                    new_chain: create Janetb
                    new\_chain: create Janet_b^*herself_a
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            simplex_pred(Janet)
            simplex_pred:
        refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.3) Janet saw her.







```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
        chaining_n_to_n(her, Janet)
            sc(her, Janet) = True
            agr(her, Janet) = True
            rnr(her, Janet) = True
            chaining_e_to_n(hera, Janet)
                agr(hera, Janet) = True
                new_chain(hera, Janet)
                    new_chain: create Janetb
                    new\_chain: create Janet_b^* here
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_n\_to\_n: exiting non\_refl\_chaining: exiting chaining\_n: exiting chaining: exiting



(10.4) \*Janet saw himself.







```
init_table: exiting
   chaining_n(himself)
       refl_chaining(himself)
            simplex_pred(himself)
            simplex_pred: Janet
            chaining_n_to_n(himself, Janet)
                sc(himself, Janet) = True
                aqr(himself, Janet) = Falsechaining_n_to_n: exiting
            simplex_pred(Janet)
            simplex_pred:
       refl_chaining: exiting
   chaining_n: exiting
chaining: exiting
```


Examples  $(10.5)-(10.11)$  are from Lees and Klima  $[22]$ .

(10.5) The men threw a smokescreen around themselves.



the men threw a smokescreen around themselves





```
init_table: exiting
chaining_n(themselves)
    refl_chaining(themselves)
        simplex_pred(themselves)
        simplex_pred: smokescreen
        chaining_n_to_n(themselves, smokescreen)
            sc(themselves, smokescreen) = True
            agr(themselves, smokescreen) = False
        chaining_n_to_n: exiting
        simplex_pred(smokescreen)
        simplex_pred: men
        chaining_n_to_n(themselves, men)
            sc(themselves, men) = True
            agr(themselves, men) = True
            rnr(themselves, men) = True
            chaining_e_to_n(themselvesa, men)
                aqr(themselves<sub>a</sub>, men) = True
                new_chain(themselvesa, men)
                    new_chain: create menb
                    new_chain: create menb^themselvesa
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            simplex_pred(men)
            simplex_pred:
        refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.6) The men found a smokescreen around them.



the men threw a smokescreen around them





```
init_table: exiting
   chaining_n(them)
        non_refl_chaining(them)
            chaining_n_to_n(them, smokescreen)
                sc(them, smokescreen) = True
                agr(them, smokescreen) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(them, men)
                sc(them, men) = Trueagr(them, men) = True
                rnr(them, men) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.7) The men found a smokescreen to be around them.







```
init_table: exiting
chaining_n(them)
    non_refl_chaining(them)
        chaining_n_to_n(them, smokescreen)
            sc(them, smokescreen) = True
            agr(them, smokescreen) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, men)
            sc(them, men) = True
            agr(them, men) = True
            rnr(them, men) = True
            chaining_e_to_n(thema, men)
                agr(hem_a, men) = Truenew_chain(thema, men)
                     new_chain: create menb
                     new_chain: create men<sub>b</sub>^thema
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.8) The men found a smokescreen and it was around them.







```
init_table: exiting
chaining_n(them)
    non_refl_chaining(them)
        chaining_n_to_n(them, it)
             sc(them, it) = Trueagr(them, it) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, smokescreen)
             sc(them, smokescreen) = True
             agr(them, smokescreen) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(them, men)
             sc(them, men) = True
             agr(them, men) = True
             rnr(them, men) = True
             chaining_e_to_n(thema, men)
                 agr(them<sub>a</sub>, men) = True
                 new_chain(thema, men)
                     new_chain: create menb
                      new\_chain: create men<sub>h</sub><sup>T</sup> them<sub>a</sub>
```


```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(it)
    non refl chaining(it)
        chaining_n_to_n(it, them)
            sc(it, them) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(it, smokescreen)
            sc(it, smokescreen) = True
            agr(it, smokescreen) = True
            rnr(it, smokescreen) = True
            chaining_e_to_n(ita, smokescreen)
                agr(it<sub>a</sub>, smokescreen) = True
                new_chain(ita, smokescreen)
                    new_chain: create smokescreenb
                    new_chain: create smokescreenboita
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            chaining_n_to_n(it, men)
                sc(it, men) = True
                agr(it, men) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
   chaining_n: exiting
chaining: exiting
```


(10.9) I told John to protect himself.







```
init_table: exiting
chaining_n(himself)
     refl_chaining(himself)
           simplex_pred(himself)
           simplex_pred: \phichaining_n_to_n(himself, \phi)
                 sc(himself, \phi) = True
                agr(himself, \phi) = True
                rnr(himself, \phi) = True
                chaining_e_to_n(himself<sub>a</sub>, \phi)
                      agr(himself<sub>a</sub>, \phi) = True
                      new_chain(himself<sub>a</sub>, \phi)
                           new_chain: create \phi<sub>b</sub>
                           new_chain: create \stackrel{\sim}{\phi_{\rm b}}<sup>^</sup>himself<sub>a</sub>
```


```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         simplex_pred(\phi)simplex_pred:
    refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non_refl\_chaining(\phi)chaining n_to_n(\phi, \text{himself})sc(\phi, himself) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = True
             rnr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phi_a, John) = Truenew_chain(\phi_a, John)
                      new_chain: create Johnb
                      new_chain: create John_b^{\hat{}}\phi_a
```


```
new_chain: exiting
chaining_e_to_n: exiting
chaining_e_to_n(\phib, John)
     agr(\phi<sub>b</sub>, John) = True
     new_chain(\phi<sub>b</sub>, John)
          new_chain: create Johnc
          new_chain: create John<sub>c</sub>^\phi<sub>b</sub>
```


```
new_chain: exiting
     chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(\phi, I)
     sc(\phi, I) = Trueagr(\phi, I) = Truernr(\phi, I) = Truechaining_e_to_n(\phi_a, I)agr(\phi_a, I) = Truenew\_chain(\phi_a, I)new_chain: create I<sub>b</sub>
                new_chain: create \text{I}_{\text{b}}\hat{\;} \phi_\text{a}
```


```
new_chain: exiting
                chaining_e_to_n: exiting
                chaining_e_to_n(\phi_b, I)agr(\phi_b, I) = Falsechaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, himself)
                sc(I, himself) = Falsechaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = Falsechaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.10) I told John to protect me.





```
chaining
```

```
init_table
```


```
init_table: exiting
chaining_n(me)
    non_refl_chaining(me)
        chaining_n_to_n(me, \phi)
            sc(me, \phi) = True
            agr(me, \phi) = True
            rnr(me, \phi) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(me, John)
            sc(me, John) = True
            agr(me, John) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(me, I)
            sc(me, I) = Trueagr(me, I) = True
            rnr(me, I) = True
            chaining_e_to_n(mea, I)
                 aqr(me_a, I) = Truenew_chain(mea, I)
                     new_chain: create Ib
                     new_chain: create I<sub>b</sub>^mea
```


```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non refl chaining(\phi)
         chaining_n_to_n(\phi, me)
             sc(\phi, me) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = Truernr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  agr(\phia, John) = True
                  new_chain(\phi_a, John)
                      new_chain: create John<sub>h</sub>
                      new_chain: create Johnh^{\hat{}}\phi_a
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining n_to_n(\phi, I)sc(\phi, I) = Trueagr(\phi, I) = Truernr(\phi, I) = True
    chaining_e_to_n(\phi_a, I)
         agr(\phi_a, I) = Truenew\_chain(\phi_a, I)new_chain: create Ic
             new_chain: create I_c^{\hat{}}\phi_a
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, me)
                sc(I, me) = Falsechaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = Falsechaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


(10.11) I told John to protect myself.







```
init_table: exiting
chaining_n(myself)
     refl_chaining(myself)
          simplex_pred(myself)
          simplex_pred: \phichaining_n_to_n(myself, \phi)
                sc(myself, \phi) = True
                agr(myself, \phi) = True
                rnr(myself, \phi) = True
                chaining_e_to_n(myself<sub>a</sub>, \phi)
                     agr(myself<sub>a</sub>, \phi) = True
                     new_chain(myself<sub>a</sub>, \phi)
                          new\_chain: \text{ create } \phi_bnew_chain: create \phi_{\rm b}^{\nu}<sup>o</sup>myself<sub>a</sub>
```


```
new_chain: exiting
             chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         simplex_pred(\phi)simplex_pred:
    refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non_refl\_changing(\phi)chaining_n_to_n(\phi, myself)
             sc(\phi, myself) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
             sc(\phi, John) = True
             agr(\phi, John) = Truernr(\phi, John) = True
             chaining_e_to_n(\phi_a, John)
                  aqr(\phi_a, John) = Truenew_chain(\phi_a, John)
                      new_chain: create John<sub>h</sub>
                      new_chain: create Johnh^{\hat{}}\phi_a
```


```
new_chain: exiting
    chaining_e_to_n: exiting
    chaining_e_to_n(\phib, John)
         agr(\phi<sub>b</sub>, John) = False
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(\phi, I)
    \text{sc}(\phi, I) = \text{True}agr(\phi, I) = Truernr(\phi, I) = True
    chaining e_to_n(\phi_a, I)agr(\phi_a, I) = Truenew_chain(\phi_a, I)
              new_chain: create Ib
              new_chain: create I_b^{\hat{}}\phi_a
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_e\_to\_n $(\phi_b, I)$  $agr(\phi_h, I) = True$  $new\_chain(\phi_b, I)$ new\_chain: create Ic new\_chain: create  $I_c^{\dagger} \phi_b$ 



```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I)
        non_refl_chaining(I)
            chaining_n_to_n(I, myself)
                sc(I, myself) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, \phi)
                sc(I, \phi) = False
            chaining_n_to_n: exiting
            chaining_n_to_n(I, John)
                sc(I, John) = Falsechaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```


## 11 Table Interpreter

The Table Interpreter module defines function interpret and has the form shown in Figure 11.1.

```
#table_interp;
from globals import *
def interpret(nnodes: list[Node]) -> list[list[list[Node]]]:
```
## Figure 11.1. Skeleton of the Table Interpreter

Basically, after the chaining table is created, a number of chains are implicitly defined by the chaining table and it is the job of the Table Interpreter to mesh these chains back into copies of the system tree, returning all trees defined by legitimate interpretations.

A nonpronominal E-node with the E-nodes that are traced by walking down chain\_link's until nil chain\_link is reached constitute a **chain**. A set of chains defined by the chaining table which cover all the pronominal N-nodes and do not intersect constitute a legitimate interpretation.

Take the table given in Figure 11.2 as an example.



## Figure 11.2. Typical Chaining Table

The chains present in Figure 11.2 are exactly (11.3)-(11.10) given below. Note that no chain begins with a pronoun. (Here, we are staying with the convention of Chapter 10 where a "<sup>^"</sup> symbol indicates a chain\_link).

(11.3) Johna

- $(11.4)$  John<sub>b</sub> hea
- $(11.5)$  John<sub>c</sub>  $\phi_a$
- $(11.6)$  John<sub>d</sub><sup> $\hat{\phi}_d$ </sup>hea

(11.7) Junea

 $(11.8)$  June<sub>b</sub>^shea  $(11.9)$  present<sub>a</sub>  $(11.10)$  present<sub>b</sub><sup>t</sup>ita

The only interpretation derivable from Figure 11.2 is (11.11).

 $(11.11)$  John<sub>d</sub><sup>2</sup> $\phi$ <sub>d</sub>  $June<sub>b</sub>^shea$  $\hat{\ }$ shea present<sub>b</sub> $\hat{\ }$ ita

Exactly how chaining information from a table to system parse tree will have to be system dependent, but we can imagine that noun phrases in a system parse tree are some kind of list elements which have linked to them, among other things, lists corresponding to their semantics. It is up to the Table Interpreter to set any chain\_link's inside the semantics of the noun phrases of the system parse tree. The Semantic Processor module of the system should then be powerful enough to be able to handle the kind of coordination that chain\_link's imply.

This strategy has a number of possibilities that simple methods of coreferencing are just not able to handle. Consider sentence (11.12) for example.

(11.12) Jack's house burned down, but he rebuilt it.

We can't really say that it **corefers** with Jack's house as Jack's house is some object that existed in the past and has stopped existing while it refers to some new object. This does not mean that it cannot chain from Jack's house, however, and indeed it should. The information the Semantic Processor module needs to give meaning to it is contained in Jack's house, and so there must be a chain\_link from Jack's house to it in order for the Semantic Processor to give meaning to it.

A similar result holds for quantifiers. We see that (11.13) is not equivalent to (11.14).

(11.13) Every connoisseur loves his wine and cheese.  $\neq$ 

(11.14) Every connoisseur loves every connoisseur's wine and cheese.

Quite clearly, his cannot be replaced by every connoisseur and preserve the meaning of the sentence. Instead, (11.13) has more the meaning given by  $(11.15).$ 

 $(11.15)$  (For all x: x is a connoisseur)(x loves x's wine and cheese.)

A number of other examples are pointed out by Bresnan [\[3\]](#page-131-0).

 $(11.16)$  All Italians think they are handsome.  $\neq$ 

(11.17) All Italians think all Italians are handsome.

(11.18) Every Italian thinks he is handsome.  $\neq$ 

(11.19) Every Italian thinks every Italian is handsome.

(11.20) Any Italian would die for his mother.  $\neq$ 

(11.21) Any Italian would die for and Italian's mother.

- $(11.22)$  Every Italian thinks that he alone is handsome.  $\neq$
- (11.23) \*Every Italian thinks that every Italian alone is handsome.
- (11.24) One girl claimed that she herself could read Homer.  $\neq$
- (11.25) \*One girl claimed that one girl herself could read Homer.

It appears that the proper interpretation for a pronoun chained to quantified noun phrase within the scope of quantification is for the pronoun to act as a bound variable.

When the pronoun is outside the scope of quantification, it is a different story. Consider  $(11.26)$  and  $(11.27)$  from Evans [\[6\]](#page-131-1).

(11.26) John owns some sheep and Harry vaccinates them.

(11.27) Mary danced with many boys and they found her interesting.

This time the pronouns are chaining to quantified noun phrases, but do not themselves lie within the scope of quantification. Instead, they appear to refer to the range of the quantification.

Similar results hold for (11.28)-(11.31) from Sidner [\[32\]](#page-132-1).

- (11.28) John lost a pen yesterday and Bill found one today.
- (11.29) John claimed to have found the solution to the problem, but Bill was sure he had found it.
- (11.30) John wants to catch a fish and eat it for supper.
- (11.31) No one would put the blame on himself.

The problems mentioned above are all rather tricky, but viewing them from the vantage point of chaining sheds more light on them than viewing them through some kind of coreference. The moral of the story seems to be that anaphora is not coreference.

Using the Table Interpreter now, we present some more examples.

(11.32) Sue told Sandy about herself.





```
init_table: exiting
chaining_n(herself)
    refl_chaining(herself)
        simplex_pred(herself)
        simplex_pred: Sandy
        chaining_n_to_n(herself, Sandy)
            sc(herself, Sandy) = True
            agr(herself, Sandy) = True
            rnr(herself, Sandy) = True
            chaining_e_to_n(herselfa, Sandy)
                agr(herselfa, Sandy) = Truenew_chain(herselfa, Sandy)
                    new_chain: create Sandyb
                    new_chain: create Sandyb^herselfa
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
simplex_pred(Sandy)
simplex_pred: Sue
chaining_n_to_n(herself, Sue)
    sc(herself, Sue) = True
    agr(herself, Sue) = True
    rnr(herself, Sue) = True
    chaining_e_to_n(herselfa, Sue)
        agr(herself<sub>a</sub>, Sue) = Truenew_chain(herselfa, Sue)
            new_chain: create Sueb
            new_chain: create Sue<sub>b</sub>^herselfa
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            simplex_pred(Sue)
            simplex_pred:
        refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```




(11.33) \*Jill kept talking about himself.





```
init_table: exiting
   chaining_n(himself)
        refl_chaining(himself)
            simplex_pred(himself)
            simplex_pred: Jill
            chaining_n_to_n(himself, Jill)
                sc(himself, Jill) = True
                agr(himself, Jill) = False
            chaining_n_to_n: exiting
            simplex_pred(Jill)
            simplex_pred:
        refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```




(11.34) Does Jack's making a pig of himself bother Bill?





```
init_table: exiting
chaining_n(himself)
    refl_chaining(himself)
        simplex_pred(himself)
        simplex_pred: pig
        chaining_n_to_n(himself, pig)
            sc(himself, pig) = True
            agr(himself, pig) = True
            rnr(himself, pig) = True
            chaining_e_to_n(himselfa, pig)
                agr(himselfa, pig) = Truenew_chain(himselfa, pig)
                    new_chain: create pigh
                    new_chain: create pig<sub>b</sub>^himselfa
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            simplex_pred(pig)
            simplex_pred: Jack's
            chaining_n_to_n(himself, Jack's)
                sc(himself, Jack's) = True
                agr(himself, Jack's) = True
                rnr(himself, Jack's) = False
            chaining_n_to_n: exiting
            simplex_pred(Jack's)
            simplex_pred:
        refl_chaining: exiting
   chaining_n: exiting
chaining: exiting
```




Features								
	<b>PNF</b>	FPF	<b>SPF</b>	TPF	PLF	<b>GNF</b>	<b>ANF</b>	<b>RPF</b>
John								
Φ	$^{+}$	?	っ	?	7	ာ	?	
June				$^+$		$^{+}$	$^{+}$	
present						?		
he	$\overline{+}$							
she	$^{+}$						$^{+}$	
it						っ		

(11.35) John wants to give June a present, but he is afraid she won't like it.

```
chaining
   init_table
```


```
init_table: exiting
chaining_n(it)
    non_refl_chaining(it)
        chaining_n_to_n(it, she)
            sc(it, she) = Trueagr(it, she) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(it, he)
            sc(it, he) = True
            agr(it, he) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(it, present)
             sc(it, present) = True
            agr(it, present) = True
            rnr(it, present) = True
            chaining_e_to_n(ita, present)
                 agr(it<sub>a</sub>, present) = Truenew_chain(ita, present)
                     new_chain: create presentb
                     new_chain: create presentb<sup>^ita</sup>
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(it, June)
    sc(it, June) = True
    agr(it, June) = False
chaining_n_to_n: exiting
chaining_n_to_n(it, \phi)
    sc(it, \phi) = True
    agr(it, \phi) = Truernr(it, \phi) = True
    chaining_e_to_n(ita, \phi)
         agr(ita, \phi) = Truenew_chain(ita, \phi)
             new_chain: create \phi_{\rm b}new_chain: create \phi_{\rm b}<sup>2</sup>ita
```


```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(it, John)
            sc(it, John) = True
            agr(it, John) = Falsechaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(she)
    non_refl_chaining(she)
        chaining_n_to_n(she, it)
            sc(she, it) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(she, he)
            sc(she, he) = True
            agr(she, he) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(she, present)
            sc(she, present) = True
            agr(she, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(she, June)
            sc(she, June) = True
            agr(she, June) = True
            rnr(she, June) = True
            chaining_e_to_n(shea, June)
                agr(she<sub>a</sub>, June) = Truenew_chain(shea, June)
                    new_chain: create Juneb
                    new_chain: create June<sub>h</sub>^shea
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(she, \phi)
    sc(she, \phi) = True
    agr(she, \phi) = True
    rnr(she, \phi) = True
    chaining_e_to_n(shea, \phi)
         agr(shea, \phi) = True
         new_chain(shea, \phi)
             new_chain: create \phi_Cnew_chain: create \phi_c<sup>o</sup>shea
```

```
new_chain: exiting
            chaining_e_to_n: exiting
        chaining_n_to_n: exiting
        chaining_n_to_n(she, John)
            sc(she, John) = True
            agr(she, John) = False
        chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(he)
    non_refl_chaining(he)
        chaining_n_to_n(he, it)
            sc(he, it) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, she)
            sc(he, she) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, present)
            sc(he, present) = True
            agr(he, present) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, June)
            sc(he, June) = True
            aqr(he, June) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(he, \phi)
            sc(he, \phi) = True
            agr(he, \phi) = True
            rnr(he, \phi) = True
            chaining_e_to_n(hea, \phi)
                 agr(hea, \phi) = True
                 new_chain(hea, \phi)
                     new_chain: create \phi_dnew_chain: create \phi_d<sup>^</sup>hea
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(he, John)
    sc(he, John) = True
    agr(he, John) = True
    rnr(he, John) = True
    chaining_e_to_n(hea, John)
        agr(hea, John) = True
        new_chain(hea, John)
             new_chain: create Johnb
             new\_chain: create John<sub>b</sub><sup>~</sup>he<sub>a</sub>
```


```
new_chain: exiting
              chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non refl chaining(\phi)
         chaining_n_to_n(\phi, it)
             \text{sc}(\phi, \text{ it}) = \text{False}chaining_n_to_n: exiting
         chaining_n_to_n(\phi, she)
              sc(\phi, she) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, he)
              sc(\phi, he) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, present)
              sc(\phi, present) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, June)
              sc(\phi, June) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, John)
              sc(\phi, John) = True
              agr(\phi, John) = Truernr(\phi, John) = True
              chaining_e_to_n(\phi_a, John)
                  agr(\phia, John) = True
                  new_chain(\phi_a, John)
                       new_chain: create Johnc
                       new_chain: create John<sub>c</sub>^\phi<sub>a</sub>
```


```
new_chain: exiting
chaining_e_to_n: exiting
chaining_e_to_n(\phi_b, John)
    agr(\phi<sub>b</sub>, John) = False
chaining_e_to_n: exiting
chaining_e_to_n(\phi_C, John)
    agr(\phi_C, John) = False
chaining_e_to_n: exiting
chaining_e_to_n(\phi_d, John)
    agr(\phi_d, John) = True
    new\_chain(\phi_d, John)new_chain: create Johnd
         new_chain: create John<sub>d</sub>^\phi_d
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_n\_to\_n: exiting non\_refl\_chaining: exiting chaining\_n: exiting chaining: exiting





(11.36) Ernie doesn't like Bernie, because he is such an asshole.

<b>Features</b>												
<b>SPF</b> <b>GNF</b> <b>PNF</b> TPF FPF <b>ANF</b> PLF <b>RPF</b>												
Ernie												
Bernie												
he												
asshole						$\sim$						

chaining init\_table



```
init_table: exiting
chaining_n(he)
    non_refl_chaining(he)
        chaining_n_to_n(he, asshole)
             sc(he, asshole) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(he, Bernie)
             sc(he, Bernie) = True
             agr(he, Bernie) = True
             rnr(he, Bernie) = True
             chaining_e_to_n(hea, Bernie)
                 agr(he<sub>a</sub>, Bernie) = Truenew_chain(hea, Bernie)
                     new_chain: create Bernieb
                     new_chain: create Bernie<sub>b</sub>^hea
```


```
new_chain: exiting
    chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(he, Ernie)
    sc(he, Ernie) = True
    agr(he, Ernie) = True
    rnr(he, Ernie) = True
    chaining_e_to_n(hea, Ernie)
        agr(hea, Ernie) = True
        new_chain(hea, Ernie)
            new_chain: create Ernieb
            new_chain: create Ernie<sub>b</sub>^hea
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
       non_refl_chaining: exiting
   chaining_n: exiting
chaining: exiting
```




# 12 Genitives

Very little modification to what has been said so far is necessary to implement attributive possessive pronouns. Recall that the attributive possessive pronouns are those pronouns listed in (12.1).

<span id="page-114-0"></span>(12.1) my, our, your, her, his, its, their

Examining sentences like (12.2)-(12.5) reveals that reflexive pronouns don't chain to genitives within the same simplex. On the other hand, nonreflexive pronouns can.

- (12.2) Mary's father killed himself.
- (12.3) \*Mary's father killed him.
- (12.4) \*Mary's father killed herself.
- (12.5) Mary's father killed her.

The same conclusions also hold for **of-genitives**. Compare sentences (12.6)- $(12.9)$  to  $(12.2)-(12.5)$ .

- (12.6) The father of Mary killed himself.
- (12.7) \*The father of Mary killed him.
- (12.8) \*The father of Mary killed herself.
- (12.9) The father of Mary killed her.

The easiest way to handle genitives, apparently, is to introduce a new Feature, GEN, for genitive and to modify the Reflexive Nonreflexive Rule to handle genitives. The new form of the Reflexive Nonreflexive Rule is shown below in Figure 12.10.

```
function rnr(n1,n2:NodePointer):boolean;
    {Reflexive Nonreflexive Rule}
    var ftr1, ftr2: features;
begin
    n1:=n1ˆ.np_link;
    n2:=n2ˆ.np_link;
    ftr1:=n1ˆ.ftr;
    ftr2:=n2ˆ.ftr:
    if ftr2[GEN]=PLUS then rnr:=false
    else case ftr1[RPF] of
        PLUS: rnr:=(n1ˆ.up_link=n2ˆ.up_link)
                        and (ftr1[GEN] ==MINUS);
        QUESTION: {doesn't occur};
        MINUS: rnr:=(n1^.up_link<>n2^.up_link)
                        or (ftr1[GEN]<>MINUS);
    end;
end;
```
#### Figure 12.10. New Reflexive Nonreflexive Rule

The examples following illustrate the interpretation of attributive possessive pronouns and pronouns in the context of genitives.

(12.11) Mary's mother cooks only for herself.

Features											
PNF <b>SPF</b> <b>ANF</b> <b>RPF</b> FPF TPF PLF <b>GNF</b> <b>GEN</b>											
Mary's											
mother											
herself											

(12.12) Mary's mother cooks only for her.



chaining init\_table



```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
        chaining_n_to_n(her, mother)
            sc(her, mother) = True
            agr(her, mother) = True
            rnr(her, mother) = True
            chaining_e_to_n(hera, mother)
                 agr(her_a, mother) = Truenew_chain(hera, mother)
                     new_chain: create motherb
                     new_chain: create mother<sub>b</sub>^hera
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            chaining_n_to_n(her, Mary's)
                sc(her, Mary's) = True
                agr(her, Mary's) = True
               rnr(her, Mary's) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```




(12.13) Mary's mother cooks only for her mother.



chaining init\_table



```
init_table: exiting
chaining_n(her)
    non_refl_chaining(her)
         chaining_n_to_n(her, mother2)
              sc(her, mother<sub>2</sub>) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(her, mother1)
              sc(her, mother<sub>1</sub>) = True
              agr(her, mother<sub>1</sub>) = Truernr(her, mother_1) = Truechaining_e_to_n(hera, mother<sub>1</sub>)
                  agr(her_a, mother_1) = Truenew_chain(hera, mother1)
                       new_chain: create mother<sub>1b</sub>
                       new_chain: create mother<sub>1b</sub>^hera
```


```
new_chain: exiting
                chaining_e_to_n: exiting
            chaining_n_to_n: exiting
            chaining_n_to_n(her, Mary's)
                sc(her, Mary's) = True
                agr(her, Mary's) = True
                rnr(her, Mary's) = False
            chaining_n_to_n: exiting
        non_refl_chaining: exiting
   chaining_n: exiting
chaining: exiting
```




## <span id="page-118-0"></span>13 Focusing

Extrasentential anaphora and ellipsis is possible through the maintenance of a focus of conversation. This maintenance is known as focusing and has been described at length by Grosz [\[10\]](#page-131-0) and Sidner [\[32\]](#page-132-0). By focus of conversation, we mean the common view of the participants of conversation of what their conversation is about. Focusing is useful because it allows the participants of conversation to avoid redundant repetition of old material. Assuming focusing is desirable in a computer natural language system, how do we implement it?

Grosz has examined task dialogues in which an expert helps an apprentice to assemble a mechanical air compressor. She finds it convenient to represent the focus of conversation as a set of overlapping focus spaces, where each focus space is a collection of objects. One focus space is active and the others are open. When a focus space is no longer needed, it is closed. One of Grosz's assumptions is that goals and subgoals are definable and recognizable in a task dialogue system with the consequence that in any conversation there is an open focus space hierarchy with the active focus space at the bottom of the hierarchy.

Sidner has approached the problem of focusing from a different perspective by analyzing monologues. For Sidner, focus is kept track of through a discourse focus, actor focus, potential discourse foci, potential actor foci, discourse focus stack, and actor focus stack. Sidner's work, which came after Grosz's, is very commendable for the algorithms she presents, although most of these are fairly sketchy.

In our approach, we will treat the focus of conversation as a collection of nonpronominal N-nodes. Among the N-nodes that we would ordinairly expect to always be in focus are the I and you of a conversation. To get a handle on the focused N-nodes, we dominate them by an S-node just as if they all had occurred in one simplex. So, for example, if  $\underline{\mathrm{I}_{0}}$  and  $\underline{\mathrm{vou}_{0}}$  are the nonpronominal N-nodes currently in focus, then the current focus representation is given by a structure like Figure 13.1.



Figure 13.1. Typical Focus Representation

When it comes time to analyze a sentence, the current focus representation is attached to the C-S-N parse tree of the sentence via a C-node which dominates them both. This makes the focused N-nodes available to the N-nodes of the C-S-N parse tree for chaining.

As an example, suppose  $I_0$  and  $\underline{you_0}$  are in focus and the current input sentence is (13.2) from Grinder [\[8\]](#page-131-1).

(13.2) It was difficult to sketch myself.



The C-S-N parse tree of (13.2) will have a form like that indicated below in Figure 13.3.



Figure 13.3. C-S-N Parse Tree of (13.2)

As the parse tree now stands in Figure 2, myself may chain to  $\phi$ , but  $\phi$  does not have an N-node to chain to. Thus, there are no legitimate interpretations without focusing. With focusing, the Parser attaches the current focus representation containing  $I_0$  and  $\underline{\text{you}}_0$  to the C-S-N parse tree by a C-node obtaining the new C-S-N parse tree shown in Figure 13.4.



Figure 13.4. C-S-N Parse Tree with Focusing

Now, <u>myself</u> can chain from  $\underline{\phi}$  and  $\underline{\phi}$  can chain from  $\underline{I_0}$  giving us a legitimate interpretation of the C-S-N tree.

This kind of strategy explains a number of other examples from Grinder. Grinder lists (13.5a)-(13.11a) as grammatical.

(13.5a) It was difficult for me to sketch myself.

- (13.6a) It was difficult for you to sketch yourself.
- (13.7a) It was difficult for him to sketch himself.
- (13.8a) It was difficult for her to sketch herself.
- (13.9a) It was difficult for us to sketch ourselves.
- (13.10a) It was difficult for you to sketch yourselves.
- (13.11a) It was difficult for them to sketch themselves.

After Equi-NP Deletion, Grinder lists only (13.5b), (13.6b), (13.9b), and (13.10b) as grammatical.

- (13.5b) It was difficult to sketch myself.
- (13.6b) It was difficult to sketch yourself.
- (13.7b) \*It was difficult to sketch himself.
- (13.8b) \*It was difficult to sketch herself.
- (13.9b) It was difficult to sketch ourselves.
- (13.10b) It was difficult to sketch yourselves.
- (13.11b) \*It was difficult to sketch themselves.

The probable reason this comes about is that we are used to thinking of I, you singular, us, and you plural as always being in focus, while referents for he, she, and they are ordinairly not in focus. Needless to say, if referents for he, she, or they are in focus, the situation changes completely. This is shown by  $(13.12)-(13.14).$ 

- (13.12a) Nurse Bob Breezy gave up drawing.
- (13.12b) [Bob] It was difficult to sketch himself.
- (13.13a) Astronaut Linda Smith gave up drawing.
- (13.13b) [Linda] It was difficult to sketch herself.
- (13.14a) The bank embezzlers gave up drawing.
- (13.14b) [the bank embezzlers] It was difficult to sketch themselves.

To indicate that various N-nodes are in focus, we bracket them at the beginning of a sentence. Thus  $(13.15)-(13.17)$  are not interpretable while  $(13.18)-$ (13.20) are.

- (13.15) \*It was difficult to sketch himself.
- (13.16) \*It was difficult to sketch herself.
- (13.17) \*It was difficult to sketch themselves.
- (13.18) [Bob] It was difficult to sketch himself.
- (13.19) [Linda] It was difficult to sketch herself.
- (13.20) [the bank embezzlers] It was difficult to sketch themselves.

The following examples involve resolution through focusing.

(13.21) It was difficult to sketch myself.

Features										
	<b>PNF</b>	FPF	<b>SPF</b>	TPF	PLF	<b>GNF</b>	<b>ANF</b>	<b>RPF</b>	<b>GEN</b>	
Ιo										
you0										
Φ		◠	っ	$\bigcap$	◠		◠			
myself										

chaining init\_table



```
init_table: exiting
chaining_n(myself)
     refl_chaining(myself)
           simplex_pred(myself)
           simplex_pred: \phichaining_n_to_n(myself, \phi)
                sc(myself, \phi) = True
                agr(myself, \phi) = True
                rnr(myself, \phi) = True
                chaining_e_to_n(myself<sub>a</sub>, \phi)
                      agr(myself<sub>a</sub>, \phi) = True
                      new_chain(myself<sub>a</sub>, \phi)
                           new_chain: create \phi_{\rm b}new_chain: create \stackrel{\sim}{\phi_{\rm D}}<sup>^</sup>myself<sub>a</sub>
```


```
new_chain: exiting
               chaining_e_to_n: exiting
          chaining_n_to_n: exiting
          simplex_pred(\phi)simplex_pred:
    refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non_refl\_changing(\phi)chaining_n_to_n(\phi, myself)
               sc(\phi, myself) = False
          chaining_n_to_n: exiting
          chaining_n_to_n(\phi, you_0)
               \text{sc}(\phi, \text{you}_0) = \text{True}agr(\phi, you_0) = Truernr(\phi, you<sup>0</sup>) = True
               chaining_e_to_n(\phi_a, you<sup>0</sup>)
                    aqr(\phi_a, you_0) = Truenew\_chain(\phi_a, you_0)new_chain: create you<sub>0b</sub>
                         new_chain: create you<sub>0b</sub>^\phia
```


```
new_chain: exiting
     chaining_e_to_n: exiting
     chaining_e_to_n(\phi_b, y \circ u_0)agr(\phi_b, you<sub>0</sub>) = False
     chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(\phi, I<sub>0</sub>)
     \text{sc}(\phi, I_0) = \text{True}agr(\phi, I_0) = Truernr(\phi, I<sub>0</sub>) = True
     chaining e_to_n(\phi_a, I_0)agr(\phi_a, I_0) = Truenew_chain(\phi_a, I<sub>0</sub>)
                new_chain: create I<sub>Ob</sub>
                new_chain: create I_{0b}^{\dagger} \phi_a
```


new\_chain: exiting chaining\_e\_to\_n: exiting chaining\_e\_to\_n( $\phi_b$ , I<sub>0</sub>)  $agr(\phi_b, I_0) = True$ new\_chain( $\phi_{\rm b}$ , I<sub>0</sub>) new\_chain: create I<sub>OC</sub> new\_chain: create  $I_{0c}^{\dagger}$ <sup>o</sup> $\phi_{b}$ 



```
new_chain: exiting
                  chaining_e_to_n: exiting
              chaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(you<sub>0</sub>)</sub>
         non_refl_chaining(you0)
              chaining_n_to_n(you<sub>0</sub>, myself)
                  sc(you<sub>0</sub>, myself) = Falsechaining_n_to_n: exiting
              chaining_n_to_n(you_0, \phi)
                  sc(you<sub>0</sub>, \phi) = False
              chaining_n_to_n: exiting
              chaining n_to_n(you_0, I_0)sc(you_0, I_0) = Trueagr(you_0, I_0) = Falsechaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I_0)
         non_refl\_chaining(I_0)chaining_n_to_n(I<sub>0</sub>, myself)
                  sc(I<sub>0</sub>, myself) = Falsechaining_n_to_n: exiting
              chaining_n_to_n(I_0, \phi)
                  \text{sc}(I_0, \phi) = \text{False}chaining_n_to_n: exiting
              chaining_n_to_n(I_0, you<sub>0</sub>)
                  sc(I_0, you_0) = Falsechaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```




(13.22) [toy] Give me that!

Features										
	<b>PNF</b>	FPF	<b>SPF</b>	TPF	PLF	<b>GNF</b>	<b>ANF</b>	<b>RPF</b>	<b>GEN</b>	
$I_0$	$^{+}$	$^{+}$				っ	$\pm$			
you0	$^{+}$		$\div$		7	ာ	$^{+}$			
toy						っ	P			
Φ	$\, +$	?	っ	7	7	ာ	?			
me	$^+$	$^{+}$				2	+			
that						2	P			

```
chaining
    init_table
```


```
init_table: exiting
chaining_n(that)
    non_refl_chaining(that)
        chaining_n_to_n(that, me)
            sc(that, me) = True
            agr(hat, me) = Falsechaining_n_to_n: exiting
        chaining_n_to_n(that, \phi)
            sc(that, \phi) = True
            agr(that, \phi) = True
            rnr(that, \phi) = False
        chaining_n_to_n: exiting
        chaining_n_to_n(that, toy)
            sc(that, toy) = Trueagr(that, toy) = True
            rnr(that, toy) = True
            chaining_e_to_n(thata, toy)
                 agr(hat_a, toy) = Truenew_chain(thata, toy)
                     new_chain: create toyb
                     new_chain: create toy<sub>b</sub>^thata
```


```
new_chain: exiting
              chaining_e_to_n: exiting
         chaining_n_to_n: exiting
         chaining n_to_n (that, you_0)
              sc(that, you<sub>0</sub>) = True
              agr(hat, you<sub>0</sub>) = Falsechaining_n_to_n: exiting
         chaining n_to_n (that, I_0)
              sc(that, I_0) = True
              agr(that, I_0) = False
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(me)
    non_refl_chaining(me)
         chaining_n_to_n(me, that)
              sc(me, that) = Falsechaining_n_to_n: exiting
         chaining n_to_n (me, \phi)
              sc(me, \phi) = True
              agr(me, \phi) = True
              rnr(me, \phi) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(me, toy)
              sc(me, toy) = True
              agr(me, toy) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(me, you<sub>0</sub>)
              sc(me, you<sub>0</sub>) = True
              agr(me, you<sub>0</sub>) = False
         chaining_n_to_n: exiting
         chaining n_to_n (me, I_0)
              sc(me, I_0) = True
             agr(me, I_0) = True
              rnr(me, I_0) = True
              chaining_e_to_n(me<sub>a</sub>, I_0)
                  agr(me<sub>a</sub>, I_0) = True
                  new_chain(me<sub>a</sub>, I_0)
                       new_chain: create I<sub>Ob</sub>
                       new_chain: create I<sub>Ob</sub>^mea
```


```
new_chain: exiting
              chaining_e_to_n: exiting
         chaining_n_to_n: exiting
    non_refl_chaining: exiting
chaining_n: exiting
chaining_n(\phi)non_refl\_changing(\phi)chaining_n_to_n(\phi, that)
              sc(\phi, that) = False
         chaining_n_to_n: exiting
         chainingn_to_n(\phi, me)sc(\phi, me) = False
         chaining_n_to_n: exiting
         chaining_n_to_n(\phi, toy)
              sc(\phi, toy) = True
              agr(\phi, toy) = Truernr(\phi, toy) = True
              chaining_e_to_n(\phi_a, toy)
                  agr(\phia, toy) = True
                  new_chain(\phi_a, toy)
                       new_chain: create toy<sub>c</sub>
                       new_chain: create toy<sub>c</sub>^\phi<sub>a</sub>
```


```
new_chain: exiting
     chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(\phi, you_0)
     \text{sc}(\phi, \text{you}_0) = \text{True}agr(\phi, you_0) = Truernr(\phi, you<sup>0</sup>) = True
     chaining_e_to_n(\phi_a, you<sup>0</sup>)
          agr(\phi_a, you_0) = Truenew\_chain(\phi_a, you_0)new_chain: create you0b
               new_chain: create you<sub>0b</sub>^\phia
```


```
new_chain: exiting
     chaining_e_to_n: exiting
chaining_n_to_n: exiting
chaining_n_to_n(\phi, I<sub>0</sub>)
     sc(\phi, I_0) = True
     agr(\phi, I_0) = Truernr(\phi, I<sub>0</sub>) = True
     chaining_e_to_n(\phi_a, I_0)agr(\phi_a, I_0) = True
          new_chain(\phi_a, I<sub>0</sub>)
               new_chain: create I_{0c}new_chain: create I_{0c}<sup>\phi</sup>a
```


```
new_chain: exiting
                  chaining_e_to_n: exiting
             chaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(you<sub>0</sub>)</sub>
         non refl chaining(you_0)
             chaining_n_to_n(you<sub>0</sub>, that)
                  sc(you<sub>0</sub>, that) = Falsechaining_n_to_n: exiting
             chaining n_to_n(you_0, me)sc(you_0, me) = Falsechaining_n_to_n: exiting
             chaining_n_to_n(you_0, \phi)
                  sc(you<sub>0</sub>, \phi) = False
             chaining_n_to_n: exiting
             chaining_n_to_n(you<sub>0</sub>, toy)
                  sc(you<sub>0</sub>, toy) = Falsechaining_n_to_n: exiting
             chaining n_to_n(you_0, I_0)sc(you_0, I_0) = Trueagr(you<sub>0</sub>, I<sub>0</sub>) = Falsechaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
    chaining_n(I_0)non_refl\_changing(I_0)chaining_n_to_n(I_0, that)
                  sc(I<sub>0</sub>, that) = Falsechaining_n_to_n: exiting
             chaining_n_to_n(I_0, me)
                  sc(I_0, me) = Falsechaining_n_to_n: exiting
             chaining_n_to_n(I_0, \phi)
                  \text{sc}(I_0, \phi) = \text{False}chaining_n_to_n: exiting
             chaining_n_to_n(I_0, toy)
                  sc(I_0, toy) = Falsechaining_n_to_n: exiting
             chaining_n_to_n(I_0, you_0)
                  sc(I_0, you_0) = Falsechaining_n_to_n: exiting
         non_refl_chaining: exiting
    chaining_n: exiting
chaining: exiting
```




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